COMP 4161
NICTA Advanced Course

Advanced Topics in Software Verification

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Isar

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Content

[a] a1 out; [b] a1 due; [c] a2 out; [d] a2 due; [e] session break; [f] a3 out; [g] a3 due

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ISAR

A LANGUAGE FOR STRUCTURED PROOFS

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Isar

apply scripts	**What about..**
➜ unreadable	➜ Elegance?
➜ hard to maintain	➜ Explaining deeper insights?
➜ do not scale	➜ Large developments?

No structure. **Isar!**

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A typical Isar proof

proof
  assume $formula_0$
  have $formula_1$ **by** simp
  ⋮
  have $formula_n$ **by** blast
  show $formula_{n+1}$ **by** ...
qed

proves $formula_0 \Longrightarrow formula_{n+1}$

(analogous to **assumes**/**shows** in lemma statements)

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Isar core syntax

proof = **proof** [method] statement* **qed**
  | **by** method

method = (simp ...) | (blast ...) | (rule ...) | ...

statement = **fix** variables (⋀)
  | **assume** proposition (⟹)
  | [**from** name+] (**have** | **show**) proposition proof
  | **next** (separates subgoals)

proposition = [name:] formula

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proof and qed

proof [method] statement* **qed**

lemma "$[\![A; B]\!] \Longrightarrow A \wedge B$"
proof (rule conjI)
  assume A: "A"
  from A **show** "A" **by** assumption
next
  assume B: "B"
  from B **show** "B" **by** assumption
qed

- ➜ **proof** (<method>) applies method to the stated goal
- ➜ **proof** applies a single rule that fits
- ➜ **proof -** does nothing to the goal

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How do I know what to Assume and Show?

Look at the proof state!

lemma "$[\![A; B]\!] \Longrightarrow A \wedge B$"
proof (rule conjI)

- ➜ **proof** (rule conjI) changes proof state to
  1. $[\![A; B]\!] \Longrightarrow A$
  2. $[\![A; B]\!] \Longrightarrow B$
- ➜ so we need 2 shows: **show** "A" and **show** "B"
- ➜ We are allowed to **assume** A,
  because A is in the assumptions of the proof state.

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The Three Modes of Isar

- ➜ **[prove]**:
  goal has been stated, proof needs to follow.
- ➜ **[state]**:
  proof block has openend or subgoal has been proved,
  new *from* statement, goal statement or assumptions can follow.
- ➜ **[chain]**:
  from statement has been made, goal statement needs to follow.

lemma "$[\![A; B]\!] \Longrightarrow A \wedge B$" **[prove]**
proof (rule conjI) **[state]**
  assume A: "A" **[state]**
  from A **[chain]** **show** "A" **[prove]** **by** assumption **[state]**
next **[state]** **. . .**

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Have

Can be used to make intermediate steps.

Example:

lemma "$(x :: \mathsf{nat}) + 1 = 1 + x$"
proof -
  have A: "$x + 1 = \mathsf{Suc}\ x$" **by** simp
  have B: "$1 + x = \mathsf{Suc}\ x$" **by** simp
  show "$x + 1 = 1 + x$" **by** (simp only: A B)
qed

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DEMO

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Backward and Forward

Backward reasoning: . . . **have** "$A \wedge B$" **proof**
- ➜ **proof** picks an **intro** rule automatically
- ➜ conclusion of rule must unify with $A \wedge B$

Forward reasoning: . . .
  assume AB: "$A \wedge B$"
  from AB **have** ". . ." **proof**
- ➜ now **proof** picks an **elim** rule automatically
- ➜ triggered by **from**
- ➜ first assumption of rule must unify with AB

General case: from $A_1 \ldots A_n$ **have** R **proof**
- ➜ first n assumptions of rule must unify with $A_1 \ldots A_n$
- ➜ conclusion of rule must unify with R

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Fix and Obtain

fix $v_1 \ldots v_n$

Introduces new arbitrary but fixed variables
($\sim$ parameters, $\bigwedge$)

obtain $v_1 \ldots v_n$ **where** <prop> <proof>

Introduces new variables together with property

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DEMO

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Fancy Abbreviations

this	=	the previous fact proved or assumed
then	=	**from** this
thus	=	**then show**
hence	=	**then have**
with $A_1 \ldots A_n$	=	**from** $A_1 \ldots A_n$ this
?thesis	=	the last enclosing goal statement

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Moreover and Ultimately

have X_1: P_1 ...	**have** P_1 ...
have X_2: P_2 ...	**moreover have** P_2 ...
⋮	⋮
have X_n: P_n ...	**moreover have** P_n ...
from $X_1 \ldots X_n$ **show** ...	**ultimately show** ...

wastes lots of brain power
on names $X_1 \ldots X_n$

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General Case Distinctions

show $formula$
proof -
  have $P_1 \vee P_2 \vee P_3$ <proof>
  moreover { **assume** P_1 ... **have** ?thesis <proof> }
  moreover { **assume** P_2 ... **have** ?thesis <proof> }
  moreover { **assume** P_3 ... **have** ?thesis <proof> }
  ultimately show ?thesis **by** blast
qed

{ ... } is a proof block similar to **proof** ... **qed**

{ **assume** P_1 ... **have** P <proof> }
stands for $P_1 \Longrightarrow P$

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Mixing proof styles

from ...
have ...
  apply - make incoming facts assumptions
  apply (...)
  ⋮
  apply (...)
  done

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