## Introduction to Information Retrieval

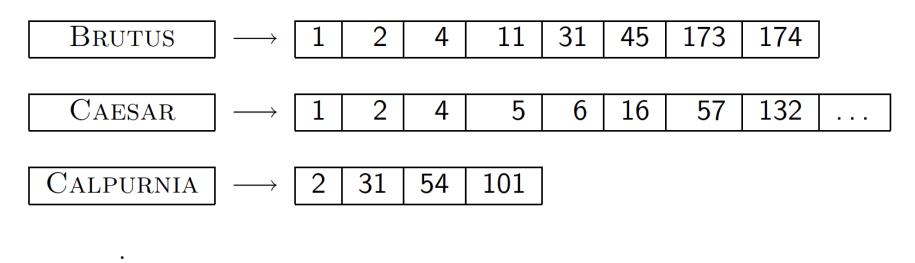
Lecture 4: Dictionaries and tolerant retrieval

#### This lecture

- Dictionary data structures
- "Tolerant" retrieval
  - Wild-card queries
  - Spelling correction

#### Sec. 3.1 Dictionary data structures for inverted indexes

The dictionary data structure stores the term vocabulary, document frequency, pointers to each postings list ... in what data structure?



postings



#### A naïve dictionary

An array of struct:

term	document	pointer to
	frequency	postings list
а	656,265	$\longrightarrow$
aachen	65	$\longrightarrow$
zulu	221	$\longrightarrow$

char[20] intPostings \*20 bytes4/8 bytes4/8 bytes

- How do we store a dictionary in memory efficiently?
- How do we quickly look up elements at query time?

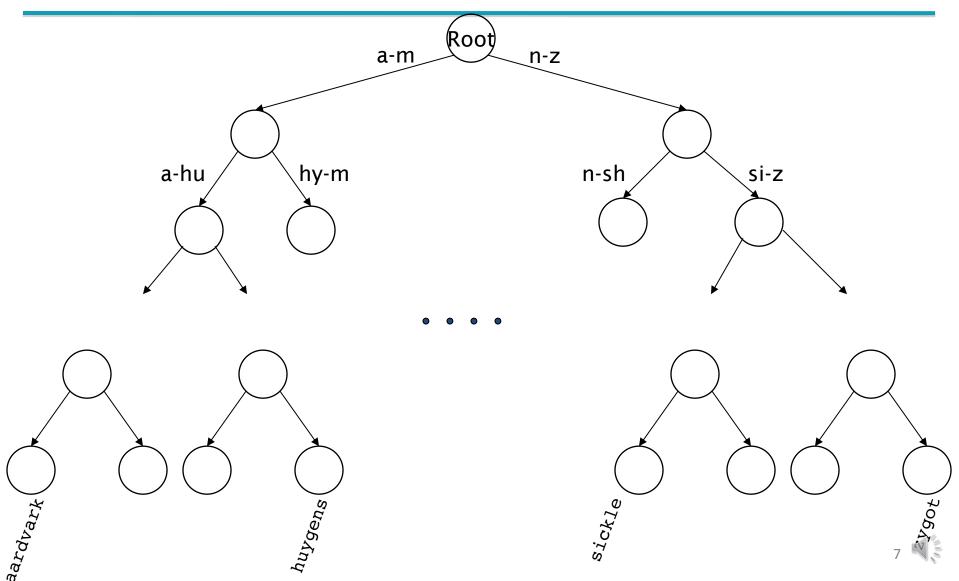
#### Dictionary data structures

- Two main choices:
  - Hash table
  - Tree
- Some IR systems use hashes, some trees

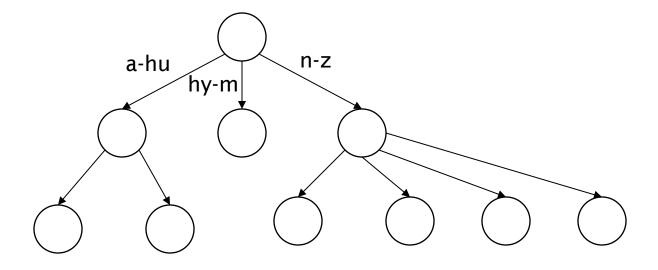
#### Hashes

- Each vocabulary term is hashed to an integer
  - (We assume you've seen hashtables before)
- Pros:
  - Lookup is faster than for a tree: O(1)
- Cons:
  - No easy way to find minor variants:
    - judgment/judgement
  - No prefix search [tolerant retrieval]
  - If vocabulary keeps growing, need to occasionally do the expensive operation of rehashing *everything*

#### Tree: binary tree



#### Tree: B-tree



 Definition: Every internal nodel has a number of children in the interval [*a*,*b*] where *a*, *b* are appropriate natural numbers, e.g., [2,4].

#### Trees

- Simplest: binary tree
- More usual: B-trees
- Trees require a standard ordering of characters and hence strings ... but we standardly have one
- Pros:
  - Solves the prefix problem (terms starting with hyp)
- Cons:
  - Slower: O(log M) [and this requires balanced tree]
  - Rebalancing binary trees is expensive
    - But B-trees mitigate the rebalancing problem

## WILD-CARD QUERIES

#### Wild-card queries: \*

- mon\*: find all docs containing any word beginning "mon".
- Easy with binary tree (or B-tree) lexicon: retrieve all words in range: mon ≤ w < moo</li>
- \*mon: find words ending in "mon": harder
  - Maintain an additional B-tree for terms backwards.
    Can retrieve all words in range: nom ≤ w < non.</li>

Exercise: from this, how can we enumerate all terms meeting the wild-card query *pro\*cent*?

#### Query processing

- At this point, we have an enumeration of all terms in the dictionary that match the wild-card query.
- We still have to look up the postings for each enumerated term.
- E.g., consider the query:

#### se\*ate AND fil\*er

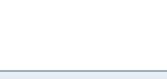
This may result in the execution of many Boolean *AND* queries.

# B-trees handle \*'s at the end of a query term

- How can we handle \*'s in the middle of query term?
  - co\*tion
- We could look up *co\** AND *\*tion* in a B-tree and intersect the two term sets
  - Expensive
  - Still need verification to remove false-positives
- The solution: transform wild-card queries so that the \*'s occur at the end
- This gives rise to the **Permuterm** Index.

#### Permuterm index

- For term *hello*, index under:
  - hello\$, ello\$h, llo\$he, lo\$hel, o\$hell
    where \$ is a special symbol.
- Queries:
  - P Exact match P\$
  - P\* Range match \$P\*
  - \*P Range match P\$\*
  - \*P\* Range match P\* #
  - P\*Q Range match Q\$P\*
  - P\*Q\*R ??? Exercise!



Q: Why not P\*\$\*



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#### Permuterm query processing

- Rotate query wild-card to the right
- Now use B-tree lookup as before.
- Permuterm problem: ≈ quadruples lexicon size



## Bigram (k-gram) indexes

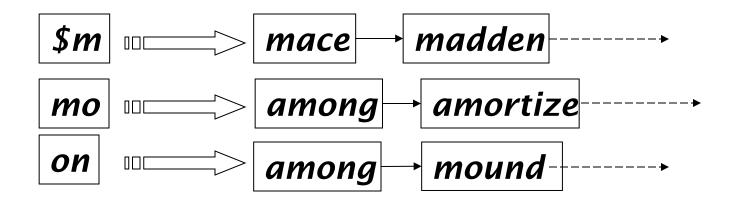
- Enumerate all k-grams (sequence of k chars) occurring in any term
- *e.g.,* from text "*April is the cruelest month*" we get the 2-grams (*bigrams*)

\$a,ap,pr,ri,il,l\$,\$i,is,s\$,\$t,th,he,e\$,\$c,cr,ru, ue,el,le,es,st,t\$, \$m,mo,on,nt,h\$

- \$ is a special word boundary symbol
- Maintain a <u>second</u> inverted index <u>from bigrams to</u> <u>dictionary terms</u> that match each bigram.

#### Bigram index example

The k-gram index finds terms based on a query consisting of k-grams (here k=2).



## Processing wild-cards

- Query mon\* can now be run as
  - \$m AND mo AND on
- Gets terms that match AND version of our wildcard query.
- But we'd enumerate moon.
- Must verify these terms against query.
- Surviving enumerated terms are then looked up in the term-document inverted index.
- Fast, space efficient (compared to permuterm).

#### Processing wild-card queries

- As before, we must execute a Boolean query for each enumerated, filtered term.
- Wild-cards can result in expensive query execution (very large disjunctions...)
  - pyth\* AND prog\*
- If you encourage "laziness" people will respond!

Type your search terms, use '\*' if you need to. E.g., Alex\* will match Alexander. Search

#### Resources

- IIR 3, MG 4.2
- Efficient spell retrieval:
  - K. Kukich. Techniques for automatically correcting words in text. ACM Computing Surveys 24(4), Dec 1992.
  - J. Zobel and P. Dart. Finding approximate matches in large lexicons. Software - practice and experience 25(3), March 1995. <u>http://citeseer.ist.psu.edu/zobel95finding.html</u>
  - Mikael Tillenius: Efficient Generation and Ranking of Spelling Error Corrections. Master's thesis at Sweden's Royal Institute of Technology. <u>http://citeseer.ist.psu.edu/179155.html</u>
- Nice, easy reading on spell correction:
  - Peter Norvig: How to write a spelling corrector

http://norvig.com/spell-correct.html