COMP2111 Week 9
Term 1, 2023
Hoare Logic
Summary

- $\mathcal{L}$: A simple imperative programming language
- Hoare triples (SYNTAX)
- Hoare logic (PROOF)
- Semantics for Hoare logic
- Handling termination
- Adding non-determinism
Aims

We’ve seen how to use Hoare logic to verify programs.

But how do we know that Hoare logic works? Do we need to take the rules on faith? Or can we prove that it works?
We’ve seen how to use Hoare logic to verify programs.

But how do we know that Hoare logic works? Do we need to take the rules on faith? Or can we prove that it works?

We’ve already asked (and answered) a similar question about a different logic (natural deduction).
Informal semantics

Hoare logic gives a proof of $\{\varphi\} P \{\psi\}$, that is: $\vdash \{\varphi\} P \{\psi\}$ (axiomatic semantics)

What does it mean for $\{\varphi\} P \{\psi\}$ to be valid, that is: $|= \{\varphi\} P \{\psi\}$?
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What does it mean for \( \{ \varphi \} P \{ \psi \} \) to be valid, that is:
\( \models \{ \varphi \} P \{ \psi \} \)?

We need a semantics for \( \mathcal{L} \).
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We need a semantics for $L$.

We could use the LTS semantics of $L$ from Week 8. We will use a denotational style instead, similar to Assignment 1 Problem 1 but systematic.
Informal semantics: Programs

We know (from Assignment 1 Problem 1) that programs can be modelled as *relations* between initial and final states.
Informal semantics: States

What is a state?

Two approaches:

Concrete: from a physical perspective
States are memory configurations, register contents, etc.
Store of variables and the values associated with them

Abstract: from a mathematical perspective
The pre-/postcondition predicates hold
⇒ States are logical interpretations (Model + Environment)
There is only one model of interest: standard interpretations of arithmetical symbols
⇒ States are fully determined by environments
⇒ States are functions that map variables to values
Informal semantics: States

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  - States are functions that map variables to values
Informal semantics: States

State space ($Env$)

- $x \leftarrow 0$
- $y \leftarrow 0$
- $z \leftarrow 0$
- $x \leftarrow 3$
- $y \leftarrow 2$
- $z \leftarrow 1$
- $x \leftarrow 1$
- $y \leftarrow 1$
- $z \leftarrow 1$
- $x \leftarrow 2$
- $y \leftarrow 2$
- $z \leftarrow 2$
- $x \leftarrow 0$
- $y \leftarrow 1$
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Informal semantics: States and Programs

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Informal semantics: States and Programs
Semantics for $L$

An environment or state is a function from variables to (numeric) values. We denote by $\text{Env}$ the set of all environments.

**NB**

An environment, $\eta$, assigns a numeric value $[e]^{\eta}$ to all expressions $e$, and a boolean value $[b]^{\eta}$ to all boolean expressions $b$. 
Semantics for $\mathcal{L}$

An **environment** or **state** is a function from variables to (numeric) values. We denote by $\text{Env}$ the set of all environments.

**NB**

An environment, $\eta$, assigns a numeric value $[e]^\eta$ to all expressions $e$, and a boolean value $[b]^\eta$ to all boolean expressions $b$.

Given a program $P$ of $\mathcal{L}$, we define $[P]$ to be a **binary relation** on $\text{Env}$ in the following manner...
Assignment

\[(\eta, \eta') \in [x := e] \text{ if, and only if } \eta' = \eta[x \mapsto [e]^{\eta}]\]
Assignment: \([z := 2]\)
Recall

If \( R \) and \( S \) are binary relations, then the \textbf{relational composition} of \( R \) and \( S \), \( R; S \) is the relation:

\[
R; S := \{(a, c) : \exists b \text{ such that } (a, b) \in R \text{ and } (b, c) \in S\}
\]

If \( R \subseteq A \times B \) is a relation, and \( X \subseteq A \), then the \textbf{image of} \( X \) \textbf{under} \( R \), \( R(X) \) is the subset of \( B \) defined as:

\[
R(X) := \{b \in B : \exists a \in X \text{ such that } (a, b) \in R\}.
\]
Sequencing

\[ [P; Q] = [P]; [Q] \]

where, on the RHS, ; is relational composition.
Conditional, first attempt

\[
\text{[if } b \text{ then } P \text{ else } Q \text{ fi]} = \begin{cases} 
 [P] & \text{if } [b]^{\eta} = \text{true} \\
 [Q] & \text{otherwise.}
\end{cases}
\]
Conditional, first attempt

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\text{[if } b \text{ then } P \text{ else } Q \text{ fi]} = \begin{cases} 
[P] & \text{if } [b]^\eta = \text{true} \\
[Q] & \text{otherwise.}
\end{cases}
\]

We’d like to avoid mentioning \( \eta \) on the LHS, so this won’t do.
Detour: Predicates as programs

A boolean expression $b$ defines a subset (or unary relation) of \( \text{Env} \):

\[
\langle b \rangle = \{ \eta : [b]^\eta = \text{true} \}
\]

This can be extended to a binary relation (i.e. a program):

\[
[b] = \{ (\eta, \eta) : \eta \in \langle b \rangle \} 
\]
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\[
[b] = \{ (\eta, \eta) : \eta \in \langle b \rangle \}
\]

Intuitively, \( b \) corresponds to the program

\[
\text{if } b \text{ then skip else abort fi}
\]
Conditional, better attempt

\[
\left[ \text{if } b \text{ then } P \text{ else } Q \text{ fi} \right] = \left[ b; P \right] \cup \left[ \neg b; Q \right]
\]
While

while \( b \) do \( P \) od

- Do 0 or more executions of \( P \) while \( b \) holds
- Terminate when \( b \) does not hold
While

while $b$ do $P$ od

- Do 0 or more executions of $(b; P)$
- Terminate with an execution of $\neg b$
While

while $b$ do $P$ od

- Do 0 or more executions of $(b; P)$
- Terminate with an execution of $\neg b$

How to do “0 or more” executions of $(b; P)$?
Reflexive and transitive closure

Given a binary relation \( R \subseteq E \times E \), the transitive closure of \( R \), \( R^* \) is defined inductively by the following rules:

\[
\begin{align*}
\text{If } x & \in E, \\
\text{then } x R^* x \\
\end{align*}
\]

\[
\begin{align*}
\text{If } x R y \\
\text{and } y R^* z, \\
\text{then } x R^* z \\
\end{align*}
\]

NB

\( R; R^* \subseteq R^* \).
While

\[ [\text{while } b \text{ do } P \text{ od}] = [b; P]^*; [\neg b] \]

- Do 0 or more executions of \((b; P)\)
- Conclude with an execution of \(\neg b\)
A Hoare triple is **valid**, written $\models \{ \varphi \} P \{ \psi \}$ if

$$\llbracket P \rrbracket (\langle \varphi \rangle) \subseteq \langle \psi \rangle.$$ 

That is, the relational image under $\llbracket P \rrbracket$ of the set of states where $\varphi$ holds is contained in the set of states where $\psi$ holds.
Validity
Validity
Validity
Validity

\[ \langle \varphi \rangle \quad [P] \quad \langle \psi \rangle \]
Validity

\[ [P] \]

\[
\langle \varphi \rangle \quad [P](\langle \varphi \rangle) \quad \langle \psi \rangle
\]
Soundness of Hoare Logic

Theorem

\[ \text{If } \vdash \{ \varphi \} P \{ \psi \} \text{ then } \models \{ \varphi \} P \{ \psi \} \]
Incompleteness

Theorem (Gödel’s Incompleteness Theorem)

There is no proof system that can prove every valid first-order sentence about arithmetic over the natural numbers.
Incompleteness

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*There is no proof system that can prove every valid first-order sentence about arithmetic over the natural numbers.*

⇒ There are true statements that do not have a proof.
Incompleteness

Theorem (Gödel’s Incompleteness Theorem)

There is no proof system that can prove every valid first-order sentence about arithmetic over the natural numbers.

⇒ There are true statements that do not have a proof.

⇒ Because of (cons) there are valid triples that result from valid, but unprovable, consequences.
Incompleteness

Theorem (Gödel’s Incompleteness Theorem)

There is no proof system that can prove every valid first-order sentence about arithmetic over the natural numbers.

⇒ There are true statements that do not have a proof.

⇒ Because of (cons) there are valid triples that result from valid, but unprovable, consequences.

⇒ Hoare Logic is not complete.
Relative completeness of Hoare Logic

Theorem (Relative completeness of Hoare Logic)

With an oracle that decides the validity of predicates,

\[ \text{if } \models \{ \varphi \} P \{ \psi \} \text{ then } \vdash \{ \varphi \} P \{ \psi \}. \]

Intuitively: Hoare logic is no more incomplete than the logic used to express the pre- and postconditions.
Summary

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Termination

Hoare triples for partial correctness:

\[ \{ \varphi \} \; P \; \{ \psi \} \]

Asserts \( \psi \) holds if \( P \) terminates.

That’s just a safety property. Let’s add liveness!
Termination

Hoare triples for partial correctness:

\[ \{ \varphi \} \text{ } P \{ \psi \} \]

Asserts \( \psi \) holds if \( P \) terminates.

That’s just a safety property. Let’s add liveness!

Hoare triples for total correctness:

\[ [\varphi] \text{ } P \ [\psi] \]

Asserts:
If \( \varphi \) holds at a starting state, and \( P \) is executed;
then \( P \) will terminate and \( \psi \) will hold in the resulting state.
Warning

Termination is hard!

- Algorithmic limitations (e.g. Halting problem)
Warning

Termination is hard!

- Algorithmic limitations (e.g. Halting problem)
- Mathematical limitations

Example

<table>
<thead>
<tr>
<th>COLLATZ</th>
</tr>
</thead>
<tbody>
<tr>
<td>while $n &gt; 1$ do</td>
</tr>
<tr>
<td>if $n % 2 = 0$</td>
</tr>
<tr>
<td>then</td>
</tr>
<tr>
<td>$n := n/2$</td>
</tr>
<tr>
<td>else</td>
</tr>
<tr>
<td>$n := 3 \times n + 1$</td>
</tr>
<tr>
<td>fi</td>
</tr>
<tr>
<td>od</td>
</tr>
</tbody>
</table>

Example of a problem with termination: Collatz Conjecture.
Total correctness

How can we show:

$$[(m \geq 0) \land (n > 0)] \text{Pow } [r = n^m]$$?
Total correctness

How can we show:

\[ ((m \geq 0) \land (n > 0)) \text{ Pow } [r = n^m]? \]

Use **Hoare Logic for total correctness:**

- (ass), (seq), (cond), and (cons) rules all the same
- Modified (loop) rule
Rules for total correctness

\[
\frac{[\varphi[e/x]] x := e [\varphi]}{(\text{ass})}
\]

\[
\frac{[\varphi] P [\psi] \quad [\psi] Q [\rho]}{[\varphi] P; Q [\rho]} (\text{seq})
\]

\[
\frac{[\varphi \land g] P [\psi] \quad [\varphi \land \neg g] Q [\psi]}{[\varphi] \text{ if } g \text{ then } P \text{ else } Q \text{ fi } [\psi]} (\text{if})
\]

\[
\frac{\varphi' \rightarrow \varphi \quad [\varphi] P [\psi] \quad \psi \rightarrow \psi'}{[\varphi'] P [\psi']} (\text{cons})
\]
Terminating while loops

{φ} while b do P od {ψ}

Partial correctness:
Find an invariant I such that:

- φ → I
- \{I ∧ b\} P \{I\}
- (I ∧ ¬b) → ψ
Terminating while loops

\[ [\varphi] \text{ while } b \text{ do } P \text{ od } [\psi] \]

**Partial correctness:**
Find an invariant \( I \) such that:
- \( \varphi \rightarrow I \) (establish)
- \( [I \land b] P [I] \) (maintain)
- \( (I \land \neg b) \rightarrow \psi \) (conclude)

**Show termination:**
Find a **variant** \( v \) such that:
- \( (I \land b) \rightarrow v > 0 \) (positivity)
- \( [I \land b \land v = N] P [v < N] \) (progress)
Loop rule for total correctness

\[
\begin{align*}
[\varphi \land g \land (v = N)] & \quad P \quad [\varphi \land (v < N)] \\
(\varphi \land g) \rightarrow (v > 0) & \\
[\varphi] & \text{while } g \text{ do } P \text{ od } [\varphi \land \neg g]
\end{align*}
\] (loop)
## Termination for Pow

<table>
<thead>
<tr>
<th>Pow</th>
<th>${\text{init: } (m \geq 0) \land (n &gt; 0)}$</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>${(1 = n^0) \land (0 \leq m) \land \text{init}}$</td>
</tr>
<tr>
<td></td>
<td>${(r = n^0) \land (0 \leq m) \land \text{init}}$</td>
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</tbody>
</table>

$r := 1;$  
$i := 0;$  

while $i < m$ do  

$r := r \times n;$  
$i := i + 1$  

end  

What is a suitable variant?
Termination for Pow

Pow

\[ \text{init}: (m \geq 0) \land (n > 0) \]
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\[ r := 1; \]
\[ i := 0; \]

while \( i < m \) do

\[ (r \times n = n^{i+1}) \land (i + 1 \leq m) \land \text{init} \]
\[ (r = n^{i+1}) \land (i + 1 \leq m) \land \text{init} \]

\[ r := r \times n; \]
\[ i := i + 1 \]
od

\[ \text{Inv} \land (i \geq m) \]
\[ r = n^m \]

What is a suitable variant? \( v := (m - i) \)
**Termination for Pow**

<table>
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| \[
\begin{align*}
\text{init: } & (m \geq 0) \land (n > 0) \\
((1 = n^0) \land (0 \leq m)) \land \text{init} \\
((r = n^0) \land (0 \leq m)) \land \text{init} \\
\{r := 1; i := 0; \} \\
\text{while } i < m \text{ do } \{r := r \ast n; i := i + 1 \} \{r = n^{i+1} \land (i + 1 \leq m) \land \text{init}\} \\
\text{od} \{r = n^m\} \\
\end{align*}
\] |

What is a suitable variant? \( v := (m - i) \)
Termination for Pow

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<td>i := 0;</td>
</tr>
<tr>
<td>while i &lt; m do</td>
</tr>
<tr>
<td>r := r * n;</td>
</tr>
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<td>i := i + 1</td>
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{init: (m ≥ 0) ∧ (n > 0)}
{(1 = n^0) ∧ (0 ≤ m) ∧ init}
{(r = n^0) ∧ (0 ≤ m) ∧ init}
{Inv}
{(r * n = n^{i+1}) ∧ (i + 1 ≤ m) ∧ init}
{(r = n^{i+1}) ∧ (i + 1 ≤ m) ∧ init}
{Inv ∧ (v = N)}
{Inv ∧ (i < m) ∧ (v = N)}
{Inv ∧ (v < N)}
{Inv ∧ (i ≥ m)}
{r = n^m}
### Termination for Pow

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\begin{align*}
\text{init: } & (m \geq 0) \land (n > 0) \\
\{(1 = n^0) \land (0 \leq m) \land \text{init}\} \\
\{(r = n^0) \land (0 \leq m) \land \text{init}\} \\
\{\text{Inv}\} \\
\{(r \times n = n^{i+1}) \land (i + 1 \leq m) \land \text{init}\} \\
\{(r = n^{i+1}) \land (i + 1 \leq m) \land \text{init} \land (v = N)\} \\
\{\text{Inv} \land (v < N)\} \\
\{\text{Inv} \land (i \geq m)\} \\
\{r = n^m\}
\end{align*}
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What is a suitable variant? $v := (m - i)$
Additional proof obligations

init: \((m \geq 0) \land (n > 0)\)

Inv: \( (r = n^i) \land (i \leq m) \land \text{init} \)

\(v: m - i\)

- \(\text{Inv} \land (i < m) \rightarrow (v > 0)\)
- \([v = N] i := i + 1 [v < N]\)
Additional proof obligations

Total correctness Hoare logic is designed to prove partial correctness and termination at the same time.

You can also do them separately:

1. Prove a partial correctness Hoare triple.
2. Find a variant for every loop.

Doing it completely separate isn’t always possible: sometimes, termination depends on the invariant.
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- Adding non-determinism
Non-determinism involves the computational model branching into one of several directions.

Any branch can happen (decision is not under our control).
Non-determinism

Why add non-determinism?

- More general than deterministic behaviour
- Sometimes useful for modelling interaction (c.f. coffee machines).
- Useful for abstraction (abstracted code is easier to reason about)
$\mathcal{L}^+$: a simple language with non-determinism

We relax the Conditional and Loop commands in $\mathcal{L}$ to give us non-deterministic behaviour.

The programs of $\mathcal{L}^+$ are defined as:

- **Assign:** $x := e$, where $x$ is a variable and $e$ is an expression
- **Predicate:** $\varphi$, where $\varphi$ is a predicate
- **Sequence:** $P; Q$, where $P$ and $Q$ are programs
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$P :: (x := e) \mid \varphi \mid P_1; P_2 \mid P_1 + P_2 \mid P_1^*$
\( \mathcal{L}^+ \): a simple language with non-determinism

\[
P :: (x := e) \mid \varphi \mid P_1; P_2 \mid P_1 + P_2 \mid P_1^*
\]

**NB**

\( \mathcal{L} \) can be defined in \( \mathcal{L}^+ \) by defining:

- if \( b \) then \( P \) else \( Q \) od = \((b; P) + (\neg b; Q)\)
- while \( b \) do \( P \) od = \((b; P)^*; \neg b\)
A program in $\mathcal{L}^+$ that non-deterministically checks if $(x \lor y) \land (\neg x \lor \neg z) \land (\neg y \lor z)$ is satisfiable:

\[
\begin{array}{|l|}
\hline
\text{SAT} \\
(x := 0) + (x := 1); \\
(y := 0) + (y := 1); \\
(z := 0) + (z := 1); \\
\hline
\end{array}
\]
A program in $\mathcal{L}^+$ that non-deterministically checks if $(x \lor y) \land (\neg x \lor \neg z) \land (\neg y \lor z)$ is satisfiable:

\[
\begin{array}{|c|}
\hline
\text{SAT} \\
(x := 0) + (x := 1); \\
(y := 0) + (y := 1); \\
(z := 0) + (z := 1); \\
\text{if} ((x = 1) \lor (y = 1)) \land \\
\quad ((x = 0) \lor (z = 0)) \land \\
\quad ((y = 0) \lor (z = 1)) \\
\hline
\end{array}
\]
A program in $\mathcal{L}^+$ that non-deterministically checks if $(x \lor y) \land (\neg x \lor \neg z) \land (\neg y \lor z)$ is satisfiable:

\[
\text{SAT}
\]

\[
\begin{align*}
  (x &:= 0) + (x := 1); \\
  (y &:= 0) + (y := 1); \\
  (z &:= 0) + (z := 1); \\
  \text{if}((x = 1) \lor (y = 1)) \land \\
    ((x = 0) \lor (z = 0)) \land \\
    ((y = 0) \lor (z = 1))
  \text{then } r := 1 \\
  \text{else } r := 0 \\
  \text{fi}
\end{align*}
\]

The formula is satisfiable if SAT could set $r$ to 1.
Proof rules

Hoare logic rules are cleaner:

\[
\begin{align*}
\{ \varphi \} P \{ \psi \} & \quad \{ \varphi \} Q \{ \psi \} \\
\{ \varphi \} P + Q \{ \psi \} \\
\{ \varphi \} P \{ \varphi \} & \quad \{ \varphi \} P^* \{ \varphi \}
\end{align*}
\]

(choice) (loop)
Semantics

Semantics is as for $\mathcal{L}$, except:

\[
[P + Q] = [P] \cup [Q] \\
[P^*] = [P]^*
\]
What follows is a proof that Hoare logic is sound.

We most likely won’t have time to do any of this in the lectures.
Summary

- Set theory revisited
- Soundness of Hoare Logic
- Completeness of Hoare Logic
Summary

- Set theory revisited
- Soundness of Hoare Logic
- Completeness of Hoare Logic
Some results on relational images

**Lemma**

For any binary relations \( R, S \subseteq X \times Y \) and subsets \( A, B \subseteq X \):

(a) If \( A \subseteq B \) then \( R(A) \subseteq R(B) \)

(b) \( R(A) \cup S(A) = (R \cup S)(A) \)

(c) \( R(S(A)) = (S; R)(A) \)
Some results on relational images

**Lemma**

For any binary relations $R, S \subseteq X \times Y$ and subsets $A, B \subseteq X$:

(a) If $A \subseteq B$ then $R(A) \subseteq R(B)$

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(c) $R(S(A)) = (S; R)(A)$

Proof (a):
Lemma

For any binary relations $R, S \subseteq X \times Y$ and subsets $A, B \subseteq X$:

(a) If $A \subseteq B$ then $R(A) \subseteq R(B)$

(b) $R(A) \cup S(A) = (R \cup S)(A)$

(c) $R(S(A)) = (S; R)(A)$

Proof (a):

$y \in R(A) \iff \exists x \in A$ such that $(x, y) \in R$

$\Rightarrow \exists x \in B$ such that $(x, y) \in R$

$\iff y \in R(B)$
Some results on relational images

Lemma

For any binary relations $R, S \subseteq X \times Y$ and subsets $A, B \subseteq X$:

(a) If $A \subseteq B$ then $R(A) \subseteq R(B)$

(b) $R(A) \cup S(A) = (R \cup S)(A)$

(c) $R(S(A)) = (S; R)(A)$

Proof (b):
Some results on relational images

**Lemma**

For any binary relations $R, S \subseteq X \times Y$ and subsets $A, B \subseteq X$:

(a) If $A \subseteq B$ then $R(A) \subseteq R(B)$

(b) $R(A) \cup S(A) = (R \cup S)(A)$

(c) $R(S(A)) = (S; R)(A)$

Proof (b):

$y \in R(A) \cup S(A) \iff y \in R(A)$ or $y \in S(A)$

$\iff \exists x \in A \text{ s.t. } (x, y) \in R$ or $\exists x \in A \text{ s.t. } (x, y) \in S$

$\iff \exists x \in A \text{ s.t. } (x, y) \in R$ or $(x, y) \in S$

$\iff \exists x \in A \text{ s.t. } (x, y) \in (R \cup S)$

$\iff y \in (R \cup S)(A)$
Some results on relational images

Lemma

For any binary relations $R, S \subseteq X \times Y$ and subsets $A, B \subseteq X$:

(a) If $A \subseteq B$ then $R(A) \subseteq R(B)$

(b) $R(A) \cup S(A) = (R \cup S)(A)$

(c) $R(S(A)) = (S; R)(A)$

Proof (c):
Some results on relational images

Lemma

For any binary relations \( R, S \subseteq X \times Y \) and subsets \( A, B \subseteq X \):

(a) If \( A \subseteq B \) then \( R(A) \subseteq R(B) \)
(b) \( R(A) \cup S(A) = (R \cup S)(A) \)
(c) \( R(S(A)) = (S; R)(A) \)

Proof (c):

\[ z \in R(S(A)) \iff \exists y \in S(A) \text{ s.t. } (y, z) \in R \]
\[ \iff \exists x \in A, y \in S(A) \text{ s.t. } (x, y) \in S \text{ and } (y, z) \in R \]
\[ \iff \exists x \in A \text{ s.t. } (x, z) \in (S; R) \]
\[ \iff z \in (S; R)(A) \]
Corollary

If $R(A) \subseteq A$ then $R^*(A) \subseteq A$

Reformulated: assuming $R(A) \subseteq A$, $x \in A$, and $x R^* y$, prove $y \in A$.

Proof is by induction on the derivation of $x R^* y$.

(B) Trivial when $x = y$.

(I) We know that $x \in A$, $x R y$ and $y R^* z$. Because $R(A) \subseteq A$, we have $y \in A$. By the induction hypothesis, $z \in A$. 
Summary

- Set theory revisited
- Soundness of Hoare Logic
- Completeness of Hoare Logic
Soundness of Hoare Logic

**Theorem**

If \( \models \{ \varphi \} P \{ \psi \} \) then \( \vdash \{ \varphi \} P \{ \psi \} \)

Proof: By induction on the structure of the proof.
Soundness of Hoare Logic

Theorem

If $\vdash \{ \varphi \} P \{ \psi \}$ then $\models \{ \varphi \} P \{ \psi \}$

Proof:
Soundness of Hoare Logic

Theorem

\[ \text{If } \vdash \{ \varphi \} \ P \ \{ \psi \} \ \text{then } \models \{ \varphi \} \ P \ \{ \psi \} \]

Proof:
By induction on the structure of the proof.
**Base case: Assignment rule**

\[
\{\varphi[e/x]\} \ x := e \{\varphi\} \quad \text{(ass)}
\]
Base case: Assignment rule

\[ \{\varphi[e/x]\} x := e \{\varphi\} \] (ass)

Need to show \( \{\varphi[e/x]\} x := e \{\varphi\} \) is always valid. That is,

\[ \llbracket x := e \rrbracket(\llangle \varphi[e/x] \rrangle) \subseteq \llangle \varphi \rrangle. \]
Base case: Assignment rule

\[
\{\varphi[e/x]\} \ x := e \ \{\varphi\} \quad \text{(ass)}
\]

Need to show \(\{\varphi[e/x]\} \ x := e \ \{\varphi\}\) is always valid. That is,

\[
\llbracket x := e \rrbracket (\llangle \varphi[e/x] \rrangle) \subseteq \llangle \varphi \rrangle.
\]

Observation: \(\llbracket \varphi[e/x] \rrbracket^\eta = \llbracket \varphi \rrbracket^{\eta'}\) where \(\eta' = \eta[x \mapsto \llbracket e \rrbracket^\eta]\)
Base case: Assignment rule

\[
\{\varphi[e/x]\} x := e \{\varphi\} \quad \text{(ass)}
\]

Need to show \(\{\varphi[e/x]\} x := e \{\varphi\}\) is always valid. That is,

\[
[x := e](\langle \varphi[e/x] \rangle) \subseteq \langle \varphi \rangle.
\]

Observation: \(\left[\varphi[e/x]\right]^{\eta} = \left[\varphi\right]^{\eta'}\) where \(\eta' = \eta[x \mapsto [e]^{\eta}]\)

So if \(\eta \in \langle \varphi[e/x] \rangle\) then \(\eta' \in \langle \varphi \rangle\)
**Base case: Assignment rule**

\[
\{\varphi[e/x]\} x := e \{\varphi\} \quad \text{(ass)}
\]

Need to show \(\{\varphi[e/x]\} x := e \{\varphi\}\) is always valid. That is,

\[
[x := e] (\langle \varphi[e/x] \rangle) \subseteq \langle \varphi \rangle.
\]

Observation: \([\varphi[e/x]]^\eta = [\varphi]^\eta'\) where \(\eta' = \eta[x \mapsto [e]^\eta]\)

So if \(\eta \in \langle \varphi[e/x] \rangle\) then \(\eta' \in \langle \varphi \rangle\)

Recall: \((\eta, \eta'') \in [x := e]\) if and only if \(\eta'' = \eta[x \mapsto [e]^\eta]\),


**Base case: Assignment rule**

\[
\{\varphi[e/x]\} x := e \{\varphi\} \quad \text{(ass)}
\]

Need to show \(\{\varphi[e/x]\} x := e \{\varphi\}\) is always valid. That is,

\[
[x := e](\langle \varphi[e/x] \rangle) \subseteq \langle \varphi \rangle.
\]

Observation: \(\{\varphi[e/x]\}^\eta = \{\varphi\}^{\eta'}\) where \(\eta' = \eta[x \mapsto [e]^\eta]\)

So if \(\eta \in \langle \varphi[e/x] \rangle\) then \(\eta' \in \langle \varphi \rangle\)

Recall: \((\eta, \eta'') \in [x := e] \) if and only if \(\eta'' = \eta[x \mapsto [e]^\eta]\),

So \([x := e](\eta) \in \langle \varphi \rangle\) for all \(\eta \in \langle \varphi[e/x] \rangle\)
Base case: Assignment rule

\[
\{\varphi[e/x]\} x := e \{\varphi\} \quad \text{(ass)}
\]

Need to show \{\varphi[e/x]\} x := e \{\varphi\} is always valid. That is,

\[
\mathbb{[}[x := e]\mathbb{]}(\langle \varphi[e/x]\rangle) \subseteq \langle \varphi \rangle.
\]

Observation: \mathbb{[}[\varphi[e/x]]]^{\eta} = \mathbb{[}[\varphi]]^{\eta'} \text{ where } \eta' = \eta[x \mapsto \mathbb{[}[e]]^{\eta}]

So if \( \eta \in \langle \varphi[e/x]\rangle \) then \( \eta' \in \langle \varphi \rangle \)

Recall: \((\eta, \eta'') \in \mathbb{[}[x := e]\mathbb{]}\) if and only if \( \eta'' = \eta[x \mapsto \mathbb{[}[e]]^{\eta}] \),

So \mathbb{[}[x := e]\mathbb{]}(\eta) \in \langle \varphi \rangle \text{ for all } \eta \in \langle \varphi[e/x]\rangle

So \mathbb{[}[x := e]\mathbb{]}(\langle \varphi[e/x]\rangle) \subseteq \langle \varphi \rangle
Inductive case 1: Sequence rule

\[
\begin{array}{c}
\{\varphi\} \ P \ \{\psi\} \\
\{\psi\} \ Q \ \{\rho\}
\end{array}
\]

\[
\frac{}
\{\varphi\} \ P; \ Q \ \{\rho\}
\]

(seq)
Inductive case 1: Sequence rule

\[
\begin{array}{c}
\{\varphi\} P \{\psi\} \\
\{\psi\} Q \{\rho\} \\
\hline
\{\varphi\} P; Q \{\rho\}
\end{array}
\] (seq)

Assume \{\varphi\} P \{\psi\} and \{\psi\} Q \{\rho\} are valid. Need to show that \{\varphi\} P; Q \{\rho\} is valid.
Inductive case 1: Sequence rule

\[
\begin{array}{c}
\{\varphi\} P \{\psi\} \quad \{\psi\} Q \{\rho\} \\
\hline
\{\varphi\} P; Q \{\rho\}
\end{array}
\]

(seq)

Assume \{\varphi\} P \{\psi\} and \{\psi\} Q \{\rho\} are valid. Need to show that \{\varphi\} P; Q \{\rho\} is valid.

Recall: \([P; Q] = [P]; [Q]\)
Inductive case 1: Sequence rule

\[
\begin{array}{c}
\{ \varphi \} \quad P \quad \{ \psi \} \\
\{ \psi \} \quad Q \quad \{ \rho \}
\end{array}
\]

\[
\{ \varphi \} \quad P; \quad Q \quad \{ \rho \}
\]

(seq)

Assume \(\{ \varphi \} \quad P \quad \{ \psi \}\) and \(\{ \psi \} \quad Q \quad \{ \rho \}\) are valid. Need to show that \(\{ \varphi \} \quad P; \quad Q \quad \{ \rho \}\) is valid.

Recall: \([P; Q] = [P]; [Q]\)

So: \([P; Q](\langle \varphi \rangle) = [Q](P(\langle \varphi \rangle))\)  

(see Lemma 1(c))
Inductive case 1: Sequence rule

\[
\begin{array}{c}
\{ \varphi \} \quad P \quad \{ \psi \} \\
\{ \psi \} \quad Q \quad \{ \rho \}
\end{array} \\
\{ \varphi \} \quad P; Q \quad \{ \rho \}
\]  
(seq)

Assume \( \{ \varphi \} \quad P \quad \{ \psi \} \) and \( \{ \psi \} \quad Q \quad \{ \rho \} \) are valid. Need to show that \( \{ \varphi \} \quad P; Q \quad \{ \rho \} \) is valid.

Recall: \([ P; Q ] = [ P ]; [ Q ]\)

So: \([ P; Q ](\langle \varphi \rangle) = [ Q ]([ P ](\langle \varphi \rangle))\)  
(see Lemma 1(c))

By IH: \([ P ](\langle \varphi \rangle) \subseteq \langle \psi \rangle\) and \([ Q ](\langle \psi \rangle) \subseteq \langle \rho \rangle\)
Inductive case 1: Sequence rule

\[
\begin{array}{c}
\{\varphi\} P \{\psi\} \\
\{\psi\} Q \{\rho\}
\end{array}
\]

(\text{seq})

\[
\{\varphi\} P; Q \{\rho\}
\]

Assume \(\{\varphi\} P \{\psi\}\) and \(\{\psi\} Q \{\rho\}\) are valid. Need to show that \(\{\varphi\} P; Q \{\rho\}\) is valid.

Recall: \([P; Q] = [P]; [Q]\)

So: \([P; Q](\langle \varphi \rangle) = [Q]( [P](\langle \varphi \rangle))\) \hspace{1cm} (see Lemma 1(c))

By IH: \([P](\langle \varphi \rangle) \subseteq \langle \psi \rangle\) and \([Q](\langle \psi \rangle) \subseteq \langle \rho \rangle\)

So: \([Q]( [P](\langle \varphi \rangle)) \subseteq [Q](\langle \psi \rangle) \subseteq \langle \rho \rangle\) \hspace{1cm} (see Lemma 1(a))
Two more useful results

Lemma

For \( R \subseteq \text{Env} \times \text{Env} \), predicates \( \varphi \) and \( \psi \), and \( X \subseteq \text{Env} \):

(a) \( \llbracket \varphi \rrbracket (X) = \langle \varphi \rangle \cap X \)

(b) \( R(\langle \varphi \land \psi \rangle) = (\llbracket \varphi \rrbracket; R)(\langle \psi \rangle) \)
Two more useful results

Lemma

For $R \subseteq \text{Env} \times \text{Env}$, predicates $\varphi$ and $\psi$, and $X \subseteq \text{Env}$:

(a) $[\varphi](X) = \langle \varphi \rangle \cap X$

(b) $R(\langle \varphi \land \psi \rangle) = ([\varphi]; R)(\langle \psi \rangle)$

Proof (a):
Two more useful results

Lemma

For $R \subseteq \text{Env} \times \text{Env}$, predicates $\varphi$ and $\psi$, and $X \subseteq \text{Env}$:

(a) $\llbracket \varphi \rrbracket(X) = \langle \varphi \rangle \cap X$

(b) $R(\langle \varphi \land \psi \rangle) = (\llbracket \varphi \rrbracket; R)(\langle \psi \rangle))$

Proof (a):

$\eta' \in \llbracket \varphi \rrbracket(X) \Leftrightarrow \exists \eta \in X \text{ s.t. } (\eta, \eta') \in \llbracket \varphi \rrbracket$

$\Leftrightarrow \exists \eta \in X \text{ s.t. } \eta = \eta' \text{ and } \eta \in \langle \varphi \rangle$

$\Leftrightarrow \eta' \in X \cap \langle \varphi \rangle$
Two more useful results

**Lemma**

For \( R \subseteq \text{Env} \times \text{Env} \), predicates \( \varphi \) and \( \psi \), and \( X \subseteq \text{Env} \):

(a) \( \llbracket \varphi \rrbracket (X) = \langle \varphi \rangle \cap X \)

(b) \( R(\langle \varphi \land \psi \rangle) = (\llbracket \varphi \rrbracket ; R)(\langle \psi \rangle) \)

Proof (b):

\[
\langle \varphi \land \psi \rangle = \langle \varphi \rangle \cap \langle \psi \rangle = \llbracket \varphi \rrbracket (\langle \psi \rangle)
\]

So \( R(\langle \varphi \land \psi \rangle) = R(\llbracket \varphi \rrbracket (\langle \psi \rangle)) = (\llbracket \varphi \rrbracket ; R)(\langle \psi \rangle) \) (see Lemma 1(b))
Inductive case 2: Conditional rule

\[
\frac{\{ \varphi \land g \} P \{ \psi \} \quad \{ \varphi \land \neg g \} Q \{ \psi \}}{\{ \varphi \} \text{ if } g \text{ then } P \text{ else } Q \text{ fi } \{ \psi \}} \quad \text{(if)}
\]
Inductive case 2: Conditional rule

\[
\begin{array}{c}
\{ \varphi \land g \} \text{ } P \{ \psi \} \quad \{ \varphi \land \neg g \} \text{ } Q \{ \psi \} \\
\{ \varphi \} \text{ if } g \text{ then } P \text{ else } Q \text{ fi } \{ \psi \}
\end{array}
\]

(if)

Assume \( \{ \varphi \land g \} \text{ } P \{ \psi \} \) and \( \{ \varphi \land \neg g \} \text{ } Q \{ \psi \} \) are valid. Need to show that \( \{ \varphi \} \text{ if } g \text{ then } P \text{ else } Q \text{ fi } \{ \psi \} \) is valid.
Inductive case 2: Conditional rule

\[
\begin{array}{c}
\{ \varphi \land g \} \quad P \{ \psi \} \\
\{ \varphi \land \neg g \} \quad Q \{ \psi \}
\end{array}
\]

(if)

\{ \varphi \} \text{ if } g \text{ then } P \text{ else } Q \text{ fi } \{ \psi \}

Assume \{ \varphi \land g \} \ P \{ \psi \} \text{ and } \{ \varphi \land \neg g \} \ Q \{ \psi \} \text{ are valid. Need to show that } \{ \varphi \} \text{ if } g \text{ then } P \text{ else } Q \text{ fi } \{ \psi \} \text{ is valid.}

Recall: \[[\text{if } g \text{ then } P \text{ else } Q \text{ fi}] = [g; P] \cup [\neg g; Q]\]
Inductive case 2: Conditional rule

\[
\begin{align*}
\{ \varphi \land g \} & \quad P \{ \psi \} \quad \{ \varphi \land \neg g \} \quad Q \{ \psi \} \\
\{ \varphi \} & \quad \text{if } g \text{ then } P \text{ else } Q \text{ fi } \{ \psi \}
\end{align*}
\]

(if)

Assume \( \{ \varphi \land g \} \quad P \{ \psi \} \) and \( \{ \varphi \land \neg g \} \quad Q \{ \psi \} \) are valid. Need to show that \( \{ \varphi \} \quad \text{if } g \text{ then } P \text{ else } Q \text{ fi } \{ \psi \} \) is valid.

Recall: \( [[ \text{if } g \text{ then } P \text{ else } Q \text{ fi} ]] = [g; P] \cup [\neg g; Q] \)

\( [[ \text{if } g \text{ then } P \text{ else } Q \text{ fi} ]](\langle \varphi \rangle) \)
Inductive case 2: Conditional rule

\[
\begin{array}{c}
\{\varphi \land g\} P \{\psi\} \quad \{\varphi \land \neg g\} Q \{\psi\} \\
\{\varphi\} \text{ if } g \text{ then } P \text{ else } Q \text{ fi } \{\psi\}
\end{array}
\] (if)

Assume \( \{\varphi \land g\} P \{\psi\} \) and \( \{\varphi \land \neg g\} Q \{\psi\} \) are valid. Need to show that \( \{\varphi\} \text{ if } g \text{ then } P \text{ else } Q \text{ fi } \{\psi\} \) is valid.

Recall: \([\text{if } g \text{ then } P \text{ else } Q \text{ fi}] = [g; P] \cup [\neg g; Q]

\[
[\text{if } g \text{ then } P \text{ else } Q \text{ fi}](\langle \varphi \rangle) \\
= [g; P](\langle \varphi \rangle) \cup [\neg g; Q](\langle \varphi \rangle)
\] (see Lemma 1(b))
Inductive case 2: Conditional rule

\[
\begin{array}{ccc}
\{ \varphi \land g \} & P & \{ \psi \} \\
\{ \varphi \} & \text{if } g \text{ then } P \text{ else } Q \text{ fi} & \{ \psi \}
\end{array}
\]

(if)

Assume \( \{ \varphi \land g \} \ P \ \{ \psi \} \) and \( \{ \varphi \land \neg g \} \ Q \ \{ \psi \} \) are valid. Need to show that \( \{ \varphi \} \text{ if } g \text{ then } P \text{ else } Q \text{ fi} \ \{ \psi \} \) is valid.

Recall: \([\text{if } g \text{ then } P \text{ else } Q \text{ fi}] = [g; P] \cup [\neg g; Q] \)

\([\text{if } g \text{ then } P \text{ else } Q \text{ fi}](\langle \varphi \rangle) \)

\[= [g; P](\langle \varphi \rangle) \cup [\neg g; Q](\langle \varphi \rangle) \] (see Lemma 1(b))

\[= [P](\langle g \land \varphi \rangle) \cup [Q](\langle \neg g \land \varphi \rangle) \] (see Lemma 2(b))
Inductive case 2: Conditional rule

\[
\begin{array}{c}
\{\varphi \land g\} P \{\psi\} \quad \{\varphi \land \neg g\} Q \{\psi\} \\
\{\varphi\} \text{ if } g \text{ then } P \text{ else } Q \text{ fi } \{\psi\}
\end{array}
\]

(if)

Assume \(\{\varphi \land g\} P \{\psi\}\) and \(\{\varphi \land \neg g\} Q \{\psi\}\) are valid. Need to show that \(\{\varphi\} \text{ if } g \text{ then } P \text{ else } Q \text{ fi } \{\psi\}\) is valid.

Recall: \[[\text{if } g \text{ then } P \text{ else } Q \text{ fi}]] = [[g; P]] \cup [[\neg g; Q]]

\[
[[\text{if } g \text{ then } P \text{ else } Q \text{ fi}]](\langle \varphi \rangle)
\]

\[
= [[g; P]](\langle \varphi \rangle) \cup [[\neg g; Q]](\langle \varphi \rangle) \quad \text{(see Lemma 1(b))}
\]
\[
= [[P]](\langle g \land \varphi \rangle) \cup [[Q]](\langle \neg g \land \varphi \rangle) \quad \text{(see Lemma 2(b))}
\]
\[
\subseteq \langle \psi \rangle \quad \text{(by IH)}
\]
Inductive case 3: While rule

\[
\begin{align*}
\{\varphi \land g\} &\implies P \{\varphi\} \\
\{\varphi\} &\implies \text{while } g \text{ do } P \text{ od } \{\varphi \land \neg g\} \\
\text{(loop)}
\end{align*}
\]
Inductive case 3: While rule

\[
\begin{align*}
\{\varphi \land g\} & \quad P \quad \{\varphi\} \\
\{\varphi\} & \quad \text{while} \quad g \quad \text{do} \quad P \quad \text{od} \quad \{\varphi \land \neg g\}
\end{align*}
\]  \quad \text{(loop)}

Assume \(\{\varphi \land g\} \quad P \quad \{\varphi\}\) is valid. Need to show that
\(\{\varphi\} \quad \text{while} \quad g \quad \text{do} \quad P \quad \text{od} \quad \{\varphi \land \neg g\}\) is valid.
Inductive case 3: While rule

\[
\begin{align*}
\{ \varphi \land g \} & P \{ \varphi \} \\
\{ \varphi \} & \text{while } g \text{ do } P \text{ od } \{ \varphi \land \neg g \}
\end{align*}
\]

(loop)

Assume \( \{ \varphi \land g \} P \{ \varphi \} \) is valid. Need to show that \( \{ \varphi \} \text{ while } g \text{ do } P \text{ od } \{ \varphi \land \neg g \} \) is valid.

Recall: \([\text{while } g \text{ do } P \text{ od}] = [g; P]^*; [\neg g]\)
Inductive case 3: While rule

\[
\{\varphi \land g\} \quad P \quad \{\varphi\} \\
\{\varphi\} \quad \text{while } g \quad \text{do } P \quad \text{od} \quad \{\varphi \land \neg g\}
\]

(loop)

Assume \(\{\varphi \land g\} \quad P \quad \{\varphi\}\) is valid. Need to show that \(\{\varphi\} \quad \text{while } g \quad \text{do } P \quad \text{od} \quad \{\varphi \land \neg g\}\) is valid.

Recall: 
\[
[\text{while } g \quad \text{do } P \quad \text{od}] = [g \mid P]^* \cdot [\neg g]
\]

\[
[g \mid P](\langle \varphi \rangle) = [P](\langle g \land \varphi \rangle)
\]

(see Lemma 2(b))
Inductive case 3: While rule

\[
\begin{array}{c}
\{ \varphi \land g \} \quad P \quad \{ \varphi \} \\
\hline
\{ \varphi \} \quad \text{while } g \text{ do } P \text{ od } \{ \varphi \land \neg g \}
\end{array}
\] (loop)

Assume \( \{ \varphi \land g \} \quad P \quad \{ \varphi \} \) is valid. Need to show that \( \{ \varphi \} \quad \text{while } g \text{ do } P \text{ od } \{ \varphi \land \neg g \} \) is valid.

Recall: \([\text{while } g \text{ do } P \text{ od}] = [g; P]^*; [\neg g] \)

\[
[g; P](\langle \varphi \rangle) = [P](\langle g \land \varphi \rangle)
\] (see Lemma 2(b))

\[
\subseteq \langle \varphi \rangle
\] (IH)
Inductive case 3: While rule

\[
\begin{align*}
\{\varphi \land g\} & \quad P \quad \{\varphi\} \\
\{\varphi\} \text{ while } g \text{ do } P \text{ od } \{\varphi \land \neg g\}
\end{align*}
\] (loop)

Assume \(\{\varphi \land g\} \quad P \quad \{\varphi\}\) is valid. Need to show that \(\{\varphi\}\) while \(g\) do \(P\) od \(\{\varphi \land \neg g\}\) is valid.

Recall: \([\text{while } g \text{ do } P \text{ od}] = [g; P]^*; [\neg g]\)

\([g; P](\langle \varphi \rangle) = [P](\langle g \land \varphi \rangle)\) \hspace{1cm} (see Lemma 2(b))

\(\subseteq \langle \varphi \rangle\) \hspace{1cm} (IH)

So \([g; P]^*(\langle \varphi \rangle) \subseteq \langle \varphi \rangle\) \hspace{1cm} (see Corollary)
Inductive case 3: While rule

\[
\begin{align*}
\{\varphi \land g\} & \quad P \{\varphi\} \\
\{\varphi\} \text{ while } g \text{ do } P \text{ od } \{\varphi \land \neg g\} 
\end{align*}
\]

(loop)

Assume \( \{\varphi \land g\} \ P \{\varphi\} \) is valid. Need to show that \( \{\varphi\} \text{ while } g \text{ do } P \text{ od } \{\varphi \land \neg g\} \) is valid.

Recall: \([\text{while } g \text{ do } P \text{ od}] = [g; P]^*; [\neg g]\)

\[
[g; P](\langle \varphi \rangle) = [P](\langle g \land \varphi \rangle) \subseteq \langle \varphi \rangle 
\]

(see Lemma 2(b)) (IH)

So \([g; P]^*(\langle \varphi \rangle) \subseteq \langle \varphi \rangle \)

(see Corollary)

So \( [g; P]^*; [\neg g](\langle \varphi \rangle) = [\neg g]([g; P]^*(\langle \varphi \rangle)) \) (see Lemma 1(c))
Inductive case 3: While rule

\[
\{\varphi \land g\} \xrightarrow{P} \{\varphi\} \\
\{\varphi\} \text{ while } g \text{ do } P \text{ od } \{\varphi \land \neg g\}
\]

(loop)

Assume \( \{\varphi \land g\} P \{\varphi\} \) is valid. Need to show that \( \{\varphi\} \text{ while } g \text{ do } P \text{ od } \{\varphi \land \neg g\} \) is valid.

Recall: \([\text{while } g \text{ do } P \text{ od}] = [g; P]^*; [\neg g]\)

\[
[g; P](\langle \varphi \rangle) = [P](\langle g \land \varphi \rangle)
\]

\(\subseteq \langle \varphi \rangle\)  

(IH)  

(see Lemma 2(b))

So \([g; P]^* (\langle \varphi \rangle) \subseteq \langle \varphi \rangle\)  

(see Corollary)

So \([g; P]^*; [\neg g](\langle \varphi \rangle) = [\neg g]( [g; P]^* (\langle \varphi \rangle))\)  

(see Lemma 1(c))

\(\subseteq [\neg g](\langle \varphi \rangle)\)  

(see Lemma 1(a))
Inductive case 3: While rule

\[
\begin{array}{c}
\{ \varphi \land g \} \quad P \quad \{ \varphi \} \\
\{ \varphi \} \quad \text{while} \quad g \quad \text{do} \quad P \quad \text{od} \quad \{ \varphi \land \neg g \}
\end{array}
\]

(loop)

Assume \( \{ \varphi \land g \} \quad P \quad \{ \varphi \} \) is valid. Need to show that \( \{ \varphi \} \quad \text{while} \quad g \quad \text{do} \quad P \quad \text{od} \quad \{ \varphi \land \neg g \} \) is valid.

Recall: \([\text{while} \quad g \quad \text{do} \quad P \quad \text{od}] = [g; P]^*; [\neg g]\)

\[
[g; P](\langle \varphi \rangle) = [P](\langle g \land \varphi \rangle)
\]

\(\subseteq \langle \varphi \rangle\)  

(IH)  

(see Lemma 2(b))

So \([g; P]^*(\langle \varphi \rangle) \subseteq \langle \varphi \rangle\)  

(see Corollary)

So \([g; P]^*; [\neg g](\langle \varphi \rangle) = [\neg g]( [g; P]^*(\langle \varphi \rangle)) \]

\(\subseteq [\neg g](\langle \varphi \rangle)\)  

(see Lemma 1(c))

\(\subseteq \langle \neg g \land \varphi \rangle\)  

(see Lemma 1(a))

(see Lemma 2(a))
Inductive case 4: Consequence rule

\[
\varphi' \rightarrow \varphi \quad \{\varphi\} \quad P \quad \{\psi\} \quad \psi \rightarrow \psi' \\
\{\varphi'\} \quad P \quad \{\psi'\}
\]

(cons)
Inductive case 4: Consequence rule

\[ \varphi' \rightarrow \varphi \quad \{ \varphi \} \ P \ \{ \psi \} \quad \psi \rightarrow \psi' \]

\[ \{ \varphi' \} \ P \ \{ \psi' \} \] (cons)

Assume \( \{ \varphi \} \ P \ \{ \psi \} \) is valid and \( \varphi' \rightarrow \varphi \) and \( \psi \rightarrow \psi' \). Need to show that \( \{ \varphi' \} \ P \ \{ \psi' \} \) is valid.
Inductive case 4: Consequence rule

\[
\frac{\varphi' \rightarrow \varphi \quad \{\varphi\} P \{\psi\} \quad \psi \rightarrow \psi'}{\{\varphi'\} P \{\psi'\}} \quad \text{(cons)}
\]

Assume \(\{\varphi\} P \{\psi\}\) is valid and \(\varphi' \rightarrow \varphi\) and \(\psi \rightarrow \psi'\). Need to show that \(\{\varphi'\} P \{\psi'\}\) is valid.

Observe: If \(\varphi' \rightarrow \varphi\) then \(\langle \varphi' \rangle \subseteq \langle \varphi \rangle\)
Inductive case 4: Consequence rule

\[
\frac{\varphi' \rightarrow \varphi \quad \{\varphi\} P \{\psi\} \quad \psi \rightarrow \psi'}{\{\varphi'\} P \{\psi'\}} \quad \text{(cons)}
\]

Assume \(\{\varphi\} P \{\psi\}\) is valid and \(\varphi' \rightarrow \varphi\) and \(\psi \rightarrow \psi'\). Need to show that \(\{\varphi'\} P \{\psi'\}\) is valid.

Observe: If \(\varphi' \rightarrow \varphi\) then \(\langle \varphi' \rangle \subseteq \langle \varphi \rangle\)

\[
\llbracket P \rrbracket(\langle \varphi' \rangle) \subseteq \llbracket P \rrbracket(\langle \varphi \rangle) \quad \text{(see Lemma 1(a))}
\]
Inductive case 4: Consequence rule

\[
\frac{\varphi' \rightarrow \varphi \quad \{\varphi\} P \{\psi\} \quad \psi \rightarrow \psi'}{\{\varphi'\} P \{\psi'\}}\quad (\text{cons})
\]

Assume \(\{\varphi\} P \{\psi\}\) is valid and \(\varphi' \rightarrow \varphi\) and \(\psi \rightarrow \psi'\). Need to show that \(\{\varphi'\} P \{\psi'\}\) is valid.

Observe: If \(\varphi' \rightarrow \varphi\) then \(\langle\varphi'\rangle \subseteq \langle\varphi\rangle\)

\[
[P](\langle\varphi'\rangle) \subseteq [P](\langle\varphi\rangle) \quad (\text{see Lemma 1(a)})
\]

\[
\subseteq \langle\psi\rangle
\]

\[
\subseteq \langle\psi'\rangle \quad \text{(IH)}
\]