# COMP4161 T3/2023 <br> Advanced Topics in Software Verification 

## Assignment 1

This assignment starts on Friday 22nd September 2023 and is due on Friday 29th September 2023 23:59:59. We will accept plain text (.txt) files, PDF (.pdf) files, and Isabelle theory (.thy) files. You are allowed to make late submissions up to five days ( 120 hours) after the deadline, but at a cost: -5 marks per day.
The assignment is take-home. This does NOT mean you can work in groups. Each submission is personal. For more information, see the plagiarism policy: https://student.unsw.edu.au/plagiarism Submit using give on a CSE machine:

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give cs4161 a1 files ...
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For example:

```
give cs4161 a1 a1.thy a1.pdf
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## $1 \lambda$-Calculus (16 marks)

(a) Simplify the term $(x y)(\lambda x \cdot(\lambda y \cdot(\lambda z \cdot(z(x y)))))$ syntactically by applying the syntactic conventions and rules. Justify your answer. (2 marks)
(b) Restore the omitted parentheses in the term $x(\lambda x y \cdot x(y z)(x y))(\lambda y \cdot y z)$ (make sure you don't change the term structure). (2 marks)
(c) Find the normal form of $(\lambda f \cdot \lambda x \cdot f(f(f x)))(\lambda g . \lambda y . g(g y))$. Justify your answer by showing the reduction sequence. Each step in the reduction sequence should be a single $\beta$-reduction step. Underline the redex being reduced for each step. (6 marks)
(d) Recall the encoding of natural numbers in lambda calculus (Church Numerals) seen in the lecture:

$$
\begin{aligned}
0 & \equiv \lambda f x . x \\
1 & \equiv \lambda f x \cdot f x \\
2 & \equiv \lambda f x \cdot f(f x) \\
3 & \equiv \lambda f x \cdot f(f(f x)) \ldots
\end{aligned}
$$

Define exp where exp mn beta-reduces to the Church Numeral representing $m^{n}$. Provide a justification of your answer. (6 marks)

## 2 Types (20 marks)

(a) Provide the most general type for the term $\lambda a b . a(c b) b$. Show a type derivation tree to justify your answer. Each node of the tree should correspond to the application of a single typing rule, and be labeled with the typing rule used. Under which contexts is the term type correct? (5 marks)
(b) Find a closed lambda term that has the following type:

(You don't need to provide a type derivation, just the term). (4 marks)
(c) Explain why $\lambda x . x x$ is not typable. (3 marks)
(d) Find the normal form of ( $\lambda x y . y)(\lambda z . z z)$ and give it a type. (3 marks)
(e) Is $(\lambda x y . y)(\lambda z . z z)$ typable? Compare this situation with the Subject Reduction that you learned in the lecture. (5 marks)

## 3 Propositional Logic (29 marks)

Prove each of the following statements, using only the proof methods: rule, erule, assumption, cases, frule, drule, rule_tac, erule_tac, frule_tac, drule_tac, rename_tac, and case_tac; and using only the proof rules: impI, impE, conjI, conjE, disjI1, disjI2, disjE, notI, notE, iffI, iffE, iffD1, iffD2, ccontr, classical, FalseE, TrueI, conjunct1, conjunct2, and mp. You do not need to use all of these methods and rules.
(a) $\mathrm{A} \longrightarrow \neg \neg \mathrm{A}$
(b) $\neg \neg \neg \mathrm{A} \longrightarrow \neg \mathrm{A}$
(c) $\neg \neg \mathrm{A} \longrightarrow \mathrm{A}$
(d) $\neg(A \wedge B) \longrightarrow \neg A \vee \neg B$
(e) $(\mathrm{A} \longrightarrow \mathrm{B}) \longrightarrow \neg \mathrm{A} \vee \mathrm{B}$
(f) $(\neg \mathrm{A} \wedge \neg \mathrm{B})=(\neg(\mathrm{A} \vee \mathrm{B}))$
$(\mathrm{g})(\mathrm{A} \longrightarrow \mathrm{B}) \longrightarrow((\mathrm{B} \longrightarrow \mathrm{C}) \longrightarrow \mathrm{A}) \longrightarrow \mathrm{B}$

## 4 Higher-Order Logic (35 marks)

Prove each of the following statements, using only the proof methods and proof rules stated in the previous question, plus any of the following proof rules: allI, alle, exI, and exE. You do not need to use all of these methods and rules. You may use rules proved in earlier parts of the question when proving later parts.
(a) $(\exists \mathrm{x} \cdot \mathrm{P} \mathrm{x} \longrightarrow \mathrm{Q}) \longrightarrow(\forall \mathrm{x}, \mathrm{P} \mathrm{x}) \longrightarrow \mathrm{Q}$
(b) $((\exists \mathrm{x}, \mathrm{P} \mathrm{x}) \longrightarrow \mathrm{Q})=(\forall \mathrm{x}, \mathrm{P} \mathrm{x} \longrightarrow \mathrm{Q})$
(c) $(\forall \mathrm{x} \cdot \mathrm{P} \mathrm{x})=(\nexists \mathrm{x} . \neg \mathrm{P} \mathrm{x})$
(d) $(\forall \mathrm{x} \cdot \mathrm{P} \mathrm{x} \wedge \mathrm{Q} \mathrm{x}) \longrightarrow(\forall \mathrm{x} \cdot \mathrm{P} \mathrm{x}) \wedge(\forall \mathrm{x} \cdot \mathrm{Q} \mathrm{x})$
(e) $(\exists x \cdot P \mathrm{x} \vee \mathrm{Q} \mathrm{x}) \longrightarrow(\exists \mathrm{x} \cdot \mathrm{P} \mathrm{x}) \vee(\exists \mathrm{x} \cdot \mathrm{Q} \mathrm{x})$
(f) $(\forall x y \cdot A y \vee B x) \longrightarrow(\forall x . B x) \vee(\forall y . A y)$

