Probabilistic Multi-agent Only-believing

Qihui Feng, Gerhard Lakemeyer RWTH Aachen University

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- **2** Logic OBL_m
- 3 Properties of the Logic

4 Conclusion

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Introduction

Knowledge and Belief:

- $K(fair(Coin) \land \neg fair(Die))$
- *B*(*fair*(Coin): 0.8)

Levesque proposed only-knowing to precisely capture the (non-)beliefs:

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- $O(fair(Coin)) \models \neg K(fair(Die)) \land \neg K(\neg fair(Die))$
- $O(fair(Coin)) \models \neg B(fair(Die): r)$ for any $r \in [0, 1]$

Research on only-knowing:

- Probabilistic only-believing: The logic \mathcal{OBL}
- Projection reasoning: $O(KB_1) \rightarrow [action]O(KB_2)$
- Multi-agent only-knowing:
 - Previous works in both propositional and first-order cases
 - No first-order account faithfully follows Levesque's principle of only-knowing.

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Levesque's notion of only-knowing: Given the only-knowing of the agent, any subjective formula will either be inferred or disproved.

 $O(\mathsf{KB}) \models \mathbf{K}\beta \text{ iff } \mathsf{KB} \models \beta; \quad O(\mathsf{KB}) \models \neg \mathbf{K}\beta \text{ iff } \mathsf{KB} \nvDash \beta.$

Desiderata for multi-agent extension:

i) non-beliefs on irrelevant items:

 $O_a(\neg fair(Coin) \land K_b(fair(Coin))) \models \neg K_a K_b(fair(Die))$

ii) non-beliefs on mental states with deeper nesting

 $O_a(\neg fair(Coin) \land K_b(fair(Coin))) \models \neg K_a K_b K_a(fair(Coin))$

To model only-knowing up to all depths is semantically difficult.

To model only-knowing up to depth k?

• Modality $O_a^{(k)}$: agent *a*'s only-knowing(believing) up to depth *k*.

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The new desiderata: for $K_a\beta$ with depth no more than k,

Either
$$O_a^{(k)} \alpha \models K_a \beta$$
 or $O_a^{(k)} \alpha \models \neg K_a \beta$?



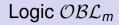


2 Logic OBL_m

3 Properties of the Logic

4 Conclusion

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A first-order modal logic with equality and featured with:

• A finite set of agents, e.g. $Ag = \{a, b\};$

Modalities for belief and only-believing for each agent.

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- Standard FO formulae;
- **B**_{*i*}(α : *r*): α is believed by agent *i* with degree *r* (written $K_i \alpha$ if r = 1).
- $O_i^{(k)}(\alpha : r)$: all that agent *i* believes up to depth *k* is α with degree *r*.
 - $O_a^{(1)}(fair(Coin))$: agent a knows fair(Coin) and nothing else about the world.
 - O_a⁽²⁾(fair(Coin)): agent a knows fair(Coin) and nothing else about the world, and nothing about Bob's beliefs about the world.

A model is a tuple (w, e_a, e_b) with world w and epistemic states e_a and e_b . A **world** $w \in W$ is a set of ground atoms.

• $w, e_a, e_b \models fair(Coin)$ iff $fair(Coin) \in w$

Epistemic states are defined inductively:

- 1-distribution assigns each world a probability: $W \to \mathbb{R}_{[0,1]}$;
- 1-epistemic state is a set of 1-distributions.

Example

$$w_1 = \{ fair(Coin) \} \text{ and } w_2 = \{ fair(Die) \}, d^1(w) = \begin{cases} 0.5 & w \in \{w_1, w_2\} \\ 0 & otherwise. \end{cases}$$

Let $e_a = \{ d^1 \}$, then $w, e_a, e_b \models B_a(fair(Coin): 0.5)$

Semantics (Nested Beliefs)

 \mathcal{E}^1 denotes the set of all 1-epistemic states. For any k > 1,

- *k*-distribution $d^k : (\mathcal{W} \times \mathcal{E}^{k-1}) \to \mathbb{R}_{[0,1]}$
- a *k*-epistemic state is a set of *k*-distributions

Example

Let
$$w_1 = \{ fair(Coin) \}, w_2 = \{ fair(Die) \}, w_3 = \{ fair(Coin), fair(Die) \}.$$

 $\tilde{d}^1(w) = \begin{cases} 0.5 & w \in \{w_1, w_3\} \\ 0 & otherwise. \end{cases}$
 $d^2(w, e_b^1) = \begin{cases} 0.3 & w = w_1, e_b^1 = \{\tilde{d}^1\} \\ 0.7 & w = w_2, e_b^1 = \{\tilde{d}^1\} \\ 0 & otherwise. \end{cases}$
Let $e_a = \{ d^2 \}$ then

 $w, e_a, e_b \models B_a(fair(Coin): 0.3)$ $w, e_a, e_b \models K_a(K_b(fair(Coin)))$

Suppose that $e_a \in \mathcal{E}^2$ (2-epistemic state)

• $w, e_a, e_b \models O_a^{(2)}(\neg fair(Coin) \land K_b(fair(Coin)))$ iff for any 2-distribution d,

$$d \in e_a \iff w, \{d\}, e_b \models K_a(\neg fair(Coin) \land K_b(fair(Coin)))$$

i.e. there is a **unique** $e_a \in \mathcal{E}^2$ which satisfies $O_a^{(2)}(\neg fair(\text{Coin}) \land K_b(fair(\text{Coin})))$ Every $e_a \in \mathcal{E}^k$ can be **uniquely** "reduced" to an $e'_a \in \mathcal{E}^{k-1}$ s.t.

 $w, e_a, e_b \models \alpha$ iff $w, e'_a, e_b \models \alpha$ for any α not deeper than k-1

For $e_a \in \mathcal{E}^3$, $w, e_a, e_b \models O_a^{(2)} (\neg fair(\text{Coin}) \land K_b(fair(\text{Coin})))$ iff e_a reduce to $e'_a \in \mathcal{E}^2$ and $w, e'_a, e_b \models O_a^{(2)} (\neg fair(\text{Coin}) \land K_b(fair(\text{Coin})))$

Compatibility:

- Formulae like $O_a^{(1)}(\neg fair(Coin) \land K_b(fair(Coin)))$ are illegal.
- *e compatible* with α : the depth of *e* is not less than the depth of α

We say Σ entails α (written $\Sigma \models \alpha$) iff:

For each model (w, e_a, e_b) compatible with Σ, α , if $(w, e_a, e_b) \models \sigma$ for all $\sigma \in \Sigma$, then $(w, e_a, e_b) \models \alpha$

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We say α is valid iff $\{\} \models \alpha$



1 Introduction

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Properties of Knowledge

 \mathcal{OBL}_m follows the KD45_n properties. For agent $i \in Ag$,

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• (Nec) If $\models \alpha$, then $\models \mathbf{K}_i \alpha$

$$\blacksquare (\mathsf{K}) \models \mathbf{K}_i \alpha \land \mathbf{K}_i (\alpha \supset \beta) \supset \mathbf{K}_i \beta$$

$$\blacksquare (\mathsf{D}) \models K_i \alpha \supset \neg K_i \neg \alpha$$

- $\blacksquare (4) \models K_i \alpha \supset K_i K_i \alpha$
- $\blacksquare (5) \models \neg K_i \alpha \supset K_i \neg K_i \alpha$
- **\mathbf{K}_{i} \alpha \wedge \neg \alpha is satisfiable**

Barcan formulae:

$$\blacksquare \models \forall x. K_i \alpha \supset K_i \forall x. \alpha$$

$$\blacksquare \models \exists \mathbf{X}. \mathbf{K}_i \alpha \supset \mathbf{K}_i \exists \mathbf{X}. \alpha$$

Properties of Beliefs

The degree of belief follows the laws of probability:

$$\blacksquare \models B_i(\alpha \colon r) \supset \neg B_i(\alpha \colon r') \text{ for } r' \neq r$$

$$\blacksquare \models B_i(\alpha \colon r) \supset B_i(\neg \alpha \colon 1 - r)$$

$$\blacksquare \models B_i(\alpha \land \beta \colon r) \land B_i(\alpha \land \neg \beta \colon r') \supset B_i(\alpha \colon r+r')$$

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Only-Believing

Modality $O_a^{(k)}$ precisely captures agent *a*'s beliefs and non-beliefs up to depth *k*.

- non-beliefs on irrelevant items: $O_a^{(2)}(\neg fair(Coin) \land K_b(fair(Coin))) \models \neg K_a K_b(fair(Die))$
- non-beliefs on deeper mental states:

$$O_a^{(2)}(\neg fair(Coin) \wedge K_b(fair(Coin))) \not\models \neg K_a K_b K_a(\neg fair(Coin))$$

$$O^{(3)}_aigl(\neg \mathit{fair}(\mathsf{Coin}) \wedge oldsymbol{K}_b(\mathit{fair}(\mathsf{Coin}))igr) \models \neg oldsymbol{K}_aoldsymbol{K}_boldsymbol{K}_a(\neg \mathit{fair}(\mathsf{Coin})igr)$$

For $i \in Ag$, given *i*-objective formulae α and β s.t. the depth of $K_i(\beta)$ not greater than k, then $O_i^{(k)}(\alpha)$ entails either $K_i\beta$ or $\neg K_i\beta$.

 OBL_m can represent defaults about another agent's beliefs:

Example

Let KB = { \neg fair(Coin)}, $\delta = \forall r.(r \neq 0 \supset \neg B_a(\neg K_b(fair(Coin)): r)) \supset K_b(fair(Coin))$ Bob believes fair(Coin) unless otherwise (Bob does not believes fair(Coin)) is believed (by Alice) with a non-zero degree

Conclusion

In this work, we

- propose an logical account for multi-agent only-believing
- prove properties on beliefs and only-believing
- explore the capability of default reasoning about nested beliefs

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For future work:

- extend to belief after actions \checkmark
- develop mechanisms for projection reasoning
- join common beliefs and only-believing

Thank you!