

# COMP4161 S2/2016

## Advanced Topics in Software Verification

### Assignment 1

This assignment starts on Mon, 2016-08-08 and is due on Mon, 2016-08-15, 23:59h. We will accept plain text (.txt) files, PDF (.pdf) files, and Isabelle theory (.thy) files.

The assignment is take-home. This does NOT mean you can work in groups. Each submission is personal. For more information, see the plagiarism policy: <https://student.unsw.edu.au/plagiarism>

Submit using `give` on a CSE machine:

```
give cs4161 a1 files ...
```

For example:

```
give cs4161 a1 a1.thy a1.pdf
```

## 1 Types (15 marks)

Construct a type derivation tree for the term  $\lambda a b c. a (x b b) (c b)$ .

Each node of the tree should correspond to the application of a *single* typing rule, indicating which typing rule is used at each step.

Under which contexts is the term type correct?

## 2 $\lambda$ -Calculus (20 marks)

Recall the encoding of booleans and booleans operations in lambda calculus seen in the lecture:

```
true  ≡ λx y. x
false ≡ λx y. y
if    ≡ λz x y. z x y
or    ≡ λx y. if x true y
```

- (a) Show that the  $\beta$  normal form for `or true false` is `true`. Justify your answer by providing the  $\beta$  reduction steps leading from the term to its normal form. Each step should only reduce *one* redex (i.e. one reduction per step). Ideally, you would underline the redex being reduced. (10 marks)

- (b) Provide a type for `true`. Justify your answer by providing a derivation tree. (5 marks)
- (c) What is a type of `or true false`? Justify your answer. (5 marks)

### 3 Propositional Logic (25 marks)

Prove each of the following statements, using only the proof methods `rule`, `erule`, `assumption`, `frule`, and `drule`; and using only the proof rules `impI`, `impE`, `conjI`, `conjE`, `disjI1`, `disjI2`, `disjE`, `notI`, `notE`, `iffI`, `iffE`, `iffD1`, `iffD2`, `ccontr`, `classical`, `FalseE`, `TrueI`, `conjunct1`, `conjunct2`, and `mp`. You do not need to use all of these methods and rules.

- (a)  $A \wedge B \longrightarrow B$  (2 marks)
- (b)  $(P \vee P) = P$  (3 marks)
- (c)  $\neg \neg P \longrightarrow P$  (3 marks)
- (d)  $(A \wedge B \longrightarrow C) = (A \longrightarrow B \longrightarrow C)$  (5 marks)
- (e)  $(\neg x) = (x = \text{False})$  (5 marks)
- (f)  $(a \longrightarrow b) = (\neg (a \wedge \neg b))$  (5 marks)

List the statements above that are provable only in a classical logic. (2 marks)

### 4 Higher-Order Logic (40 marks)

Prove each of the following statements, using only the proof methods and rules from Question 3 plus you may also use the additional methods `rule_tac`, `erule_tac`, `drule_tac`, `frule_tac`, `case_tac`, and `rename_tac`, and the additional rules `allI`, `allE`, `exI`, `exE`, and `spec`. You may use rules proved in earlier parts of the question when proving later parts.

- (a)  $(\forall x y. R x y) \longrightarrow (\forall y x. R x y)$  (3 marks)
- (b)  $(\exists x. P x \wedge Q x) \longrightarrow (\exists x. P x) \wedge (\exists x. Q x)$  (3 marks)
- (c)  $\text{False} = (\forall P. P)$  (4 marks)
- (d)  $\exists R. (\forall x. \exists y. R x y) \wedge \neg (\exists y. \forall x. R x y)$  (6 marks)
- (e)  $\forall x. \neg R x \longrightarrow RM x \implies \forall x. \neg RM x \longrightarrow R x$  (6 marks)

(f)  $\llbracket \forall x. \neg R x \longrightarrow R (M x); \exists x. R x \rrbracket \implies \exists x. R x \wedge R (M (M x))$   
(8 marks)

Formalise and prove the following statement using only the proof methods and rules as earlier in this question. (10 marks)

If every poor person has a rich mother, then there is a rich person with a rich grandmother.