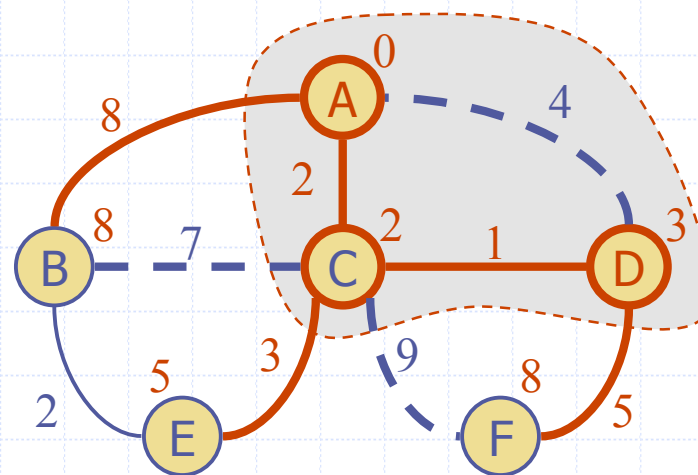
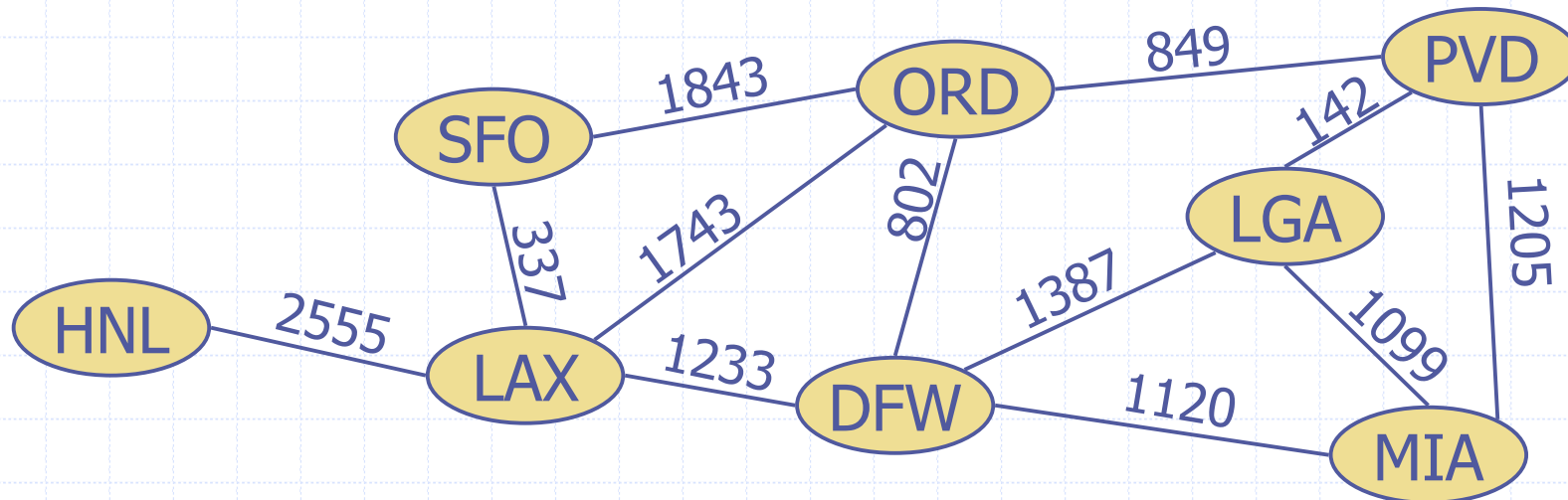


# Shortest Paths



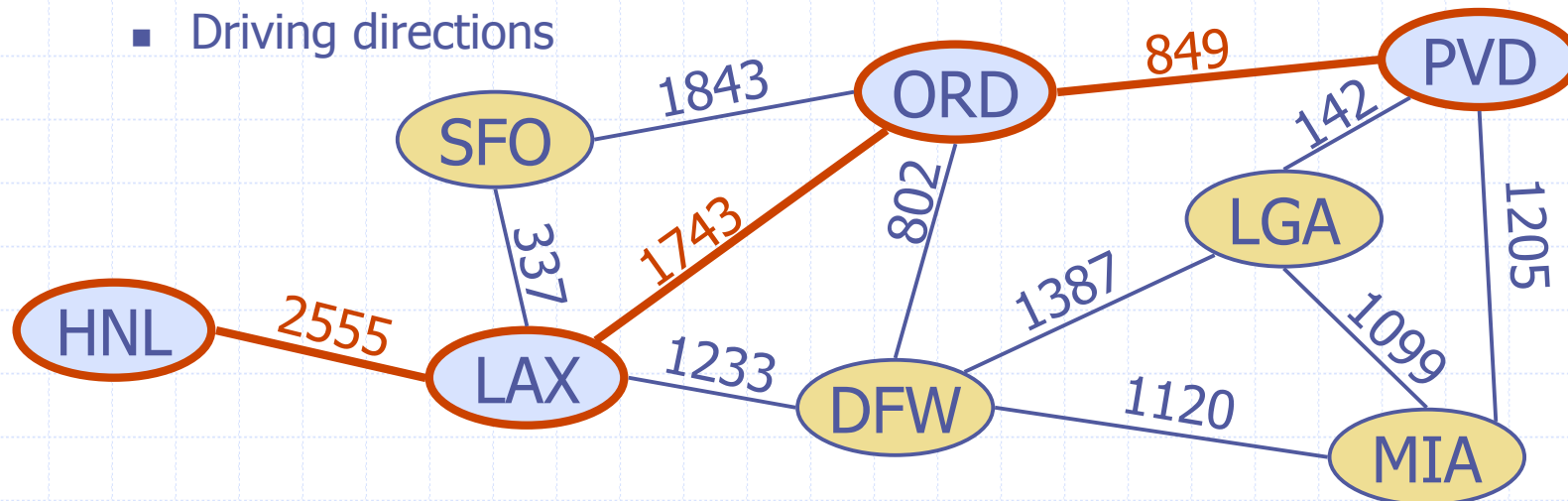
# Weighted Graphs (§ 12.5)

- ◆ In a weighted graph, each edge has an associated numerical value, called the weight of the edge
- ◆ Edge weights may represent, distances, costs, etc.
- ◆ Example:
  - In a flight route graph, the weight of an edge represents the distance in miles between the endpoint airports



# Shortest Paths (§ 12.6)

- ◆ Given a weighted graph and two vertices  $u$  and  $v$ , we want to find a path of minimum total weight between  $u$  and  $v$ .
  - Length of a path is the sum of the weights of its edges.
- ◆ Example:
  - Shortest path between Providence and Honolulu
- ◆ Applications
  - Internet packet routing
  - Flight reservations
  - Driving directions



# Shortest Path Properties

## Property 1:

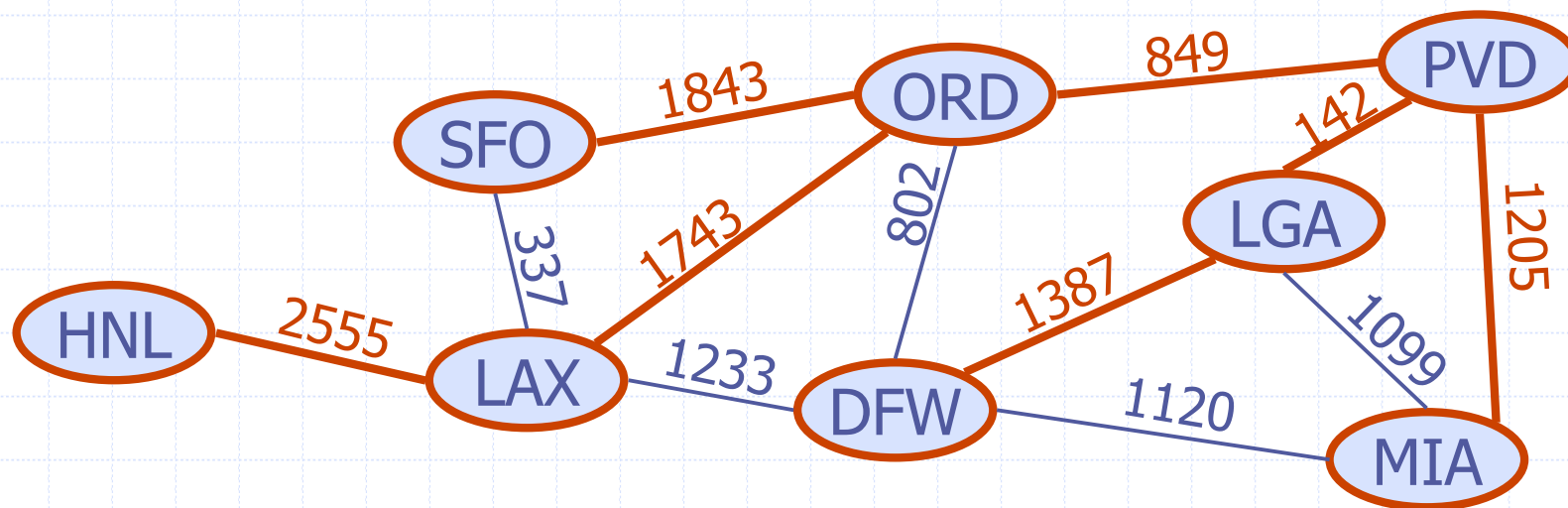
A subpath of a shortest path is itself a shortest path

## Property 2:

There is a tree of shortest paths from a start vertex to all the other vertices

## Example:

Tree of shortest paths from Providence

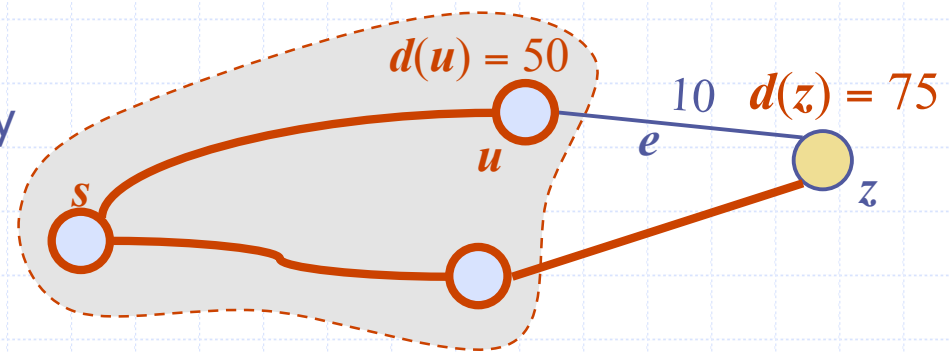


# Dijkstra's Algorithm (§ 12.6.1)

- ◆ The distance of a vertex  $v$  from a vertex  $s$  is the length of a shortest path between  $s$  and  $v$
- ◆ Dijkstra's algorithm computes the distances of all the vertices from a given start vertex  $s$
- ◆ Assumptions:
  - the graph is connected
  - the edges are undirected
  - the edge weights are **nonnegative**
- ◆ We grow a "**cloud**" of vertices, beginning with  $s$  and eventually covering all the vertices
- ◆ We store with each vertex  $v$  a label  $d(v)$  representing the distance of  $v$  from  $s$  in the subgraph consisting of the cloud and its adjacent vertices
- ◆ At each step
  - We add to the cloud the vertex  $u$  outside the cloud with the smallest distance label,  $d(u)$
  - We update the labels of the vertices adjacent to  $u$

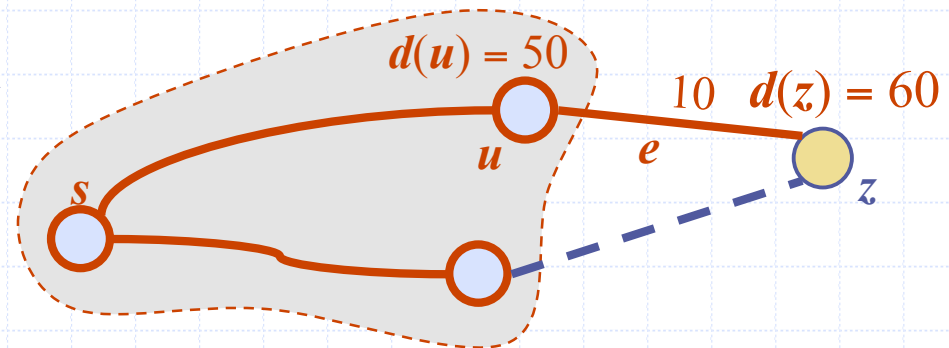
# Edge Relaxation

- ◆ Consider an edge  $e = (u, z)$  such that
  - $u$  is the vertex most recently added to the cloud
  - $z$  is not in the cloud

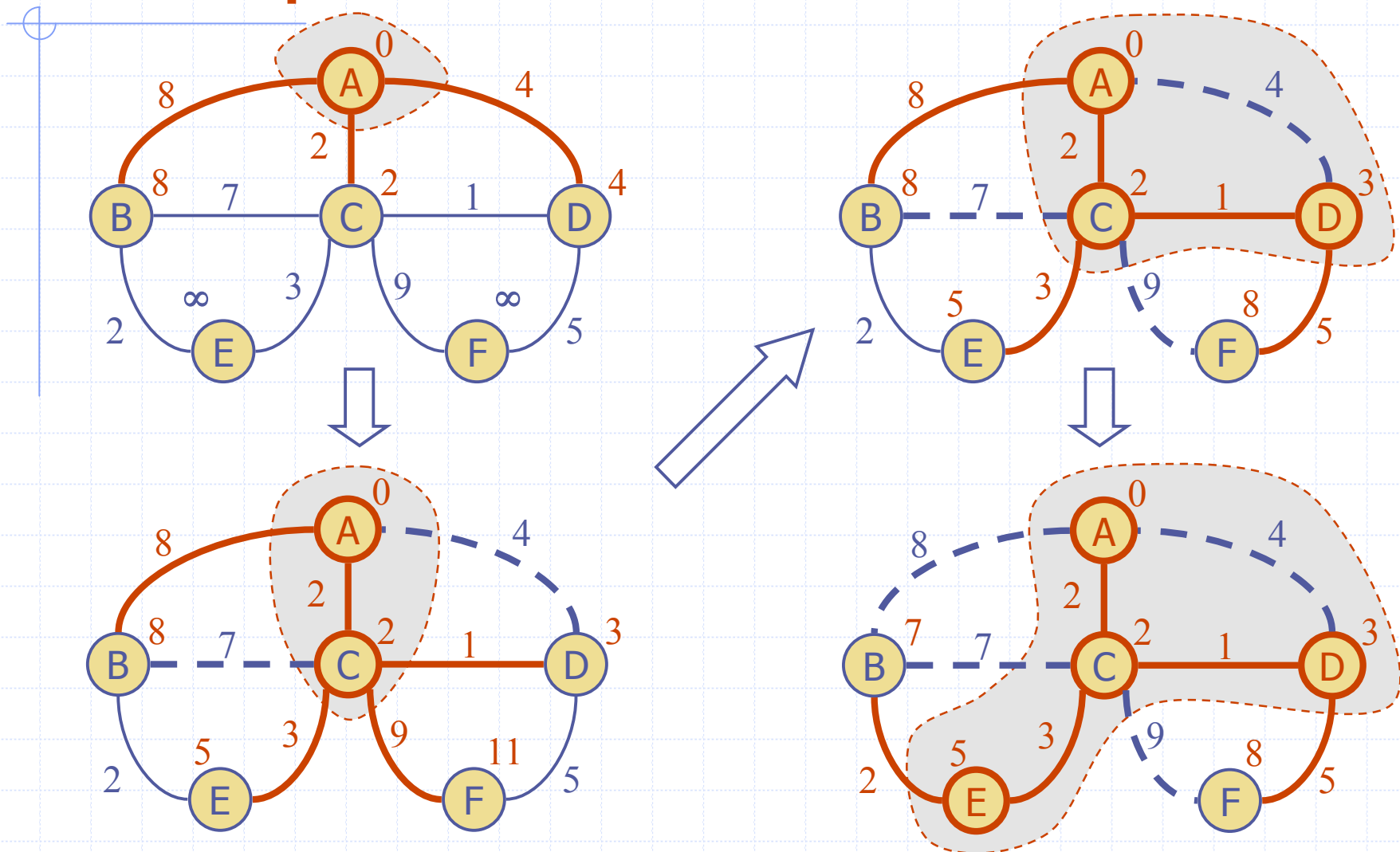


- ◆ The relaxation of edge  $e$  updates distance  $d(z)$  as follows:

$$d(z) \leftarrow \min\{d(z), d(u) + \text{weight}(e)\}$$



# Example



# Example (cont.)

