COMP 4161

NICTA Advanced Course

Advanced Topics in Software Verification

Toby Murray, June Andronick, Gerwin Klein

```
1:      public static int binarySearch(int[] a, int key) {
2:          int low = 0;
3:          int high = a.length - 1;
4:
5:          while (low <= high) {
6:              int mid = (low + high) / 2;
7:              int midVal = a[mid];
8:
9:              if (midVal < key)
10:                   low = mid + 1
11:              else if (midVal > key)
12:                   high = mid - 1;
13:              else
14:                   return mid; // key found
15:          }
16:         return -(low + 1);  // key not found.
17:     }
```

6:                int mid = (low + high) / 2;

http://googleresearch.blogspot.com/2006/06/
extra-extra-read-all-about-it-nearly.html

Organisatorials

When   Mon   14:00 – 15:30

            Thu   15:00 – 16:30

Where   Mon: Quadrangle G044 (E15-G044)

            Thu: Mathews 309 (F23-309)

`http://www.cse.unsw.edu.au/~cs4161/`

NICTA

Members of the seL4 verification team

➜ Functional correctness and security of a C microkernel

Security ↔ Isabelle/HOL model ↔ Haskell model ↔ C code

➜ 10 000 LOC / 500 000 lines of proof script (!)

➜ a bit under 30 person years of effort

It's all being open sourced, tomorrow!

`http://sel4.systems`

We are always embarking on exciting new projects.
We offer

➜ summer student scholarship projects

➜ honours and PhD theses

➜ research assistant and verification engineer positions

What you will learn

➔ how to use a theorem prover

➔ background, how it works

➔ how to prove and specify

➔ how to reason about programs

Health Warning

Theorem Proving is addictive

Content — Using Theorem Provers

Rough timeline

➜ Intro & motivation, getting started [today]

➜ Foundations & Principles

- Lambda Calculus, natural deduction [1,2]
- Higher Order Logic [3[a]]
- Term rewriting [4]

➜ Proof & Specification Techniques

- Inductively defined sets, rule induction [5]
- Datatypes, recursion, induction [6, 7]
- Hoare logic, proofs about programs, C verification [8[b],9]
- (mid-semester break)
- Writing Automated Proof Methods [10]
- Isar, codegen, typeclasses, locales [11[c],12]

[a]a1 due; [b]a2 due; [c]a3 due

What you should do to have a chance at succeeding

➜ attend lectures

➜ try Isabelle early

➜ redo all the demos alone

➜ try the exercises/homework we give, when we do give some

➜ **DO NOT CHEAT**

- Assignments and exams are take-home. This does NOT mean you can work in groups. Each submission is personal.
- For more info, see Plagiarism Policy[a]

---

[a] `www.cse.unsw.edu.au/about-us/organisational-structure/student-services/policies/`

Credits

This course was originally written by

Gerwin Klein

Credits

some material (in using-theorem-provers part) shamelessly stolen from

Tobias Nipkow, Larry Paulson, Markus Wenzel

David Basin, Burkhardt Wolff

Don't blame them, errors are ours

What is a proof?

NICTA

to prove   (Merriam-Webster)

➜ from Latin probare (test, approve, prove)

➜ to learn or find out by experience (archaic)

➜ to establish the existence, truth, or validity of
(by evidence or logic)
prove a theorem, the charges were never proved in court

pops up everywhere

➜ politics (weapons of mass destruction)

➜ courts (beyond reasonable doubt)

➜ religion (god exists)

➜ science (cold fusion works)

What is a mathematical proof?

In mathematics, a proof is a demonstration that, given certain axioms, some statement of interest is necessarily true. (Wikipedia)

Example: $\sqrt{2}$ is not rational.

Proof: assume there is $r \in \mathbb{Q}$ such that $r^2 = 2$.

Hence there are mutually prime p and q with $r = \frac{p}{q}$.

Thus $2q^2 = p^2$, i.e. p^2 is divisible by 2.

2 is prime, hence it also divides p, i.e. $p = 2s$.

Substituting this into $2q^2 = p^2$ and dividing by 2 gives $q^2 = 2s^2$. Hence, q is also divisible by 2. Contradiction. Qed.

Nice, but..

➜ still not rigorous enough for some

- what are the rules?
- what are the axioms?
- how big can the steps be?
- what is obvious or trivial?

➜ informal language, easy to get wrong

➜ easy to miss something, easy to cheat

Theorem. A cat has nine tails.

Proof. No cat has eight tails. Since one cat has one more tail than no cat, it must have nine tails.

A derivation in a formal calculus

Example: $A \wedge B \longrightarrow B \wedge A$ derivable in the following system

Rules:

$$\frac{X \in S}{S \vdash X} \text{ (assumption)} \qquad \frac{S \cup \{X\} \vdash Y}{S \vdash X \longrightarrow Y} \text{ (impI)}$$

$$\frac{S \vdash X \quad S \vdash Y}{S \vdash X \wedge Y} \text{ (conjI)} \qquad \frac{S \cup \{X, Y\} \vdash Z}{S \cup \{X \wedge Y\} \vdash Z} \text{ (conjE)}$$

Proof:

1. $\{A, B\} \vdash B$ (by assumption)

2. $\{A, B\} \vdash A$ (by assumption)

3. $\{A, B\} \vdash B \wedge A$ (by conjI with 1 and 2)

4. $\{A \wedge B\} \vdash B \wedge A$ (by conjE with 3)

5. $\{\} \vdash A \wedge B \longrightarrow B \wedge A$ (by impI with 4)

What is a theorem prover?

Implementation of a formal logic on a computer.

➜ fully automated (propositional logic)

➜ automated, but not necessarily terminating (first order logic)

➜ with automation, but mainly interactive (higher order logic)

➜ based on rules and axioms

➜ can deliver proofs

There are other (algorithmic) verification tools:

➜ model checking, static analysis, ...

➜ usually do not deliver proofs

➜ See COMP3153: Algorithmic Verification

Why theorem proving?

➜ Analysing systems/programs thoroughly

➜ Finding design and specification errors early

➜ High assurance (mathematical, machine checked proof)

➜ it's not always easy

➜ it's fun

Isabelle

➜ used here for applications, learning how to prove

What is Isabelle?

A generic interactive proof assistant

➜ **generic:**

not specialised to one particular logic

(two large developments: HOL and ZF, will mainly use HOL)

➜ **interactive:**

more than just yes/no, you can interactively guide the system

➜ **proof assistant:**

helps to explore, find, and maintain proofs

Why Isabelle?

- ➜ free
- ➜ widely used systems
- ➜ active development
- ➜ high expressiveness and automation
- ➜ reasonably easy to use
- ➜ (and because we know it best ;-))

If I prove it on the computer, it is correct, right?

If I prove it on the computer, it is correct, right?

No, because:

① hardware could be faulty

② operating system could be faulty

③ implementation runtime system could be faulty

④ compiler could be faulty

⑤ implementation could be faulty

⑥ logic could be inconsistent

⑦ theorem could mean something else

NICTA

No, but:

probability for

➜ OS and H/W issues reduced by using different systems

➜ runtime/compiler bugs reduced by using different compilers

➜ faulty implementation reduced by having the right prover architecture

➜ inconsistent logic reduced by implementing and analysing it

➜ wrong theorem reduced by expressive/intuitive logics

No guarantees, but assurance immensly higher than manual proof

If I prove it on the computer, it is correct, right?

Soundness architectures

careful implementation   PVS

LCF approach, small proof kernel   HOL4

                                                Isabelle

explicit proofs + proof checker   Coq

                                                Twelf

                                                Isabelle

                                                HOL4

Meta Logic

Meta language:

The language used to talk about another language.

Examples:

English in a Spanish class, English in an English class

Meta logic:

The logic used to formalize another logic

Example:

Mathematics used to formalize derivations in formal logic

Syntax:

Formulae: $F ::= V \mid F \longrightarrow F \mid F \wedge F \mid False$

$V ::= [A - Z]$

Derivable: $S \vdash X$ X a formula, S a set of formulae

logic / meta logic

$$\frac{X \in S}{S \vdash X}$$

$$\frac{S \cup \{X\} \vdash Y}{S \vdash X \longrightarrow Y}$$

$$\frac{S \vdash X \quad S \vdash Y}{S \vdash X \wedge Y}$$

$$\frac{S \cup \{X, Y\} \vdash Z}{S \cup \{X \wedge Y\} \vdash Z}$$

$$\bigwedge \qquad \Longrightarrow \qquad \lambda$$

Syntax: $\bigwedge x.\, F$ (F another meta level formula)

in ASCII: `!!x. F`

➜ universal quantifier on the meta level

➜ used to denote parameters

➜ example and more later

$$\Longrightarrow$$

Syntax: $A \Longrightarrow B$　　　　$(A, B$ other meta level formulae$)$

in ASCII: `A ==> B`

Binds to the right:

$$A \Longrightarrow B \Longrightarrow C \quad = \quad A \Longrightarrow (B \Longrightarrow C)$$

Abbreviation:

$$[\![A; B]\!] \Longrightarrow C \quad = \quad A \Longrightarrow B \Longrightarrow C$$

➜ read: A and B implies C

➜ used to write down rules, theorems, and proof states

NICTA

mathematics: if $x < 0$ and $y < 0$, then $x + y < 0$

formal logic: $\vdash\ x < 0 \wedge y < 0 \longrightarrow x + y < 0$

variation: $x < 0; y < 0 \ \vdash\ x + y < 0$

Isabelle: **lemma** "$x < 0 \wedge y < 0 \longrightarrow x + y < 0$"

variation: **lemma** "$[\![x < 0; y < 0]\!] \Longrightarrow x + y < 0$"

variation: **lemma**

assumes "$x < 0$" and "$y < 0$" shows "$x + y < 0$"

Example: a rule

logic:
$$\frac{X \quad Y}{X \wedge Y}$$

variation:
$$\frac{S \vdash X \quad S \vdash Y}{S \vdash X \wedge Y}$$

Isabelle:
$$[\![X; Y]\!] \Longrightarrow X \wedge Y$$

logic:

$$\frac{X \vee Y \quad \begin{matrix} X \\ \vdots \\ Z \end{matrix} \quad \begin{matrix} Y \\ \vdots \\ Z \end{matrix}}{Z}$$

variation:

$$\frac{S \cup \{X\} \vdash Z \quad S \cup \{Y\} \vdash Z}{S \cup \{X \vee Y\} \vdash Z}$$

Isabelle:

$$[\![X \vee Y; X \Longrightarrow Z; Y \Longrightarrow Z]\!] \Longrightarrow Z$$

λ

Syntax: $\lambda x.\, F$ (F another meta level formula)

in ASCII: `%x. F`

➜ lambda abstraction

➜ used for functions in object logics

➜ used to encode bound variables in object logics

➜ more about this in the next lecture

ENOUGH THEORY!

GETTING STARTED WITH ISABELLE

System Architecture

Prover IDE (jEdit) – user interface

HOL, ZF – object-logics

Isabelle – generic, interactive theorem prover

Standard ML – logic implemented as ADT

User can access all layers!

System Requirements

➜ **Linux**, **Windows**, or **MacOS X (10.7 +)**

➜ **Standard ML**

   (PolyML fastest, SML/NJ supports more platforms)

➜ **Java** (for jEdit)

Premade packages for Linux, Mac, and Windows + info on:

`http://mirror.cse.unsw.edu.au/pub/isabelle/`

Documentation

Available from `http://isabelle.in.tum.de`

➜ Learning Isabelle

- Tutorial on Isabelle/HOL (LNCS 2283)
- Tutorial on Isar
- Tutorial on Locales

➜ Reference Manuals

- Isabelle/Isar Reference Manual
- Isabelle Reference Manual
- Isabelle System Manual

➜ Reference Manuals for Object-Logics

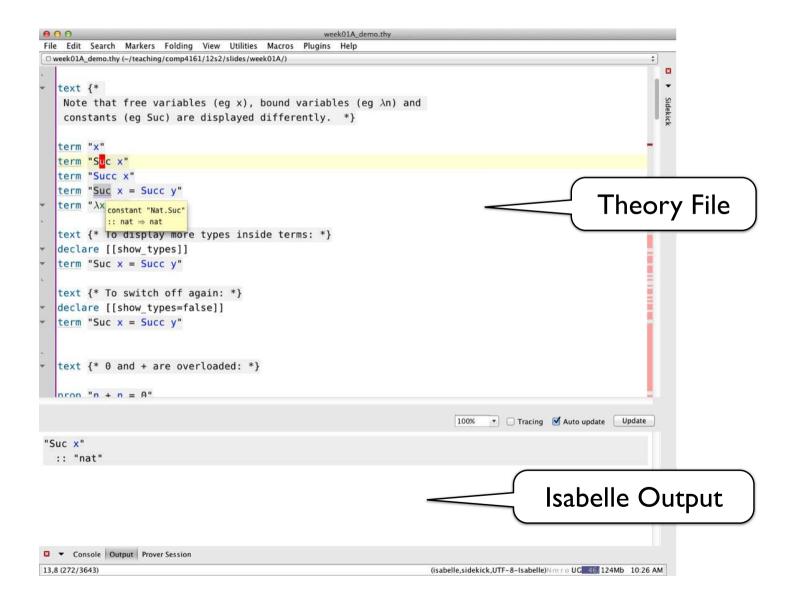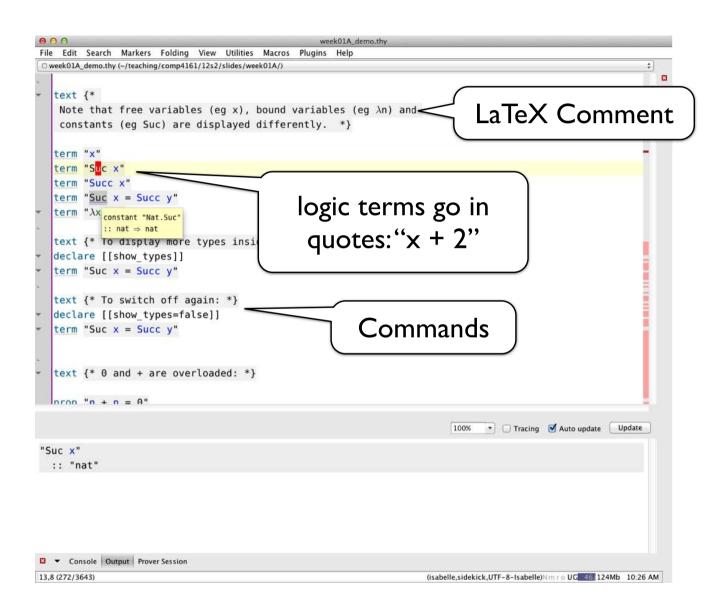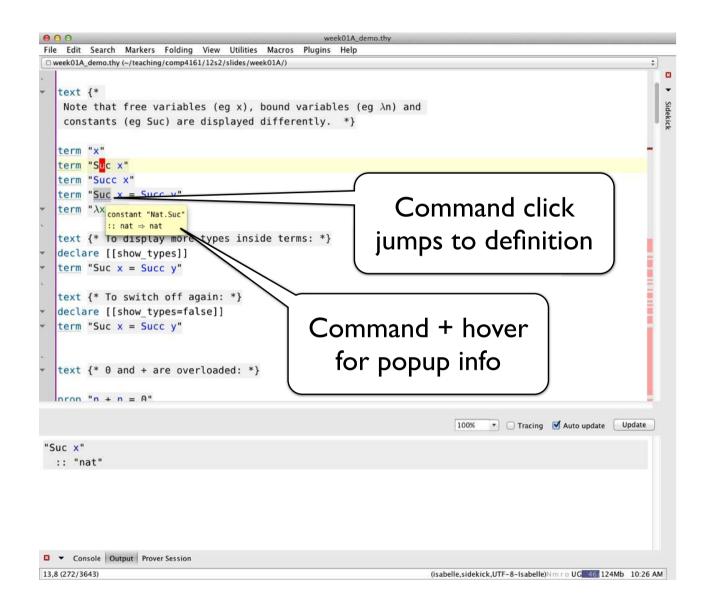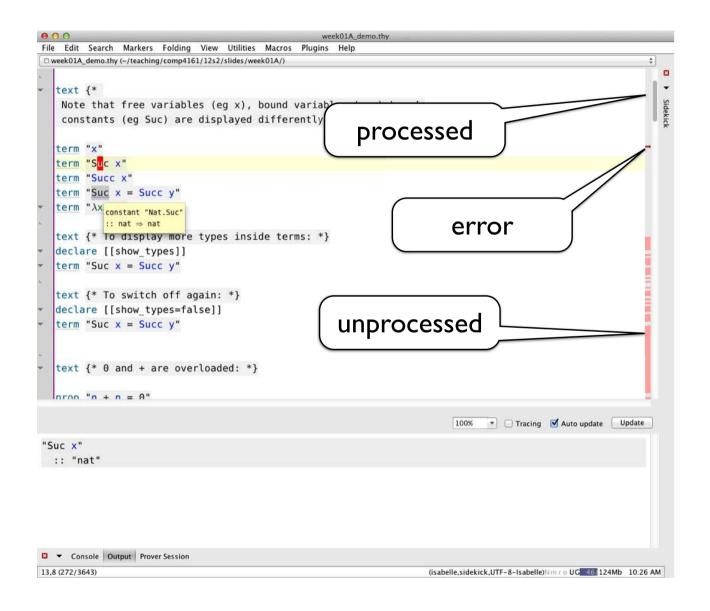# jEdit/PIDE

Theory File

Isabelle Output

jEdit/PIDE

jEdit/PIDE

jEdit/PIDE

DEMO

➜ Download and install Isabelle from

   `http://mirror.cse.unsw.edu.au/pub/isabelle/`

➜ Step through the demo files from the lecture web page

➜ Write your own theory file, look at some theorems in the library, try 'find_theorems'

➜ How many theorems can help you if you need to prove something like "Suc(Suc x))"?

➜ What is the name of the theorem for associativity of addition of natural numbers in the library?