

Assignment 3 Adv



Advance Assignment

- Shared pages and copy-on-write
- `sbrk()`
- Demand loading and `mmap`
- Paging

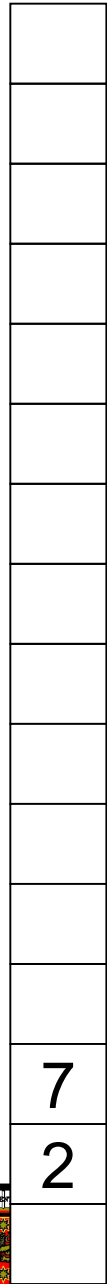


Shared pages and Copy-on-write

- What are they
- Why are they useful
- What they are not
 - Shared memory



Proc 1 Address Space



Two (or more) processes running the same program and sharing a section

7

2

Page Table



Proc 2 Address Space



Physical Address Space

B

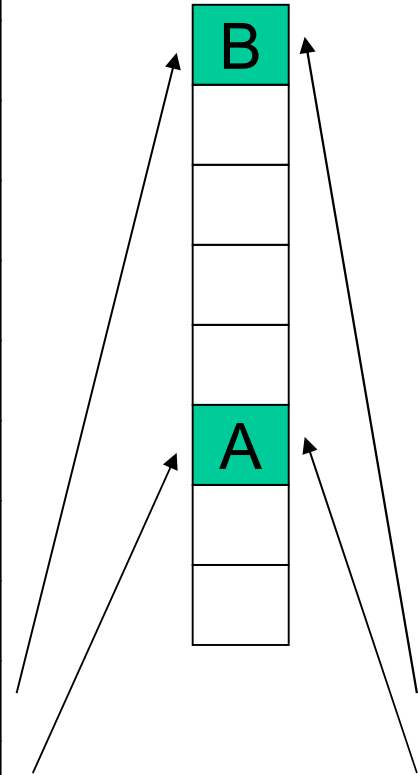
A

7

2

Page Table

4



COW

- fork() can be more efficient
- as_copy is underlying routine
- set pages read_only
 - Keep reference count in frame table
 - On write-fault, vm_fault copies, decrement count.



sbrk

- The "break" is the end address of a process's heap region.
- The sbrk call adjusts the "break" by the amount.
- It returns the old "break". Thus, to determine the current "break", call sbrk(0).
- The heap region is initially empty, so at process startup, the beginning of the heap region is the same as the end and may thus be retrieved using sbrk(0).

0xffffffff

0xc0000000

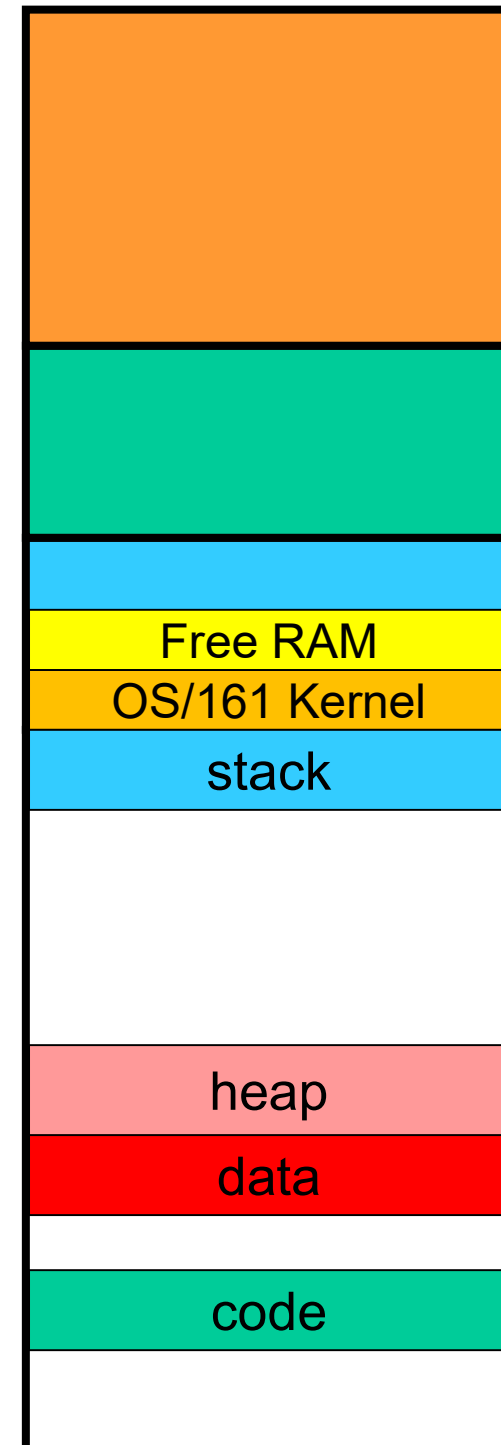
0xa0000000

0x80000000

0x10000000

0x04000000

0x00000000

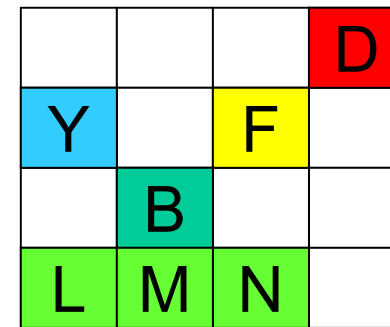
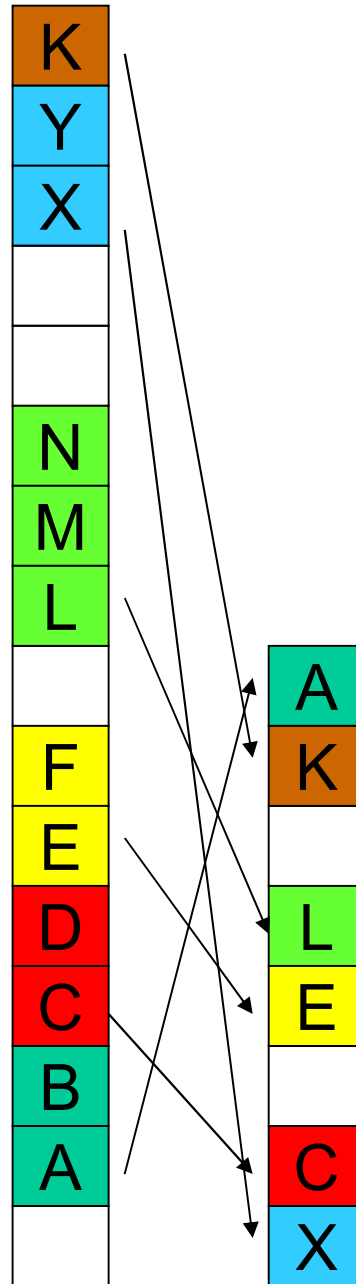


mmap() and demand loading



Memory-mapped files and paging

Memory mapped file



Disk

Physical Address Space



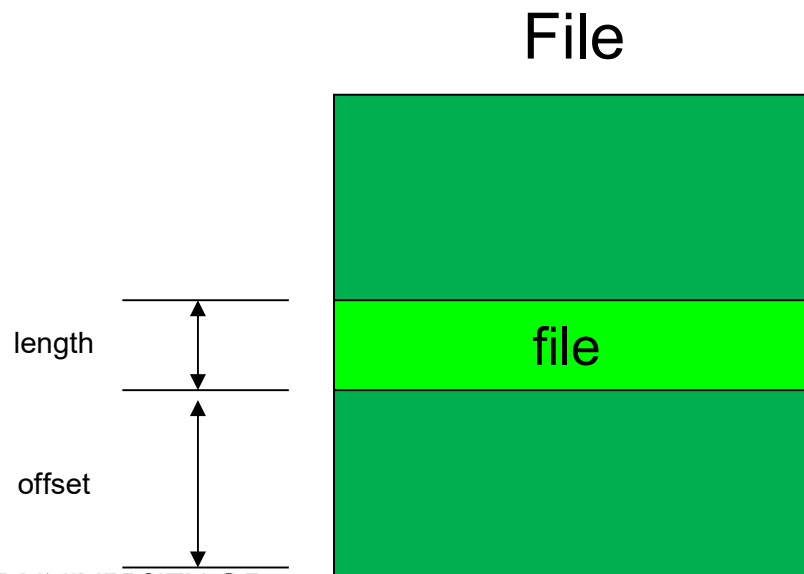
mmap/munmap semantics

```
void *mmap(size_t length, int prot, int fd, off_t offset);  
int munmap(void *addr);
```

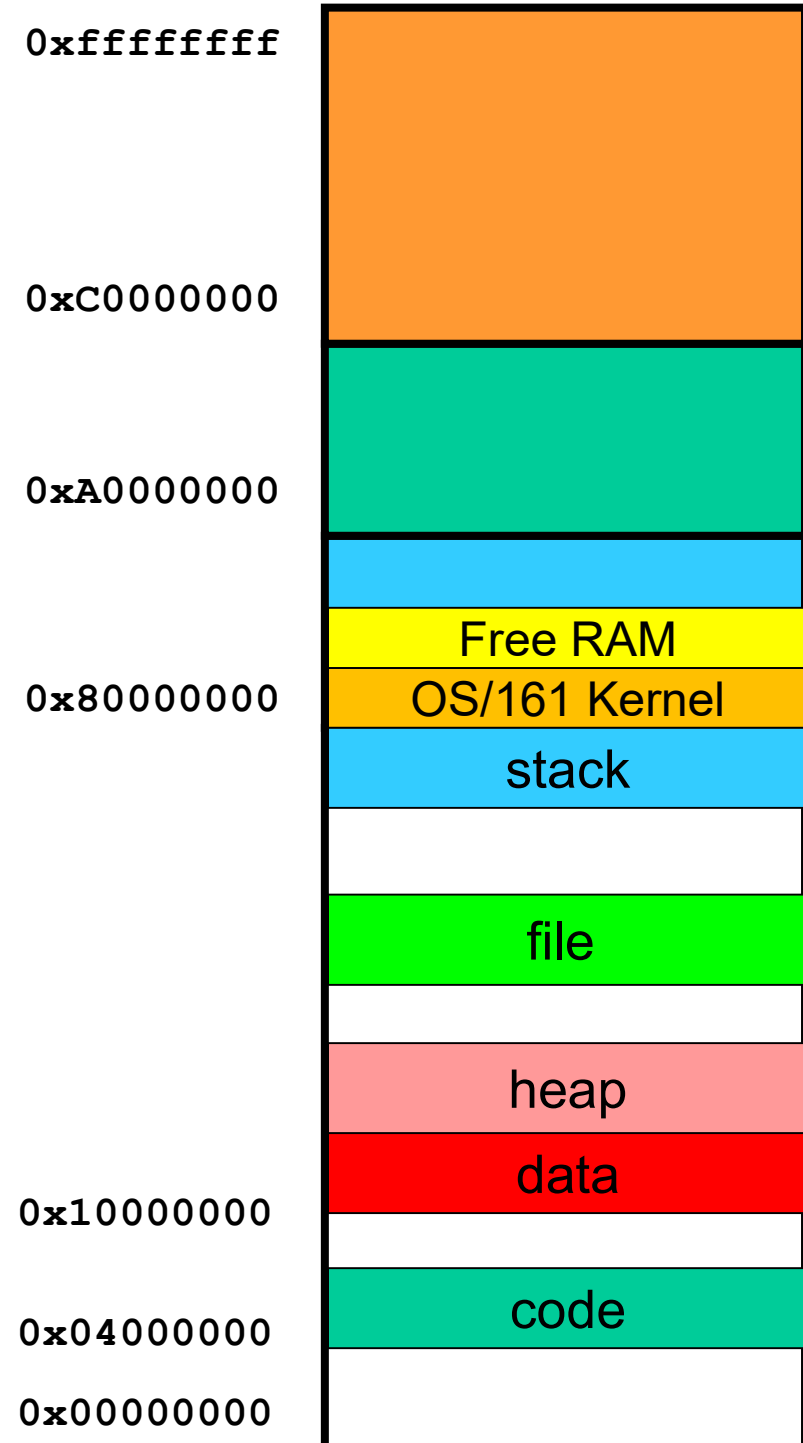


mmap

```
void *mmap(size_t length, int prot, int fd, off_t offset);
```



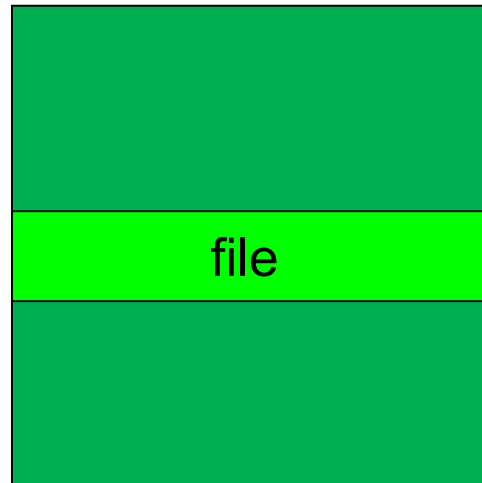
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munmap

```
int munmap(void *addr);
```

File



0xffffffff

0xC0000000

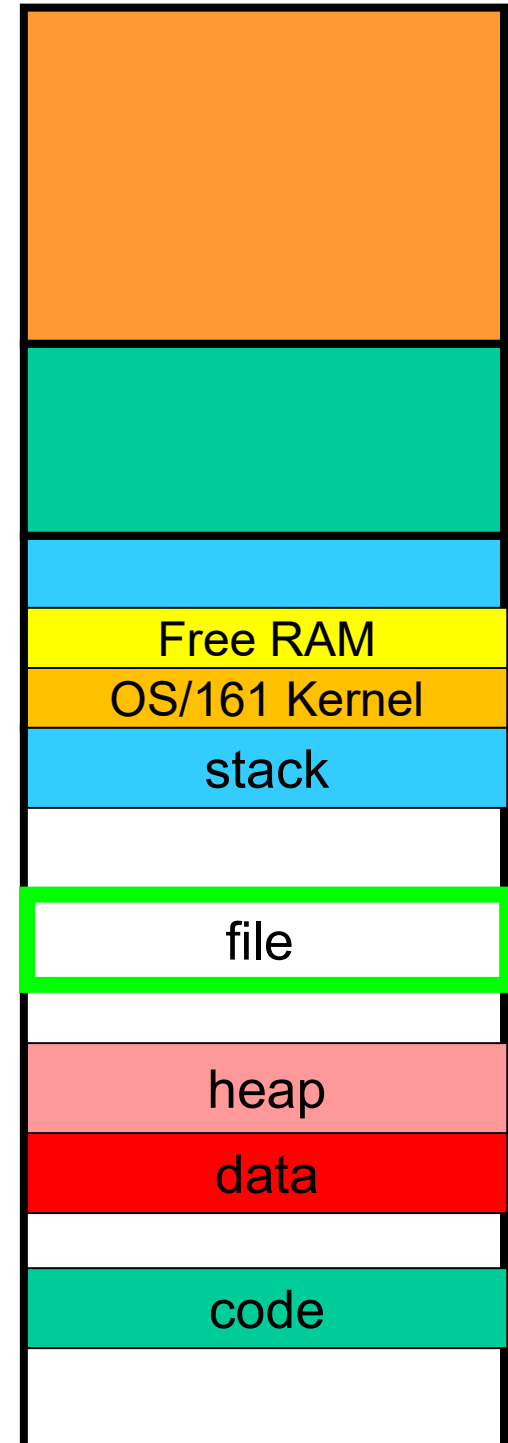
0xA0000000

0x80000000

0x10000000

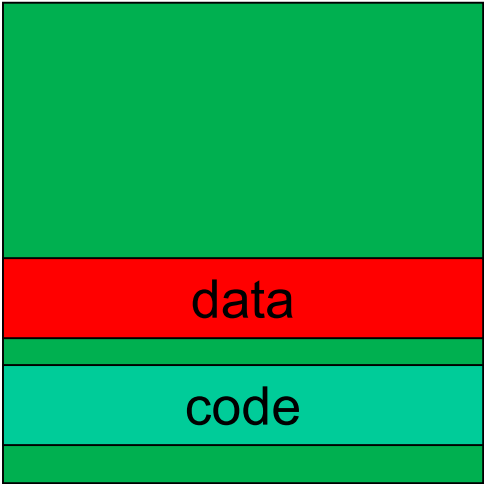
0x04000000

0x00000000

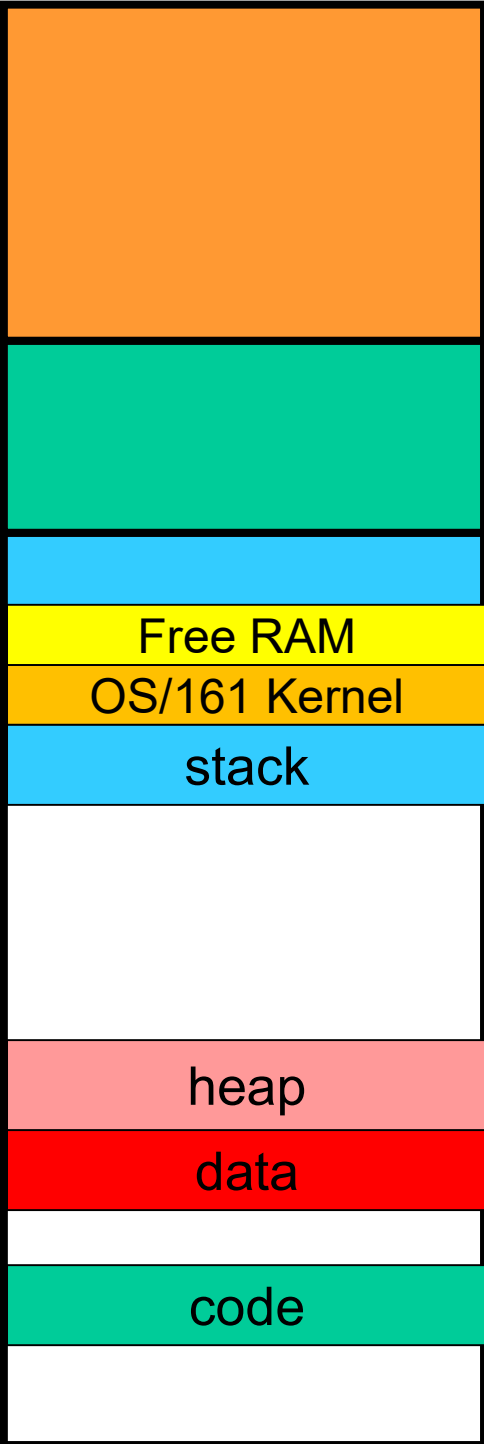


demand loading

Executable
file



0xffffffff
0xc0000000
0xa0000000
0x80000000
0x10000000
0x04000000
0x00000000





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