Processes and Threads

Major Requirements of an Operating System

- Interleave the execution of several processes to maximize processor utilization while providing reasonable response time
- Allocate resources to processes
- Support interprocess communication and user creation of processes

Processes and Threads

- Processes:
  - Also called a task or job
  - Execution of an individual program
  - "Owner" of resources allocated for program execution
  - Encompasses one or more threads
- Threads:
  - Unit of execution
  - Can be traced
    - list the sequence of instructions that execute
  - Belongs to a process

Execution snapshot of three single-threaded processes (No Virtual Memory)

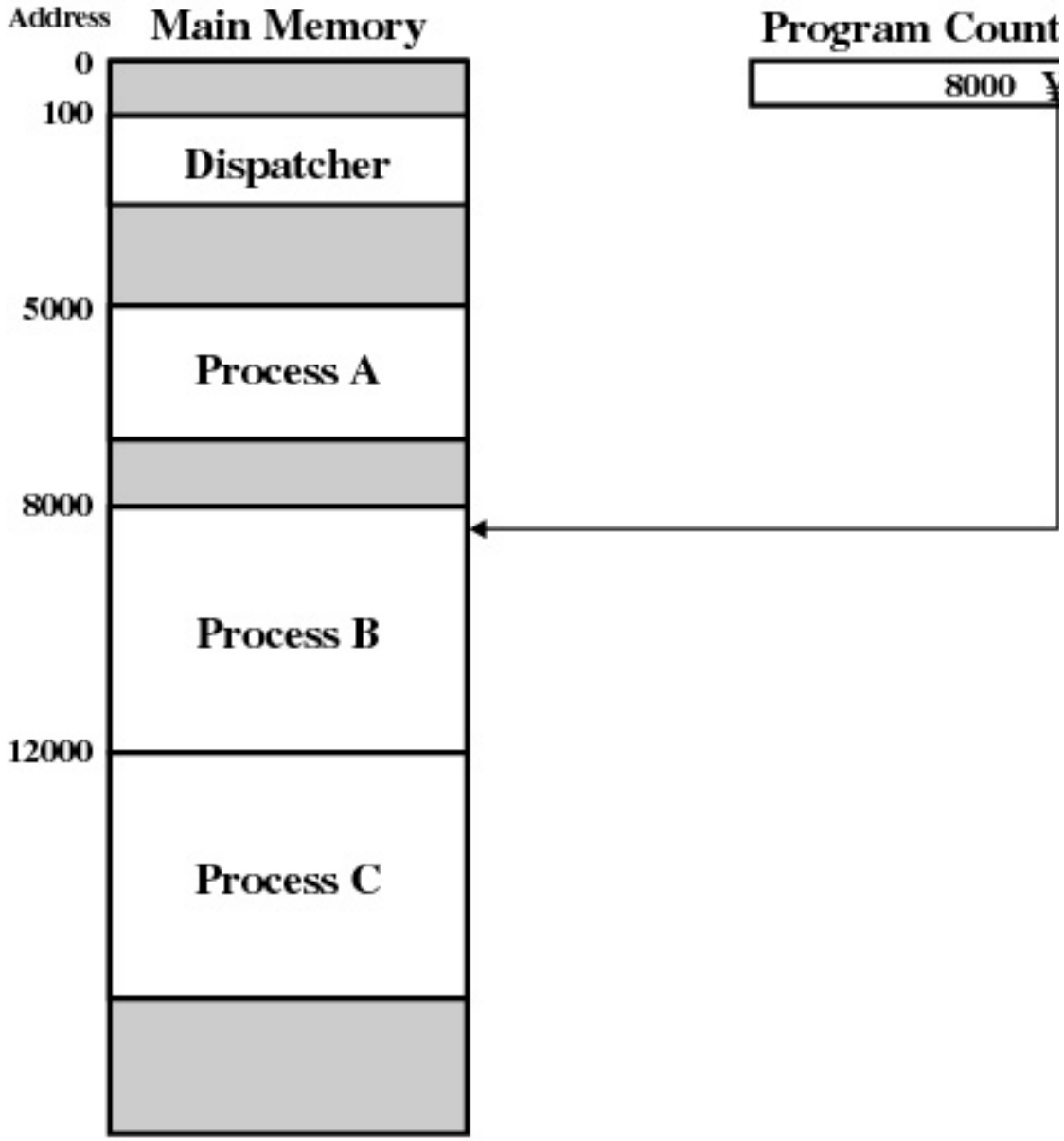


Figure 3.1 Snapshot of Example Execution (Figure 3
at Instruction Cycle 13

Logical Execution Trace

(a) Trace of Process A	(b) Trace of Process B	(c) Trace of Process C
5000	8000	12000
5001	8001	12001
5002	8002	12002
5003	8003	12003
5004		12004
5005		12005
5006		12006
5007		12007
5008		12008
5009		12009
5010		12010
5011		12011

5000 = Starting address of program of Process A
8000 = Starting address of program of Process B
12000 = Starting address of program of Process C

Figure 3.2 Traces of Processes of Figure 3.1

Combined Traces

(Actual CPU Instructions)

What are the shaded sections?

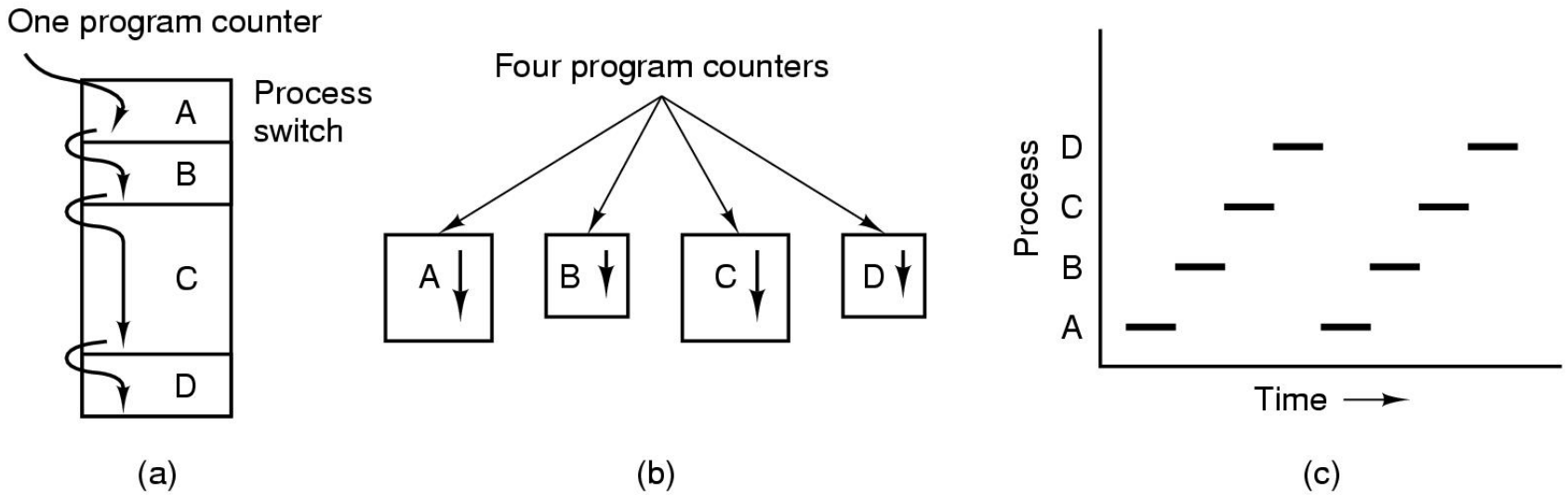
1	5000	27	12004
2	5001	28	12005
3	5002	----------	Time out
4	5003	29	100
5	5004	30	101
6	5005	31	102
----------	Time out	32	103
7	100	33	104
8	101	34	105
9	102	35	5006
10	103	36	5007
11	104	37	5008
12	105	38	5009
13	8000	39	5010
14	8001	40	5011
15	8002	----------	Time out
16	8003	41	100
----------	I/O request	42	101
17	100	43	102
18	101	44	103
19	102	45	104
20	103	46	105
21	104	47	12006
22	105	48	12007
23	12000	49	12008
24	12001	50	12009
25	12002	51	12010
26	12003	52	12011
		----------	Time out

100 = Starting address of dispatcher program

shaded areas indicate execution of dispatcher process;
first and third columns count instruction cycles;
second and fourth columns show address of instruction being executed

Figure 3.3 Combined Trace of Processes of Figure 3.1

Summary: The Process Model



- Multiprogramming of four programs
- Conceptual model of 4 independent, sequential processes (with a single thread each)
- Only one program active at any instant

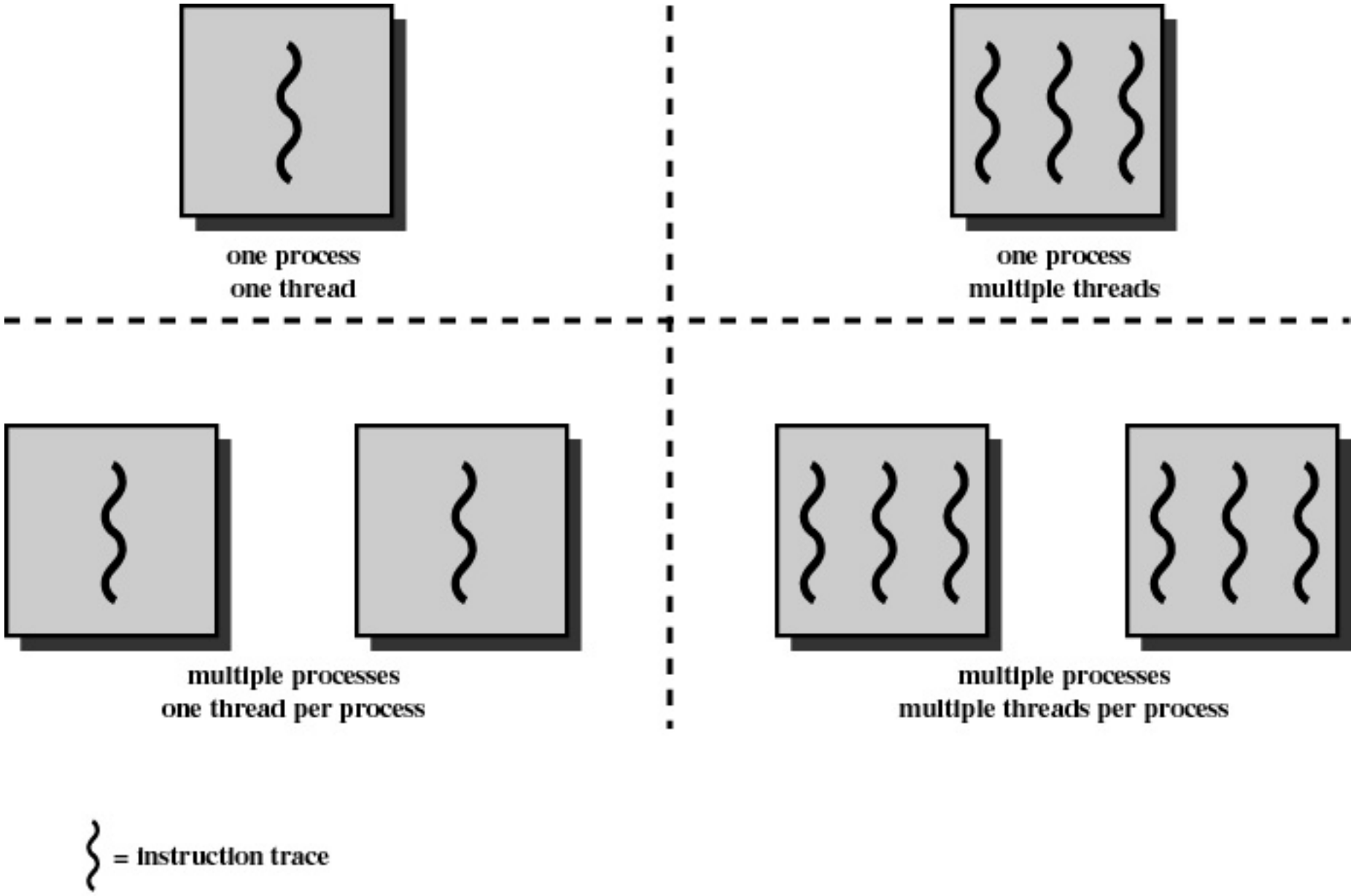


Figure 4.1 Threads and Processes [ANDE97]

Process and thread models of selected OSes

- Single process, single thread
  - MSDOS
- Single process, multiple threads
  - OS/161 as distributed
- Multiple processes, single thread
  - Traditional unix
- Multiple processes, multiple threads
  - Modern Unix (Linux, Solaris), Windows 2000

Note: Literature (incl. Textbooks) often do not cleanly distinguish between processes and threads (for historical reasons)

Process Creation

Principal events that cause process creation

1. System initialization
   - Foreground processes (interactive programs)
   - Background processes
     - Email server, web server, print server, etc.
     - Called a *daemon* (unix) or *service* (Windows)
2. Execution of a process creation system call by a running process
   - New login shell for an incoming telnet connection
3. User request to create a new process
4. Initiation of a batch job

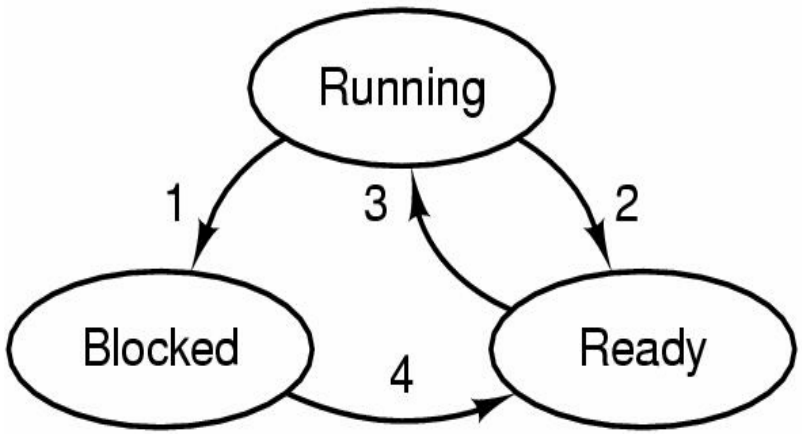
Note: Technically, all these cases use the same system mechanism to create new processes.

Process Termination

Conditions which terminate processes

1. Normal exit (voluntary)
2. Error exit (voluntary)
3. Fatal error (involuntary)
4. Killed by another process (involuntary)

Process/Thread States



1. Process blocks for input
2. Scheduler picks another process
3. Scheduler picks this process
4. Input becomes available

- Possible process/thread states
  - running
  - blocked
  - ready
- Transitions between states shown

Some Transition Causing Events

Running ▷Ready

- Voluntary `Yield()`
- End of timeslice

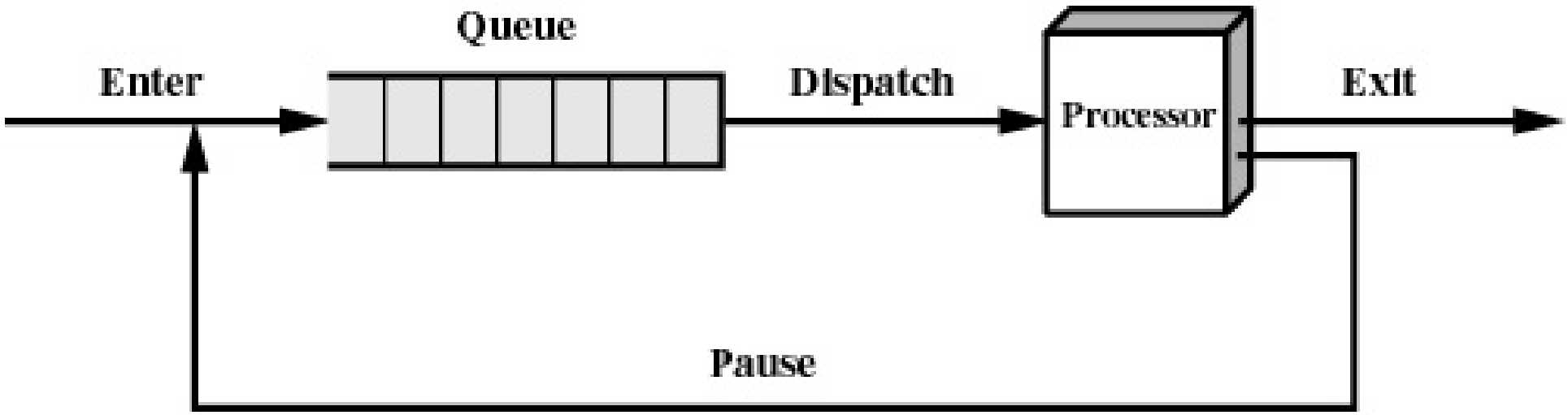
Running ▷Blocked

- Waiting for input
  - File, network,
- Waiting for a timer (alarm signal)
- Waiting for a resource to become available

Dispatcher

- Sometimes also called the scheduler
  - The literature is also a little inconsistent on this point
- Has to choose a *Ready* process to run
  - How?
  - It is inefficient to search through all processes

The Ready Queue

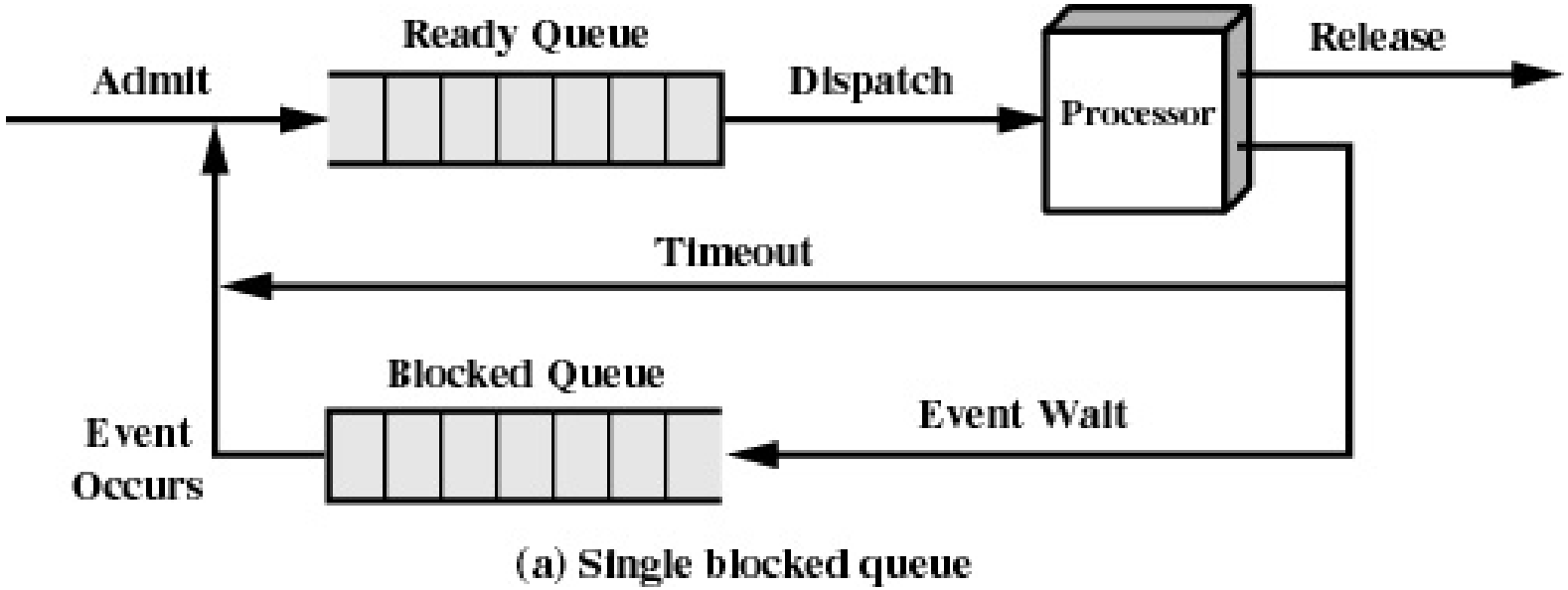


(b) Queuing diagram

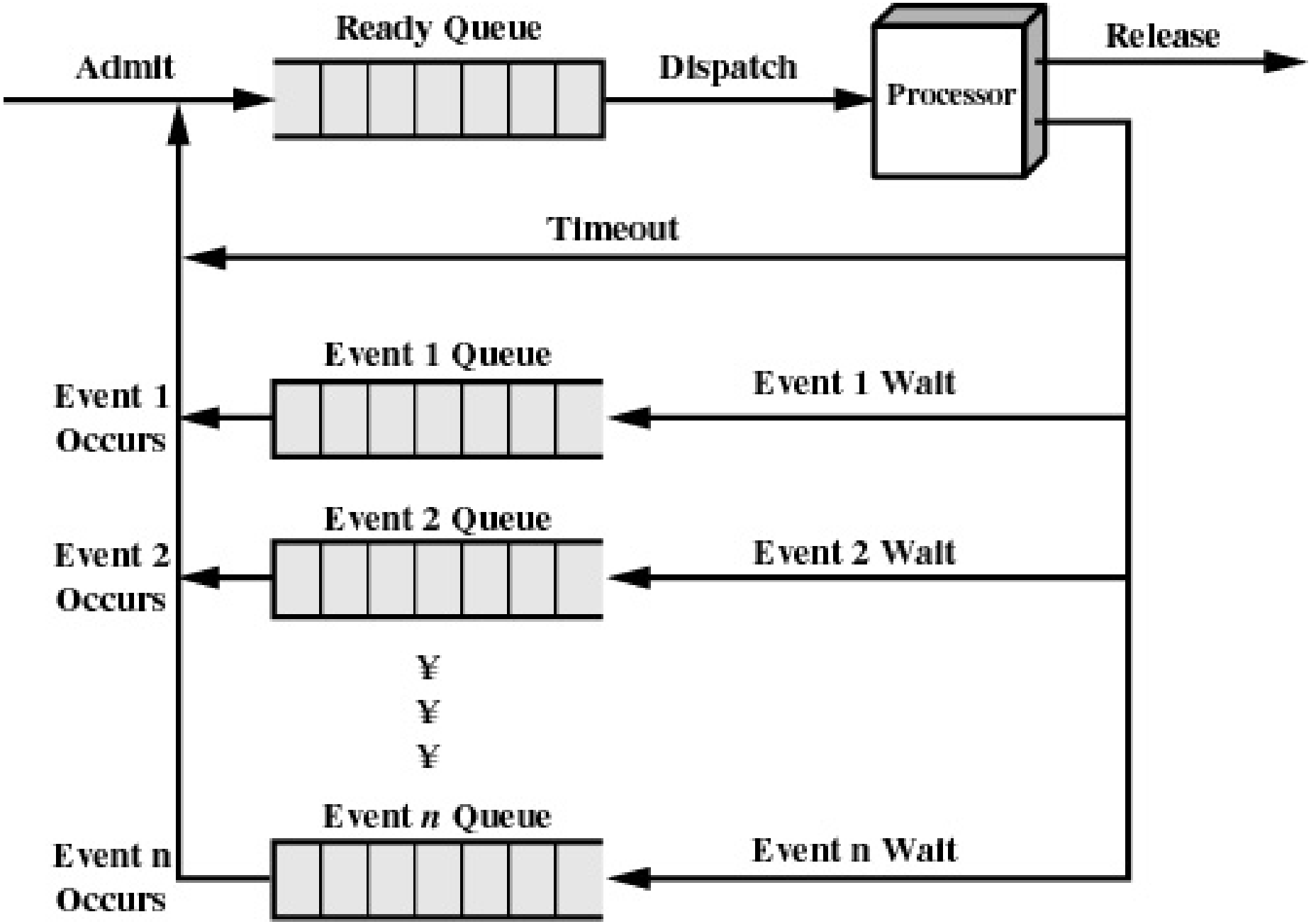
What about blocked processes?

- When an *unblocking* event occurs, we also wish to avoid scanning all processes to select one to make *Ready*

Using Two Queues



(a) Single blocked queue



(b) Multiple blocked queues