



COMP4161: Advanced Topics in Software Verification

INV

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^aa1 due; ^ba2 due; ^ca3 due

Today



Practice with invariants!

Recall:

- it needs to be an invariant
- it needs to be enough

Example 1



```
A := 0;  
B := 0;  
  
WHILE A  $\neq$  a  
DO  
    B := B + b;  
    A := A + 1  
OD
```

Example 1



```
A := 0;  
B := 0;  
  
WHILE A  $\neq$  a  
DO  
    B := B + b;  
    A := A + 1  
OD  
  
{ B = b * a }
```

Example 1



$\{ a \geq 0 \wedge b \geq 0 \}$

$A := 0;$

$B := 0;$

WHILE $A \neq a$

DO

$B := B + b;$

$A := A + 1$

OD

$\{ B = b * a \}$

Example 1



$\{ a \geq 0 \wedge b \geq 0 \}$

$A := 0;$

$B := 0;$

$\text{INV } \{ B = b * A \}$

$\text{WHILE } A \neq a$

DO

$\quad B := B + b;$

$\quad A := A + 1$

OD

$\{ B = b * a \}$

Example 1

$\{ a \geq 0 \wedge b \geq 0 \}$

$A := 0;$

$B := 0;$

$\text{INV } \{ B = b * A \}$

$\text{WHILE } A \neq a$

DO

$B := B + b;$

$A := A + 1$

OD

$\{ B = b * a \}$

$$0 = b * 0$$

$$B = b * A \wedge A \neq a \longrightarrow B + b = b * (A + 1)$$

$$B = b * A \wedge A = a \longrightarrow B = b * a$$

Example 2



$A := 0;$

$B := 0;$

WHILE $A < a$

DO

$B := B + b;$

$A := A + 1$

OD

Example 2



```
A := 0;  
B := 0;  
  
WHILE A < a  
DO  
    B := B + b;  
    A := A + 1  
OD  
{ B = b * a }
```

Example 2



$\{ a \geq 0 \wedge b \geq 0 \}$

$A := 0;$

$B := 0;$

WHILE $A < a$

DO

$B := B + b;$

$A := A + 1$

OD

$\{ B = b * a \}$

Example 2



```
{  $a \geq 0 \wedge b \geq 0$  }  
A := 0;  
B := 0;  
INV {  $B = b * A$  }  
WHILE  $A < a$   
DO  
     $B := B + b$ ;  
     $A := A + 1$   
OD  
{  $B = b * a$  }
```

Example 2

$\{ a \geq 0 \wedge b \geq 0 \}$

$A := 0;$

$B := 0;$

$\text{INV } \{ B = b * A \}$

$\text{WHILE } A < a$

DO

$B := B + b;$

$A := A + 1$

OD

$\{ B = b * a \}$

$$0 = b * 0$$

$$B = b * A \wedge A < a \longrightarrow B + b = b * (A + 1)$$

$$B = b * A \wedge A \geq a \longrightarrow B = b * a$$

Example 2

$\{ a \geq 0 \wedge b \geq 0 \}$

$A := 0;$

$B := 0;$

$\text{INV } \{ B = b * A \} \wedge A \leq a$

WHILE $A < a$

DO

$B := B + b;$

$A := A + 1$

OD

$\{ B = b * a \}$

$$0 = b * 0 \wedge 0 \leq a$$

$$B = b * A \wedge A < a \longrightarrow B + b = b * (A + 1) \\ \wedge A \leq a \qquad \wedge A + 1 \leq a$$

$$B = b * A \wedge A \geq a \longrightarrow B = b * a \\ \wedge A \leq a$$

Example 3



$A := a;$

$B := 1;$

WHILE $A \neq 0$

DO

$B := B * b;$

$A := A - 1$

OD

Example 3



$A := a;$

$B := 1;$

WHILE $A \neq 0$

DO

$B := B * b;$

$A := A - 1$

OD

$\{ B = b^a \}$

Example 3



$\{ a \geq 0 \wedge b \geq 0 \}$

$A := a;$

$B := 1;$

WHILE $A \neq 0$

DO

$B := B * b;$

$A := A - 1$

OD

$\{ B = b^a \}$

Example 3



$\{ a \geq 0 \wedge b \geq 0 \}$

$A := a;$

$B := 1;$

$\text{INV } \{ B = b^{a-A} \}$

WHILE $A \neq 0$

DO

$B := B * b;$

$A := A - 1$

OD

$\{ B = b^a \}$

Example 3

$\{ a \geq 0 \wedge b \geq 0 \}$

$A := a;$

$B := 1;$

$\text{INV } \{ B = b^{a-A} \}$

WHILE $A \neq 0$

DO

$B := B * b;$

$A := A - 1$

OD

$\{ B = b^a \}$

$$1 = b^{a-a}$$

$$B = b^{a-A} \wedge A \neq 0 \longrightarrow B * b = b^{a-(A-1)}$$

$$B = b^{a-A} \wedge A = 0 \longrightarrow B = b^a$$

Example 4



$X := x;$

$Y := [];$

WHILE $X \neq []$

DO

$Y := (hd\ X \# Y);$

$X := tl\ X$

OD

Example 4



$X := x;$

$Y := [];$

WHILE $X \neq []$

DO

$Y := (hd\ X \# Y);$

$X := tl\ X$

OD

$\{ Y = rev\ x \}$

Example 4



$\{ \textit{True} \}$

$X := x;$

$Y := [];$

WHILE $X \neq []$

DO

$Y := (\textit{hd } X \# Y);$

$X := \textit{tl } X$

OD

$\{ Y = \textit{rev } x \}$

Example 4



```
{ True }  
X := x;  
Y := [];  
INV { (rev X)@Y = rev x }  
WHILE X ≠ []  
  
DO  
  Y := (hd X#Y);  
  X := tl X  
OD  
{ Y = rev x }
```

Example 4

```
{ True }  
X := x;  
Y := [];      (rev x)@[] = rev x  
INV { (rev X)@Y = rev x }  
WHILE X ≠ []  
    (rev X)@Y = rev x ∧ X ≠ [] →  
    (rev (tl X))@((hd X)#Y) = rev x  
DO  
    Y := (hd X)#Y;  
    X := tl X  
OD      (rev X)@Y = rev x ∧ X = [] → Y = rev x  
{ Y = rev x }
```


Example 5



$A := a; B := b; C := 1;$

WHILE $B \neq 0$
DO

 WHILE $(B \bmod 2 = 0)$

 DO

$A := A * A;$

$B := B \text{ div } 2;$

 OD

$C := C * A;$

$B := B - 1$

 OD

Example 5

Try with $b = 10 = 2^1 + 2^3$ or $b = 12 = 2^2 + 2^3$ (and e.g. $a=3$)



```
A := a; B := b; C := 1;
```

```
WHILE B ≠ 0  
DO
```

```
    WHILE (B mod 2 = 0)
```

```
        DO
```

```
            A := A * A;
```

```
            B := B div 2;
```

```
        OD
```

```
        C := C * A;
```

```
        B := B - 1
```

```
    OD
```

Example 5



Try with $b = 10 = 2^1 + 2^3$ or $b = 12 = 2^2 + 2^3$ (and e.g. $a=3$)

$\{ a \geq 0 \wedge b \geq 0 \}$

$A := a; B := b; C := 1;$

WHILE $B \neq 0$

DO

WHILE $(B \bmod 2 = 0)$

DO

$A := A * A;$

$B := B \text{ div } 2;$

OD

$C := C * A;$

$B := B - 1$

OD

$\{ C = a^b \}$

Example 5



Try with $b = 10 = 2^1 + 2^3$ or $b = 12 = 2^2 + 2^3$ (and e.g. $a=3$)

$\{ a \geq 0 \wedge b \geq 0 \}$

$A := a; B := b; C := 1;$

$\text{INV } \{ a^b = C * A^B \}$

WHILE $B \neq 0$

DO

$\text{INV } \{ a^b = C * A^B \}$

WHILE $(B \bmod 2 = 0)$

DO

$A := A * A;$

$B := B \text{ div } 2;$

OD

$C := C * A;$

$B := B - 1$

OD

$\{ C = a^b \}$

Example 5



Try with $b = 10 = 2^1 + 2^3$ or $b = 12 = 2^2 + 2^3$ (and e.g. $a=3$)

$\{ a \geq 0 \wedge b \geq 0 \}$

$A := a; B := b; C := 1;$

$$a^b = 1 * a^b$$

INV $\{ a^b = C * A^B \}$

WHILE $B \neq 0$

$$a^b = C * A^B \wedge B \neq 0 \longrightarrow a^b = (C * A) * A^{B-1}$$

DO

INV $\{ a^b = C * A^B \}$

WHILE $(B \bmod 2 = 0)$

$$a^b = C * A^B \wedge B \bmod 2 = 0 \longrightarrow a^b = C * (A * A)^{B \div 2}$$

DO

$A := A * A;$

$B := B \div 2;$

OD

$C := C * A;$

$B := B - 1$

OD

$$a^b = C * A^B \wedge B = 0 \longrightarrow C = a^b$$

$\{ C = a^b \}$

Example 6



$l := 0; u := \text{length } A - 1; A := a$

WHILE $l \leq u$
DO

 WHILE $l < \text{length } A \wedge A[l] \leq \text{piv}$ DO $l := l + 1$ OD;

 WHILE $0 < u \wedge \text{piv} \leq A[u]$ DO $u := u - 1$ OD;

 IF $l \leq u$ THEN $A := A[l := A[u], u := A[l]]$ ELSE SKIP FI
OD

Example 6



$l := 0; u := \text{length } A - 1; A := a$

WHILE $l \leq u$
DO

 WHILE $l < \text{length } A \wedge A[l] \leq \text{piv}$ DO $l := l + 1$ OD;

 WHILE $0 < u \wedge \text{piv} \leq A[u]$ DO $u := u - 1$ OD;

 IF $l \leq u$ THEN $A := A[l := A[u], u := A[l]]$ ELSE SKIP FI
OD

$\{ \text{LEQ } A \ u \wedge \text{EQ } A \ u \ l \wedge \text{GEQ } A \ l \wedge A \text{ permutes } a \}$

Example 6



$LEQ\ A\ n = \forall k. k < n \longrightarrow A!k \leq piv$

$GEQ\ A\ n = \forall k. n < k < length\ A \longrightarrow A!k \geq piv$

$EQ\ A\ n\ m = \forall k. n \leq k \leq m \longrightarrow A!k = piv$

$\{ 0 < length\ A \}$

$l := 0; u := length\ A - 1; A := a$

WHILE $l \leq u$

DO

WHILE $l < length\ A \wedge A!l \leq piv$ DO $l := l + 1$ OD;

WHILE $0 < u \wedge piv \leq A!u$ DO $u := u - 1$ OD;

IF $l \leq u$ THEN $A := A[l := A!u, u := A!l]$ ELSE SKIP FI

OD

$\{ LEQ\ A\ u \wedge EQ\ A\ u\ l \wedge GEQ\ A\ l \wedge A\ \text{permutes}\ a \}$

Example 6



$LEQ\ A\ n = \forall k. k < n \longrightarrow A!k \leq piv$

$GEQ\ A\ n = \forall k. n < k < length\ A \longrightarrow A!k \geq piv$

$EQ\ A\ n\ m = \forall k. n \leq k \leq m \longrightarrow A!k = piv$

$\{ 0 < length\ A \}$

$l := 0; u := length\ A - 1; A := a$

$INV\ \{ LEQ\ A\ l \wedge GEQ\ A\ u \wedge u < length\ A \wedge l \leq length\ A \wedge A\ \text{permutes}\ a \}$

WHILE $l \leq u$

DO

$INV\ \{ LEQ\ A\ l \wedge GEQ\ A\ u \wedge u < length\ A \wedge l \leq length\ A \wedge A\ \text{permutes}\ a \}$

WHILE $l < length\ A \wedge A!l \leq piv$ DO $l := l + 1$ OD;

$INV\ \{ LEQ\ A\ l \wedge GEQ\ A\ u \wedge u < length\ A \wedge l \leq length\ A \wedge A\ \text{permutes}\ a \}$

WHILE $0 < u \wedge piv \leq A!u$ DO $u := u - 1$ OD;

IF $l \leq u$ THEN $A := A[l := A!u, u := A!l]$ ELSE SKIP FI

OD

$\{ LEQ\ A\ u \wedge EQ\ A\ u\ l \wedge GEQ\ A\ l \wedge A\ \text{permutes}\ a \}$

Example 7

Reminder:

datatype ref = Ref int | Null

Pointer access: $p \rightarrow \text{field}$

Pointer update: $p \rightarrow \text{field} ::= v$

Definition:

"*List* *nxt* *p* *Ps*" is a linked list, starting at pointer *p* following the next pointer through the function *nxt*, and where *Ps* contains the list of the pointers of the linked list.

$\{ \text{List } \text{nxt } p \ Ps \wedge X \in Ps \}$

WHILE $p \neq \text{Null} \wedge p \neq \text{Ref } X$

DO

$p := p \rightarrow \text{nxt};$

OD



Example 7

Reminder:

datatype ref = Ref int | Null

Pointer access: $p \rightarrow \text{field}$

Pointer update: $p \rightarrow \text{field} ::= v$

Definition:

"*List* *nxt* *p* *Ps*" is a linked list, starting at pointer *p* following the next pointer through the function *nxt*, and where *Ps* contains the list of the pointers of the linked list.

$\{ \text{List } \textit{nxt } p \textit{ Ps} \wedge X \in \textit{Ps} \}$

WHILE $p \neq \text{Null} \wedge p \neq \text{Ref } X$

DO

$p := p \rightarrow \textit{nxt};$

OD

$\{ p = \text{Ref } X \}$



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Definition:

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$\{ \text{List } \textit{nxt } p \textit{ Ps} \wedge X \in \textit{Ps} \}$

WHILE $p \neq \text{Null} \wedge p \neq \text{Ref } X$

DO

$p := p \rightarrow \textit{nxt};$

OD

$\{ p = \text{Ref } X \}$



Example 7

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Pointer access: $p \rightarrow \text{field}$

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Definition:

"*List* *nxt* *p* *Ps*" is a linked list, starting at pointer *p* following the next pointer through the function *nxt*, and where *Ps* contains the list of the pointers of the linked list.

$\{ \text{List } \text{nxt } p \ Ps \wedge X \in Ps \}$

INV $\{ \exists Qs. \text{List } \text{nxt } p \ Qs \wedge X \in Qs \}$

WHILE $p \neq \text{Null} \wedge p \neq \text{Ref } X$

DO

$p := p \rightarrow \text{nxt};$

OD

$\{ p = \text{Ref } X \}$



Example 7

Reminder:

datatype ref = Ref int | Null

Pointer access: $p \rightarrow \text{field}$

Pointer update: $p \rightarrow \text{field} ::= v$

Definition:

"*List* *nxt* *p* *Ps*" is a linked list, starting at pointer *p* following the next pointer through the function *nxt*, and where *Ps* contains the list of the pointers of the linked list.



$\{ \text{List } \text{nxt } p \ Ps \wedge X \in Ps \}$	$\exists Qs. \text{List } \text{nxt } p \ Qs \wedge X \in Qs$
INV $\{ \exists Qs. \text{List } \text{nxt } p \ Qs \wedge X \in Qs \}$	
WHILE $p \neq \text{Null} \wedge p \neq \text{Ref } X$	$\exists Qs. \text{List } \text{nxt } p \ Qs \wedge X \in Qs$ $\wedge p \neq \text{Null} \wedge p \neq \text{Ref } X \rightarrow$ $\exists Qs. \text{List } \text{nxt } (p \rightarrow \text{nxt}) \ Qs \wedge X \in Qs$
DO	
$p := p \rightarrow \text{nxt};$	
OD	$\exists Qs. \text{List } \text{nxt } p \ Qs \wedge X \in Qs$ $\wedge (p = \text{Null} \vee p = \text{Ref } X) \rightarrow p = \text{Ref } X$
$\{ p = \text{Ref } X \}$	

Example 8



What is is Isabelle function doing?

fun f :: 'a list \Rightarrow ' a list \Rightarrow ' a list where
f [] ys = ys|
f xs [] = xs|
f (x#xs) (y#ys) = x#y# f xs ys

Example 8



What is Isabelle function doing?

```
fun splice :: 'a list  $\Rightarrow$  ' a list  $\Rightarrow$  ' a list where
  splice [] ys = ys|
  splice xs [] = xs|
  splice (x#xs) (y#ys) = x#y# f xs ys
```


Example 8



What is Isabelle function doing?

```
fun splice :: 'a list  $\Rightarrow$  'a list  $\Rightarrow$  'a list where
  splice [] ys = ys|
  splice xs [] = xs|
  splice (x#xs) (y#ys) = x#y# f xs ys
```

Let's write it with linked lists!

Example 8



$$\{ \text{List } \text{nxt } p \ Ps \wedge \text{List } \text{nxt } q \ Qs \wedge (\text{set } Ps \cap \text{set } Qs) = \{\} \wedge \text{size } Qs \leq \text{size } Ps \}$$

$$\{ \text{List } \text{nxt } p \ (\text{splice } Ps \ Qs) \}$$

Example 8



$\{ \text{List } \text{next } p \ Ps \wedge \text{List } \text{next } q \ Qs \wedge (\text{set } Ps \cap \text{set } Qs) = \{\} \wedge \text{size } Qs \leq \text{size } Ps \}$
 $pp := p;$

WHILE $q \neq \text{Null}$
DO
 $qq := q \rightarrow \text{next}; q \rightarrow \text{next} := pp \rightarrow \text{next}; pp \rightarrow \text{next} = q; pp := q \rightarrow \text{next}; q := qq;$
OD
 $\{ \text{List } \text{next } p \ (\text{splice } Ps \ Qs) \}$

Example 8



List *nxt* *p* *Ps* = *Path* *nxt* *p* *Ps* *Null*

Path *nxt* *p* *Ps* *Null* is a linked list from *p* to *q* following function *nxt* and containing list of pointers *Ps*

$\{ \text{List } \text{nxt } p \ Ps \wedge \text{List } \text{nxt } q \ Qs \wedge (\text{set } Ps \cap \text{set } Qs) = \{\} \wedge \text{size } Qs \leq \text{size } Ps \}$

pp := *p*;

INV {

}

WHILE *q* ≠ *Null*

DO

qq := *q* → *nxt*; *q* → *nxt* := *pp* → *nxt*; *pp* → *nxt* = *q*; *pp* := *q* → *nxt*; *q* := *qq*;

OD

{ *List* *nxt* *p* (*splice* *Ps* *Qs*) }

Example 8



List *nxt* *p* *Ps* = *Path* *nxt* *p* *Ps* *Null*

Path *nxt* *p* *Ps* *Null* is a linked list from *p* to *q* following function *nxt* and containing list of pointers *Ps*

$\{ \text{List } \text{nxt } p \ Ps \wedge \text{List } \text{nxt } q \ Qs \wedge (\text{set } Ps \cap \text{set } Qs) = \{\} \wedge \text{size } Qs \leq \text{size } Ps \}$

pp := *p*;

INV { $\exists PPs$

List *nxt* *pp* *PPs*

}

WHILE *q* ≠ *Null*

DO

qq := *q* → *nxt*; *q* → *nxt* := *pp* → *nxt*; *pp* → *nxt* = *q*; *pp* := *q* → *nxt*; *q* := *qq*;

OD

{ *List* *nxt* *p* (*splice* *Ps* *Qs*) }

Example 8



List *nxt* *p* *Ps* = *Path* *nxt* *p* *Ps* *Null*

Path *nxt* *p* *Ps* *Null* is a linked list from *p* to *q* following function *nxt* and containing list of pointers *Ps*

$\{ \text{List } \text{nxt } p \ Ps \wedge \text{List } \text{nxt } q \ Qs \wedge (\text{set } Ps \cap \text{set } Qs) = \{\} \wedge \text{size } Qs \leq \text{size } Ps \}$

$pp := p;$

INV $\{ \exists PPs \ QQs$

$\text{List } \text{nxt } pp \ PPs \wedge \text{List } \text{nxt } q \ QQs$

}

WHILE $q \neq \text{Null}$

DO

$qq := q \rightarrow \text{nxt}; q \rightarrow \text{nxt} := pp \rightarrow \text{nxt}; pp \rightarrow \text{nxt} = q; pp := q \rightarrow \text{nxt}; q := qq;$

OD

$\{ \text{List } \text{nxt } p \ (\text{splice } Ps \ Qs) \}$

Example 8



List *nxt* *p* *Ps* = *Path* *nxt* *p* *Ps* *Null*

Path *nxt* *p* *Ps* *Null* is a linked list from *p* to *q* following function *nxt* and containing list of pointers *Ps*

$\{ \text{List } \text{nxt } p \ Ps \wedge \text{List } \text{nxt } q \ Qs \wedge (\text{set } Ps \cap \text{set } Qs) = \{\} \wedge \text{size } Qs \leq \text{size } Ps \}$

$pp := p;$

INV $\{ \exists PPPs \ QQs \ PPPs.$

$\text{List } \text{nxt } pp \ PPs \wedge \text{List } \text{nxt } q \ QQs \wedge \text{Path } \text{nxt } p \ PPPs \ pp$

}

WHILE $q \neq \text{Null}$

DO

$qq := q \rightarrow \text{nxt}; q \rightarrow \text{nxt} := pp \rightarrow \text{nxt}; pp \rightarrow \text{nxt} = q; pp := q \rightarrow \text{nxt}; q := qq;$

OD

$\{ \text{List } \text{nxt } p \ (\text{splice } Ps \ Qs) \}$

Example 8



List *nxt* *p* *Ps* = *Path* *nxt* *p* *Ps* *Null*

Path *nxt* *p* *Ps* *Null* is a linked list from *p* to *q* following function *nxt* and containing list of pointers *Ps*

$\{ \text{List } \text{nxt } p \ Ps \wedge \text{List } \text{nxt } q \ Qs \wedge (\text{set } Ps \cap \text{set } Qs) = \{\} \wedge \text{size } Qs \leq \text{size } Ps \}$

pp := *p*;

INV { $\exists PPs \ QQs \ PPPs.$

$\text{List } \text{nxt } pp \ PPs \wedge \text{List } \text{nxt } q \ QQs \wedge \text{Path } \text{nxt } p \ PPPs \ pp$
 $\wedge \text{PPPs@splice } PPs \ QQs = \text{splice } Ps \ Qs$

}

WHILE *q* ≠ *Null*

DO

qq := *q* → *nxt*; *q* → *nxt* := *pp* → *nxt*; *pp* → *nxt* = *q*; *pp* := *q* → *nxt*; *q* := *qq*;

OD

{ *List* *nxt* *p* (*splice* *Ps* *Qs*) }

Example 8



List *nxt* *p* *Ps* = *Path* *nxt* *p* *Ps* *Null*

Path *nxt* *p* *Ps* *Null* is a linked list from *p* to *q* following function *nxt* and containing list of pointers *Ps*

$\{ \text{List } \text{nxt } p \ Ps \wedge \text{List } \text{nxt } q \ Qs \wedge (\text{set } Ps \cap \text{set } Qs) = \{\} \wedge \text{size } Qs \leq \text{size } Ps \}$

$pp := p;$

INV $\{ \exists PPs \ QQs \ PPPs. \ \text{size } QQs \leq \text{size } PPs \wedge$
 $\text{List } \text{nxt } pp \ PPs \wedge \text{List } \text{nxt } q \ QQs \wedge \text{Path } \text{nxt } p \ PPPs \ pp$
 $\wedge \text{PPPs@splice } PPs \ QQs = \text{splice } Ps \ Qs$

}

WHILE $q \neq \text{Null}$

DO

$qq := q \rightarrow \text{nxt}; q \rightarrow \text{nxt} := pp \rightarrow \text{nxt}; pp \rightarrow \text{nxt} = q; pp := q \rightarrow \text{nxt}; q := qq;$

OD

$\{ \text{List } \text{nxt } p \ (\text{splice } Ps \ Qs) \}$

Example 8



List *nxt* *p* *Ps* = *Path* *nxt* *p* *Ps* *Null*

Path *nxt* *p* *Ps* *Null* is a linked list from *p* to *q* following function *nxt* and containing list of pointers *Ps*

```
{ List nxt p Ps  $\wedge$  List nxt q Qs  $\wedge$  (set Ps  $\cap$  set Qs) = {}  $\wedge$  size Qs  $\leq$  size Ps }  
pp := p;  
INV {  $\exists$  PPs QQs PPPs. size QQs  $\leq$  size PPs  $\wedge$   
      List nxt pp PPs  $\wedge$  List nxt q QQs  $\wedge$  Path nxt p PPPs pp  
       $\wedge$  PPPs@splice PPs QQs = splice Ps Qs  $\wedge$   
      set PPs  $\cap$  set QQs = {}  $\wedge$  distinct PPPs  $\wedge$  set PPPs  $\cap$  (set PPs  $\cup$  set QQs) = {}  
    }  
WHILE q  $\neq$  Null  
DO  
  qq := q  $\rightarrow$  nxt; q  $\rightarrow$  nxt := pp  $\rightarrow$  nxt; pp  $\rightarrow$  nxt = q; pp := q  $\rightarrow$  nxt; q := qq;  
OD  
{ List nxt p (splice Ps Qs) }
```

Demo