COMP 3221

## Microprocessors and Embedded Systems

Lectures 18 : Pointers \& Arrays in C/ Assembly
http://www.cse.unsw.edu.au/~cs3221
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## Review: Register Convention

- Caller Saved Registers:
- Return address
lr
- Arguments
a1, a2, a3, a4
- Return values
a1, a2, a3, a4
- Callee Saved Registers:
-v Registers v1 - v8


## Overview

${ }^{\circ}$ Arrays, Pointers, Functions in C
${ }^{\circ}$ Example
${ }^{\circ}$ Pointers, Arithmetic, and Dereference
${ }^{\circ}$ Conclusion

## Review: Function Call Bookkeeping

${ }^{\circ}$ Big Ideas:

- Follow the procedure conventions and nobody gets hurt.
- Data is just 1's and 0 's, what it represents depends on what you do with it
- Function Call Bookkeeping:
- Caller Saved Registers are saved by the caller, that is, the function that includes the bl instruction
- Callee Saved Registers are saved by the callee, that is, the function that includes the mov pc, lr instruction
- Some functions are both a caller and a callee


## Argument Passing Options

## ${ }^{\circ} 2$ choices

- "Call by Value": pass a copy of the item to the function/procedure
- $x$... $f(x)$... $x$. Call to $f$ does not change $x$
- "Call by Reference": pass a pointer to the item to the function/procedure
${ }^{\circ}$ Single word variables passed by value
${ }^{\circ}$ What about passing an array? e.g., a [100]
- Pascal--call by value--copies 100 words of a [] onto the stack
- C--call by reference--passes a pointer (1 word) to the array a [] in a register

Pointers Implementation in ARM
${ }^{\circ} \mathrm{c}$ is int, has value 100, in memory at address $0 \times 10000000, p$ in v1, $x$ in v2
$\mathrm{p}=\& \mathrm{c} ; ~ / * \mathrm{p}$ gets $0 \times 10000000$ */
x = *p; /* x gets 100 */
*p $=200 ; / *$ c gets 200 */

```
; p = &c; /* p gets 0x10000000 */
mov v1,0x1000000 ; p = 0x10000000
; x = *p; /* x gets 100 */
ldr v2, [v1] ; dereferencing p
; *p = 200; /* c gets 200 */
mov a1, #200
str a1, [v1] ; dereferencing p
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\section*{Arrays, Pointers, Functions in C}
\({ }^{\circ} 4\) versions of array function that adds two arrays and puts sum in a third array (sumarray)
- Third array is passed to function
- Using a local array (on stack) for result and passing a pointer to it
- Third array is allocated on heap
- Third array is declared static
\({ }^{\circ}\) Purpose of example is to show interaction of C statements, pointers, and memory allocation

\section*{Calling sumarray, Version 1}
```

int x[100], y[100], z[100];
sumarray(x, y, z);

```

\section*{\({ }^{\circ} \mathrm{C}\) calling convention means above the same as}
sumarray (\&x[0], \&y[0], \&z[0]);

\section*{\({ }^{\circ}\) Really passing pointers to arrays}
\begin{tabular}{llll} 
mov & \(\mathrm{a} 1, \mathrm{sb}\) & \(; \mathrm{x}[0]\) & starts at sb \\
add & \(\mathrm{a} 2, \mathrm{sb}, \# 400\) & \(; \mathrm{y}[0]\) & above \(\mathrm{x}[100]\) \\
add & \(\mathrm{a} 3, \mathrm{sb}, \# 800\) & \(; \mathrm{z}[0]\) above \(\mathrm{y}[100]\) \\
bl & sumarray & &
\end{tabular}

\section*{Version 2 to Fix Weakness of Version 1}
\({ }^{\circ}\) Would like recursion to work
int * sumarray(int a[],int b[]); /* adds 2 arrays and returns sum */
sumarray ( \(x, \operatorname{sumarray}(y, z)\) );

\section*{\({ }^{\circ}\) Cannot do this with Version 1 style solution: what about this}
```

int * sumarray(int a[],int b[]) {
int i, c[100];
for(i=0;i<100;i=i+1)
c[i] = a[i] + b[i];
return c;
}

```

\section*{Version 1: Optimized Compiled Code}
```

void sumarray(int a[],int b[],int c[]) {
int i;
for(i=0;i<100;i=i+1)
c[i] = a[i] + b[i];
}
sumarray: stmfd sp!,{v1- v2};save v1-v2 on stack
add a4, a1,\#400 ; beyond end of a[]
Loop: cmp a1, a4
beq Exit
ldr v1, [a1], \#4 ;a1=a[i], a1=a1+4
ldr v2, [a2], \#4 ;a2=b[i], a2=a2+4
add v2, v2, v1 ;v2=a[i] + b[i]
str v2, [a3], \#4 ;c[i]=a[i] + b[i]
; a3 = a3+4
b Loop
Exit: ldmfd sp!,{v1-v2}; restore v1-v2
mov pc, lr

```

Pointers, Arithmetic, and Dereference


\section*{Version 2: Revised Compiled Code}
```

for(i=0;i<100;i=i+1)
c[i] = a[i] + b[i];
return c;}

```
sumarray: stmfd sp!,\{v1-v2\};save v1-v2 on stack
        add \(a 4, a 1, \# 400 ;\) beyond end of \(a[]\)
        sub sp, sp,\#400 ; space for c
        mov a3, sp ; ptr for c
Loop:
    cmp al! a4
    beq Exit
    ldr v1, [a1], \#4 ;a1=a[i], a1=a1+4
    ldr v2, [a2], \#4 ;a2=b[i], a2=a2+4
    add v2, v2, v1 ; v2=a[i] + b[i]
    str v2, [a3], \#4 ;c[i]=a[i] + b[i]
    b Loop

Exit:
\begin{tabular}{ll} 
mov a1, sp & \(;\) \&c[0] \\
\hline add sp,sp, 400 & \(;\) pop stack \\
ldmfd sp!,\{v1-v2\}; restore v1-v2
\end{tabular}
        mov pc, lr

Weakness of Version 2
\({ }^{\circ}\) Legal Syntax; What's Wrong?
Address
\({ }^{\circ}\) Will work until call another function that uses stack
\({ }^{\circ} \mathrm{c}\) [100] Won't be reused instantly(e.g, add a printf)
\({ }^{\circ}\) Stack allocated + unrestricted pointer is problem


\section*{Version 3: Revised Compiled Code}
```

sumarray: stmfd sp!,{a1-a2,v1- v2,lr}
;save a1-a2, v1-v2 \& lr on stack
add a4, a1,\#400 ; beyond end of a[]
mov a1,\#400 bl malloc ; get space for c
mov a3, a1 ; get \&c
Idmfd sp!,{a1-a2} ; restor a1-a2
Loop:
cmp a1, a4
beq Exit
ldr v1, [a1], \#4 ;a1=a[i], a1=a1+4
ldr v2, [a2], \#4 ;a2=b[i], a2=a2+4
add v2, v2, v1 ;v2=a[i]' + b[i]
str v2, [a3], \#4 ;c[i]=a[i] + b[i]
; a3 = a3+4
b Loop
Exit:
sub a1, a3, \#400; ; \&c[0]
; and return

```

Version 3 to Fix Weakness of Version 2
\begin{tabular}{|c|c|}
\hline \({ }^{\circ}\) Solution: allocate c [] on heap & \\
\hline ```
int * sumarray(int a[],int b[]) {
    int i;
    int *c;
``` &  \\
\hline \(c=(\) int \(*\) ) malloc (100) ; & 4 \\
\hline \[
\begin{aligned}
& \text { for }(i=0 ; i<100 ; i=i+1) \\
& c[i]=a[i]+b[i] ; \\
& \text { return } c ;
\end{aligned}
\] & Heap
\[
c[100]
\] \\
\hline Not reused unless freed & Static \\
\hline \begin{tabular}{l}
- Can lead to memory leaks \\
- Java, has garbage collectors to reclaim free space
\end{tabular} & Code \\
\hline
\end{tabular}

\section*{Lifetime of storage \& scope}
\({ }^{\circ}\) automatic (stack allocated)
- typical local variables of a function
- created upon call, released upon return
- scope is the function
\({ }^{\circ}\) heap allocated
- created upon malloc, released upon free
- referenced via pointers
\({ }^{\circ}\) external / static
- exist for entire program
"What's This Stuff Good For?"


\section*{Version 4 : Alternative to Version 3}
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|l|}{\({ }^{\circ}\) Static declaration} \\
\hline & Stack \\
\hline ```
int * sumarray(int a[],int b[]) {
    int i;
    static int c[100];
``` &  \\
\hline for (i=0; i<100;i=i+1) & 1 \\
\hline \[
\underset{\mathrm{return}}{\mathrm{c} ; \mathrm{i}} \underset{\mathrm{c}}{ }=\mathrm{a}[\mathrm{i}]+\mathrm{b}[\mathrm{i}] ;
\] & Heap \\
\hline \multicolumn{2}{|l|}{\}} \\
\hline Compiler allocates once for function, space is reused & Static
\[
c[100]
\] \\
\hline - Will be changed next time sumarray invoked & Code \\
\hline - Used in C libraries & \\
\hline
\end{tabular}

\footnotetext{
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}

\section*{What about Structures?}
\({ }^{\circ}\) Scalars passed by value
\({ }^{\circ}\) Arrays passed by reference (pointers)
\({ }^{\circ}\) Structures by value too
\({ }^{\circ}\) Can think of C passing everything by value, just that arrays are simply a notation for pointers and the pointer is passed by value
"And in Conclusion.."
\({ }^{\circ}\) In C:
- Scalars passed by value
- Arrays passed by reference
\({ }^{\circ}\) In C functions we can return a pointer to Arrays defined in Static, Heap or stack area.
\({ }^{\circ}\) Returning a pointer to an array in stack gives rise to unrestricted pointers```

