COMP9242 Advanced OS

S2/2017 W08: **Operating System Security**

@GernotHeiser

Incorporating Material from Toby Murray

UNSW AUSTRALIA

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What is security?

Different things to different people:

On June 8, as the investigation into the initial intrusion proceeded, the response team shared with relevant agencies that there was a high degree of confidence that OPM systems containing information related to the background investigations of prospective Federal governm whom a Federal background may have been compromisec

Computer Security

- Protecting **my** interests that are under computer control from malign threats
- Inherently subjective
  - Different people have different interests
  - Different people face different threats
- Don't expect one-size-fits-all solutions
  - Grandma doesn't need an air gap
  - Windows alone is insufficient for protecting TOP SECRET (TS) classified data on an Internet-connected machine

Claiming system "*security*" only makes sense with respect to well-defined *security objectives*:

- Identify *threats*
- Identify set of *secure system states*

State of OS Security

- Traditionally:
  - Has not kept pace with evolving user demographics
    - Focused on e.g. Defence and Enterprise
  - Has not kept pace with evolving threats
    - Focused on protecting users from users, not apps they run
- Is getting better in some ways
  - Eg smartphone OSes implement stricter security than desktops
  - Less blatant holes in mainstream Oses
- Is getting worse in others:
  - OSes are still growing in size, functionality and complexity
  - Too few people understand how to write secure code
  - More people know how to attack

OS Security

- What is the role of the OS for security?
- Minimum:
  - provide **mechanisms** to allow the construction of secure systems
  - that are capable of securely implementing the intended users'/ administrators' **policies**
  - while ensuring these mechanisms cannot be subverted

Good security mechanisms

- Are widely applicable
- Support general security principles
- Are easy to use correctly and securely
- Do not hinder non-security priorities (e.g. productivity, generativity)
  - Principle of “do not pay for what you don’t need”
- Lend themselves to correct implementation and verification

Security Design Principles

- Saltzer & Schroeder [SOSP ’73, CACM ’74]
  - **Economy of mechanism** – KISS
  - **Fail-safe defaults** – as in good engineering
  - **Complete mediation** – check everything
  - **Open design** – not security by obscurity
  - **Separation of privilege** – defence in depth
  - **Least privilege** – aka *principle of least authority* (POLA)
  - **Least common mechanism** – minimise sharing
  - **Psychological acceptability** – if it’s hard to use it won’t be

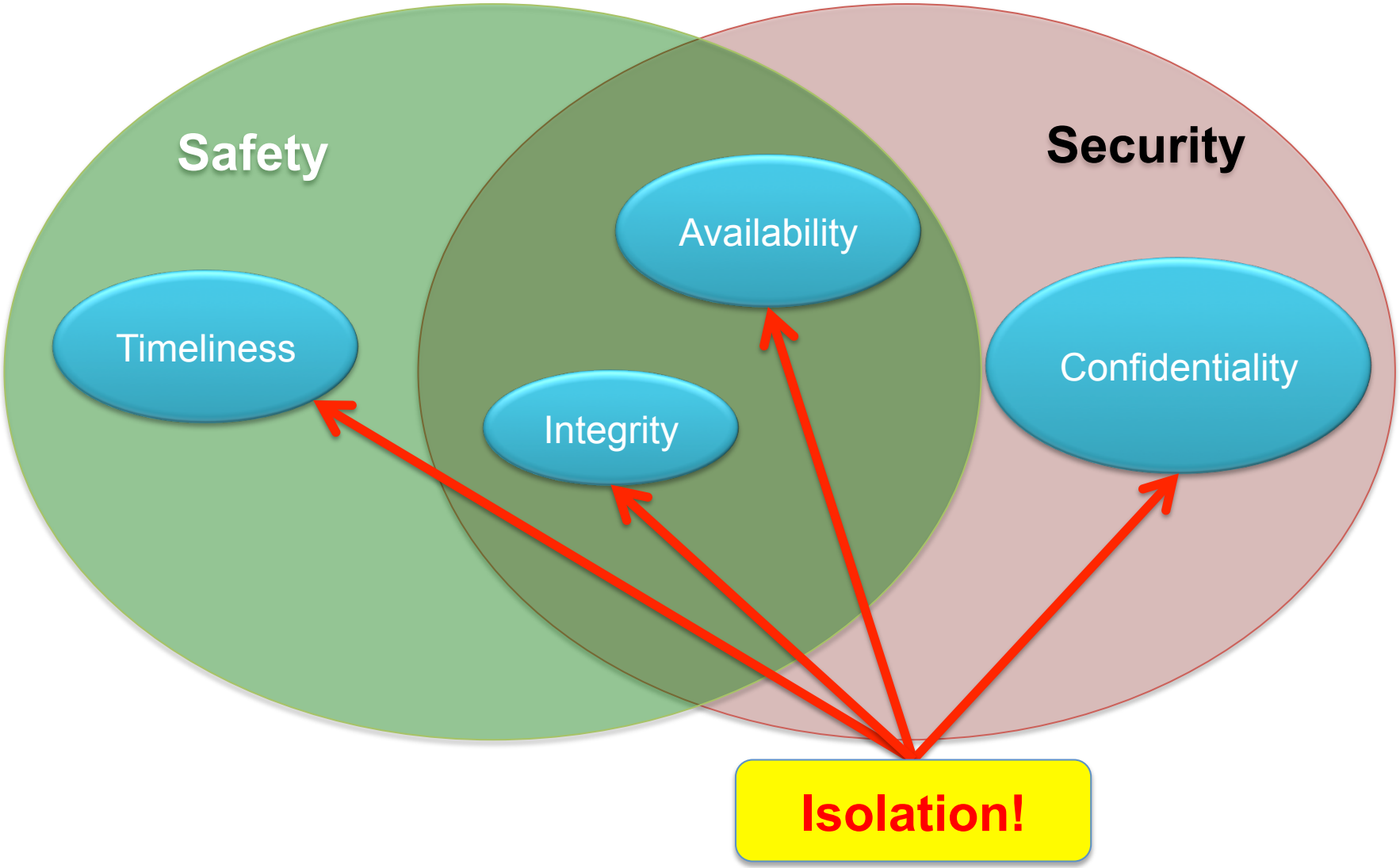
Common OS Security Mechanisms

- Access Control Systems
  - control what each process can access
- Authentication Systems
  - confirm the identity on whose behalf a process is running
- Logging
  - for audit, detection, forensics and recovery
- Filesystem Encryption
- Credential Management
- Automatic Updates

Security Policies

- Define what should be protected
  - and from whom
- Often in terms of common security goals (*CIA properties*):
  - **Confidentiality**
    - X should not be learnt by Y
  - **Integrity**
    - X should not be tampered with by Y
  - **Availability**
    - X should not be made unavailable to Z by Y

Security vs Safety



Policy vs. Mechanism

- Policies accompany mechanisms:
  - **access control** policy
    - who can access what?
  - **authentication** policy
    - is password sufficient to authenticate TS access?
- Policy often restricts the applicable mechanisms
- One person's policy is another's mechanism

Assumptions

- All policies and mechanisms operate under certain **assumptions**
  - **e.g.** TS-cleared users can be trusted not to write TS data into the UNCLASS window
- Problem: implicit or poorly understood assumptions
- Good assumptions:
  - clearly identified
  - verifiable

Risk Management

- Comes down to **risk management**
  - At the heart of all security
  - Assumptions: no absolute security, what risks we are willing to tolerate?
  - Cost & likelihood of violation vs. cost of prevention
  - Gain vs cost for attacker
- Other risks:
  - we mitigate (using security mechanisms)
  - or transfer (e.g. by buying insurance)
- Security policy should distinguish which is appropriate for each risk
  - Based on a thorough **risk assessment**

Trust

- Systems always have **trusted** entites
  - whose misbehaviour can cause insecurity
  - hardware, OS, sysadmin ...
- **Trusted Computing Base** (TCB):
  - the set of all such entities
- Secure systems require **trustworthy** TCBs
  - achieved through **assurance** and **verification**
  - shows that the TCB is unlikely to misbehave
  - *Minimising the TCB is key* for ensuring correct behaviour

Assurance and Formal Verification

- **Assurance**:
  - systematic evaluation and testing
  - essentially an intensive and onerous form of quality assurance
- **Formal verification**:
  - mathematical proof
- Together trying to establish correctness of:
  - the **design** of the mechanisms
  - and their **implementation**
- **Certification**: independent examination confirming that the assurance or verification was done right

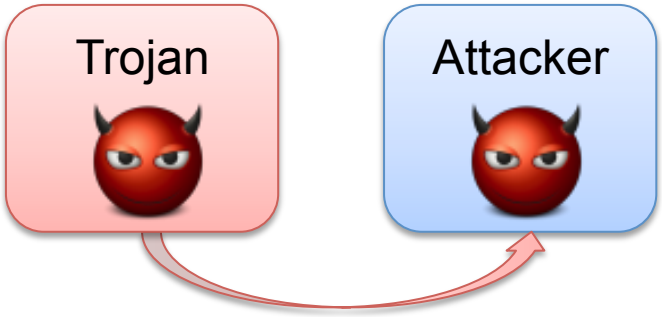
Covert Channels

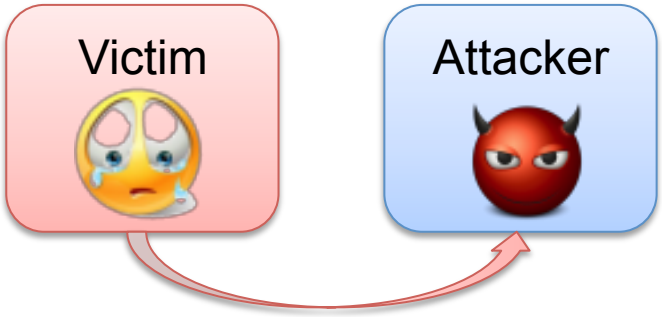
- Information flow not controlled by security mechanisms
  - Confidentiality requires absence of all such channels
- **Storage** Channel:
  - Attribute of shared resource used as channel
    - time stamp on file, existence of file
  - Controllable by access control
- **Timing** Channel:
  - Temporal order of shared resource accesses
  - Outside of access-control system
  - Much more difficult to control and analyse
- Other **physical** channels:
  - Power draw
  - Temperature (fan speed)
  - Electromagnetic emanation
  - Acoustic emanation

Covert Timing Channels

- Created by shared resource whose timing-related behaviour can be monitored
  - network bandwidth, CPU load, memory latency ...
- Requires access to a time source
  - anything that allows processes to synchronise
  - Generally compare relative occurrence of two event sequences (clocks)
- Critical issue is channel bandwidth
  - low bandwidth limits damage
    - why DRM ignores low bandwidth channels
  - beware of amplification
    - e.g. leaking passwords, encryption keys etc.

Covert Channels vs Side Channels





- Trojan intentionally creates signal through targeted resource use
- Worst-case bandwidth

- Attacker uses signal created by victim's innocent operations
- Much lower bandwidth

Summary: Introduction

- Security is very subjective, needs well-defined objectives
- OS security:
  - provide good security **mechanisms**
  - that support users' **policies**
- Security depends on establishing **trustworthiness** of trusted entities
  - **TCB:** set of all such entities
    - should be as small as possible
  - Main approaches: assurance and verification
- The OS is necessarily part of the TCB

ACCESS-CONTROL PRINCIPLES

Access Control

- **who** can access **what** in which ways
  - the “who” are called **subjects** (or **agents**)
    - e.g. users, processes etc.
  - the “what” are called **objects**
    - e.g. individual files, sockets, processes etc.
    - includes all subjects
  - the “ways” are called **permissions**
    - e.g. read, write, execute etc.
    - are usually specific to each kind of object
    - include those meta-permissions that allow modification of the protection state
      - e.g. own

AC Mechanisms and Policies

- AC Policy
  - Specifies allowed accesses
  - And how these can change over time
- AC Mechanism
  - Implements the policy
- Certain mechanisms lend themselves to certain kinds of policies
  - Some policies cannot be expressed using your OS's mechanisms

Protection State

Access control matrix [Lampson'71] defines the **protection state** at particular time

Subjects are also objects

	Obj1	**Obj2**	**Obj3**	**Subj2**
Subj1	R	RW		send
Subj2		RX		control
Subj3	RW		RWX own	recv

Storing Protection State

- Not usually as access control matrix
  - too sparse, inefficient, dynamic
- Two obvious choices:
  - store individual columns with each object
    - defines the subjects that can access each object
    - each such column is called the object's **access control list**
  - store individual rows with each subject
    - defines the objects each subject can access
      aka subject's **protection domain**
    - each such row is called the subject's **capability list**

Access Control Lists (ACLs)

- Subjects usually aggregated into classes
  - e.g. UNIX: owner, group, everyone
  - more general lists in Windows, recent Linux
  - Can have negative rights
    eg. to overwrite group rights
- Meta-permissions (e.g. own)
  - control class membership
  - allow modifying the ACL
- Implemented in almost all commercial OSes

Obj1

Subj1	R
Subj2	
Subj3	RW

Capabilities

- A **capability** [Dennis & Van Horn, 1966] is a capability list element

	Obj1	Obj2	Obj3	Subj2
Subj1	R	RW		send

  - **Names** an object to which the capability refers
  - **Confers** permissions over that object
- Capability is **prima facie authority** to perform an operation
  - System will perform operation iff appropriate capability is presented
- Less common in commercial systems
  - IBM System/38→AS/400→i-Series
  - KeyKOS (VISA transaction processing) [Bomberger et al, 1992]
  - More common in research: EROS [Shapiro'99], Cheri, seL4

Capability-Based Access Control

Cap = Access Token:
Prima-facie evidence of privilege

Obj reference

Access rights

Eg. read, write, send, execute…

Eg. thread, address space

Capabilities provide
- Fine-grained access control
- Reasoning about information flow

Any system call is invoking a capability:

err = method(cap, args);

Capabilities: Implementations

- Capabilities must be unforgeable
  - Traditionally protected by hardware (tagged memory), eg System-38
  - Can be copied etc like data
- On conventional hardware, either:
  - Stored as ordinary user-level data, but unguessable due to sparseness
    - contains password or secure hash: PCS [Anderson'86], Mungi
    - **"sparse" capabilies**
  - Stored separately (in-kernel), referred to by user programs by index/ address, eg Mach [Accetta'86], EROS
    - **"partitioned" or "segregated" capabilities**
    - like UNIX file descriptors
- Sparse capabilities can be leaked more easily
  - Huge amplification of covert channels!

ACLs and Capabilities: Duals?

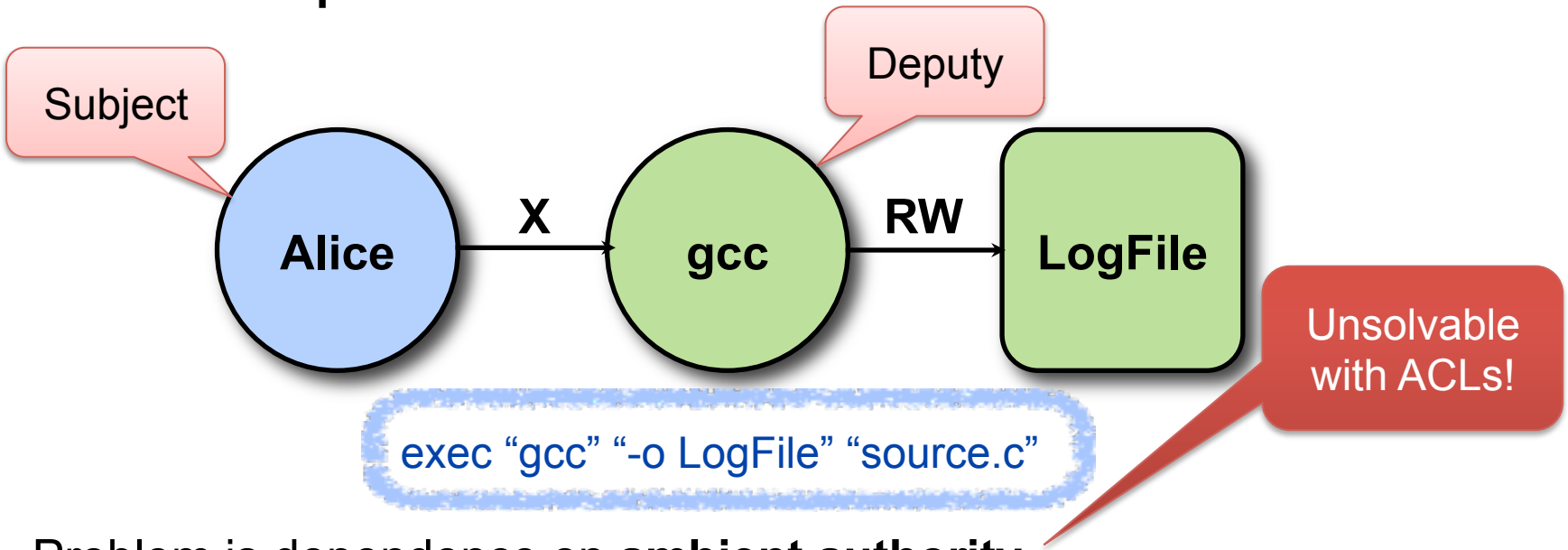
- In theory:
  - Dual representations of access control matrix
- Practical differences:
  - Naming and namespaces
    - Ambient authority
    - Deputies
  - Evolution of protection state
  - Forking
  - Auditing of protection state

Duals: Naming and Namespaces

- ACLs:
  - objects referenced by **name**
    - e.g. open(“/etc/passwd”,O_RDONLY)
  - require a subject (class) namespace
    - e.g. UNIX users and groups
- Capabilities:
  - objects referenced by **capability**
  - no further namespace required

Duals: Confused Deputies

- ACLs: separation of object naming and permission can lead to **confused deputies**



- Problem is dependence on **ambient authority**
  - Deputy uses its own authority when performing action on behalf of client
- Capabilities are both names and permissions
  - You can’t name something without having permission to it
  - Presentation is normally explicit (not ambient)

Duals: Evolution of Protection State

- ACLs:
  - Protection state changes by modifying ACLs
    - Requires certain meta-permissions on the ACL
- Capabilities:
  - Protection state changes by delegating and revoking capabilities
  - Fundamental properties enable reasoning about **information flow**:
    - A can send message to B only if A holds cap to B
    - A can obtain access to C only if it receives message with cap to C
  - **Right to delegate** may also be controlled by capabilities
    - e.g. A can delegate to B only if A has a capability to B that carries appropriate permissions
    - A can delegate X to B only if it has **grant** authority on X

Duals: Forking

- What permissions should children get?
- ACLs: depends on the child's subject
  - UNIX etc.: child inherits parent's subject
    - Inherits **all** of the parent's permissions
    - Any program you run inherits all of your authority
  - Bad for least privilege
- Capabilities: child has no caps by default
  - Parent gets a capability to the child upon fork
  - Used to delegate explicitly the necessary authority
  - **Defaults to least privilege**

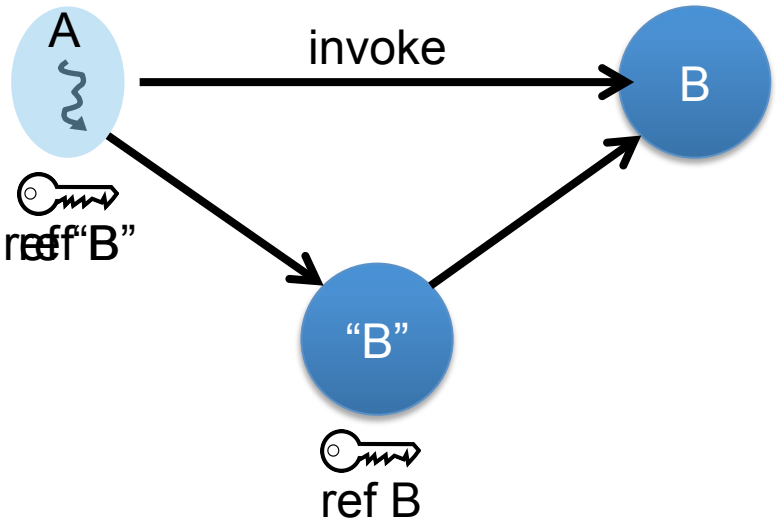
Duals: Auditing of Protection State

- Who has permission to access a particular object (right now)?
  - ACLs: Just look at the ACL
  - Caps: hard to determine with sparse or tagged caps, or for partitioned
- What objects a can particular subject access (right now)?
  - Capabilities: Just look at its capabilities
  - ACLs: may be impossible to determine without full scan
- “Who can access my stuff?” vs. “How much damage can X do?”

Interposing Object Access

Caps are opaque object references (pure names)

- Holder cannot tell which object a cap references nor the authority
- Supports transparent interposition (virtualisation)



Nice student project!

Usage:

- API virtualisation
- Security monitor
  - Security policy enforcement
  - Info flow tracing
  - Packet filtering…
- Secure logging
- Debugging
- Lazy object creation
  - Initial cap to constructor
  - Replace by proper object cap

Duals: Saltzer & Schroeder Principles

Security Principle	ACLs	Capabilities
Economy of Mechanism	Dubious	Yes!
Fail-safe defaults	Generally not	Yes!
Complete mediation	Yes (if properly done)	Yes (if properly done)
Open design	Neutral	Neutral
Separation of privilege	No	Doable
Least privilege	No	Yes
Least common mechanism	No	Yes
Psychological acceptability	Neutral	Neutral

Mandatory vs. Discretionary AC

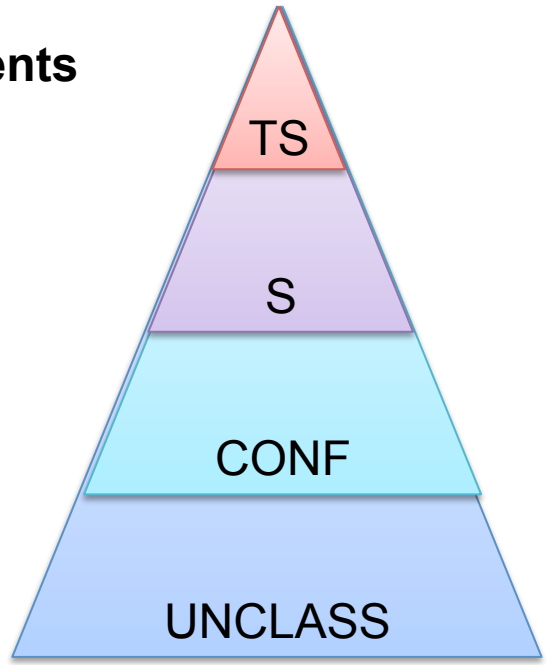
- Discretionary Access Control:
  - Users can make access control decisions
    - delegate their access to other users etc.
- Mandatory Access Control (MAC):
  - enforcement of administrator-defined policy
  - users cannot make access control decisions (except those allowed by mandatory policy)
  - can prevent untrusted applications running with user's privileges from causing damage

MAC

- Common in areas with global security requirements
  - e.g. national security classifications
- Less useful for general-purpose settings:
  - hard to support different kinds of policies
  - all policy changes must go through sysadmin
  - hard to dynamically delegate only specific rights required at runtime

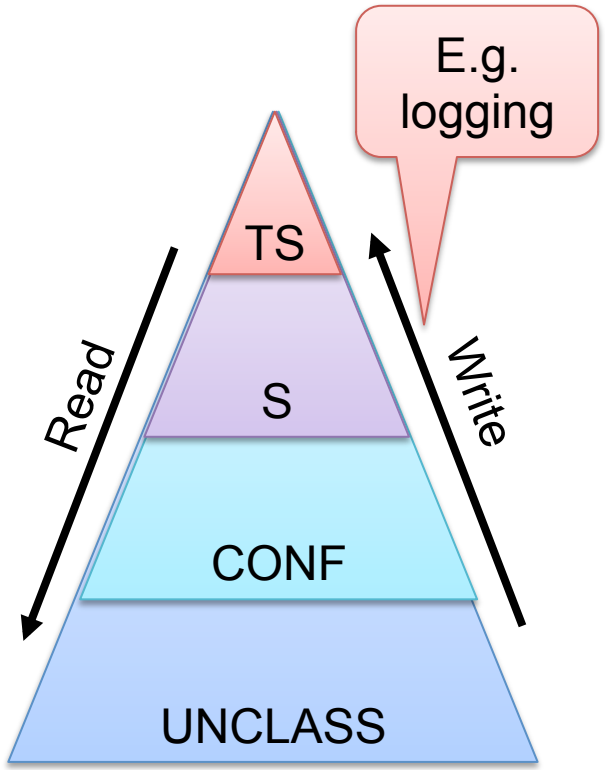
Bell-LaPadula [1966] (BLP) Model

- MAC Policy/Mechanism
  - Formalises National Security Classifications
- Every object assigned a **classification**
  - e.g. TS, S, C, U
  - may also have orthogonal security **compartments**
    - Support need-to-know
- Classifications ordered in a **lattice**
  - e.g. TS > S > C > U
- Every subject assigned a **clearance**
  - Highest classification they're allowed to learn



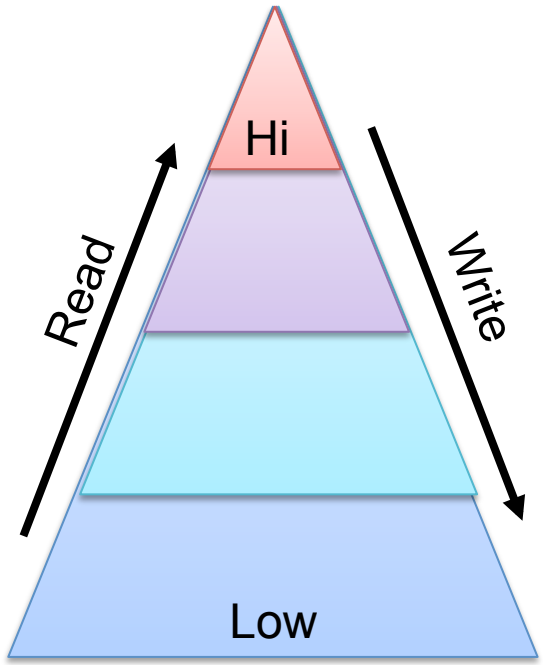
BLP: Rules

- **Simple Security Property** ("no read up"):
  - s can read o iff clearance(s) >= class(o)
  - S-cleared subject can read U,C,S but not TS
  - standard confidentiality
- ***-Property** ("no write down"):
  - s can write o iff clearance(s) <= class(o)
  - S-cleared subject can write TS,S, but not C,U
  - to prevent accidental or malicious leakage of data to lower levels



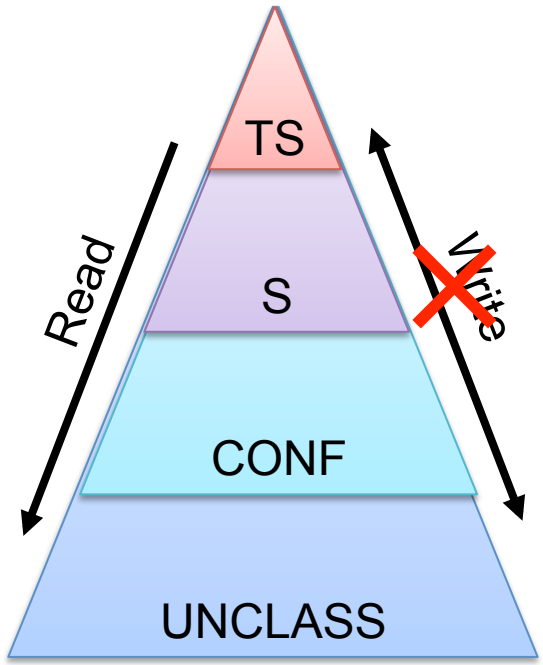
Biba Integrity Model

- Bell-LaPadula enforces **confidentiality**
- **Biba:** Its dual, enforces **integrity**
- Objects now carry **integrity** classification
- Subjects labelled by **lowest** level of data each subject is allowed to learn
- BLP order is inverted:
  - s can read o iff clearance(s) <= class(o)
  - s can write o iff clearance(s) >= class(o)



Confidentiality + Integrity

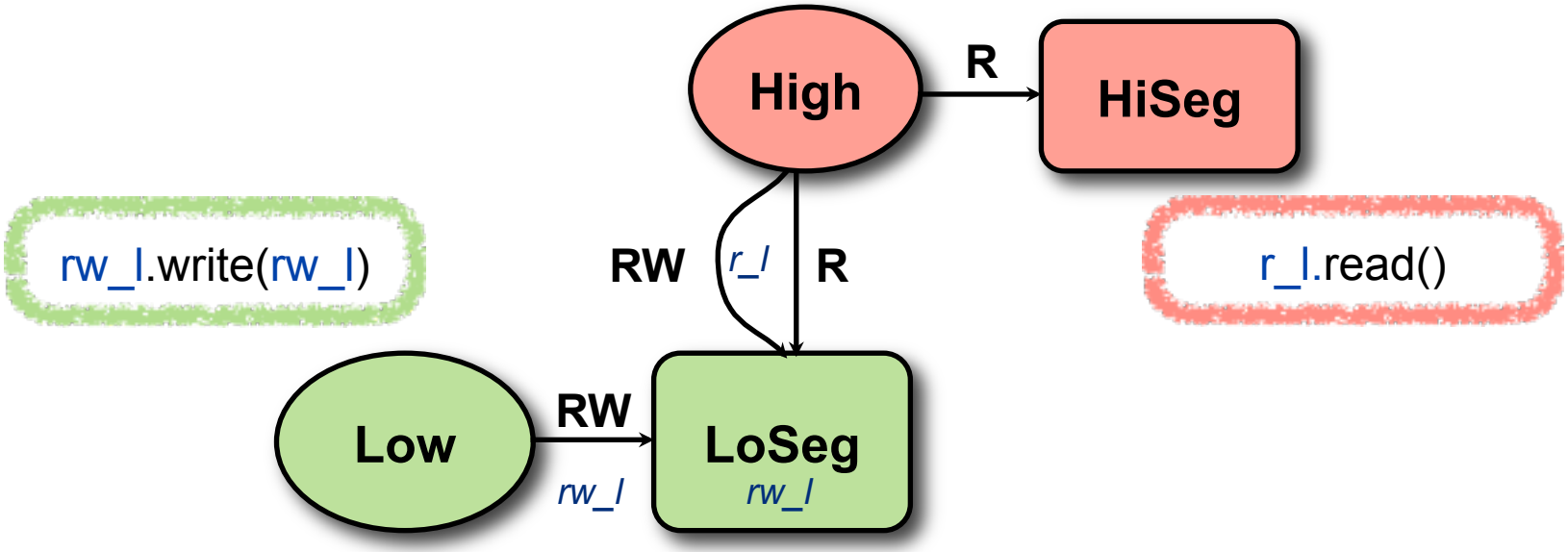
- BLP+Bibra allows no information flow across classes
  - Assume high-classified subject to treat low-integrity info responsibly
  - Allow read-down
- **Strong *-Property** ("matching writes only"):
  - s can write o iff clearance(s) = class(o)
  - Eg for logging, high reads low data and logs



Boebert's Attack

"On the inability of an unmodified capability cachine to enforce the *-property" [Boebert 1984]

- Shows an attack on capability systems that violates the *-property
  - Low passes cap to write buffer to High, which can then write down
  - Where caps and data are indistinguishable (sparse, tagged)
  - Does not work against **partitioned** capability systems



Boebert's Attack: Lessons

- Not all mechanisms can support all policies
- Many policies treat data- and access-propagation differently
  - Eg explicit **grant** capability (Take-grant model)
  - Cannot be expressed using sparse capability systems
- This does **not** mean that capability systems and MAC are incompatible in general

Decideability

- Boebert's attack highlights the need for **decideability** of safety in an AC system
- **Safety Problem:** given an initial protection state s, and a possible future protection state s', can s' be reached from s?
  - i.e. can an arbitrary (unwanted) access propagation occur?
- **Harrison, Ruzzo, Ullman [1975] (HRU):**
  - undecideable in general
  - equivalent to the halting problem

Decideable AC systems

- The safety problem for an AC system is **decideable** if we can always answer this question mechanically
- Most capability-based AC systems are decideable:
  - instances of Lipton-Snyder **Take-Grant** access control model [1977]
  - Take-Grant is decideable in linear time
- Less clear for many common ACL systems

Summary: AC Principles

- ACLs and Capabilities:
  - Capabilities tend to better support least privilege
  - But ACLs can be better for auditing
- MAC good for global security requirements
- Certain kinds of policies cannot be enforced with certain kinds of mechanisms
  - e.g. *-property with sparse capabilities
- AC systems should be decideable
  - so we can reason about them

ASSURANCE AND VERIFICATION

Assurance: Substantiating Trust

- Specification
  - unambiguous description of desired behaviour
- System design
  - justification that it meets specification
    - by mathematical proof or compelling argument
- Implementation
  - justification that it implements the design
    - by proof, code inspection, rigorous testing
- Maintenance
  - justifies that system use meets assumptions

Common Criteria

- Common Criteria for IT Security Evaluation [ISO/IEC 15408, 99]
  - ISO standard, for general use
  - evaluates QA used to ensure systems meet their requirements
  - Developed out of the famous US DOD "Orange Book": *Trusted Computer System Evaluation Criteria* [1985]
- **Target of Evaluation (TOE)** evaluated against **Security Target (ST)**
  - **ST:** statement of desired security properties based on **Protection Profiles**

Common Criteria: EALs

- 7 **Evaluated Assurance Levels**
  - **higher levels = more thorough evaluation**
    - **higher cost**
    - **not necessarily better security**

EAL 1–4 "not for use in hostile environments"

Level	Requirements	Specification	Design	Implementation
EAL1	not evaluated	**Informal**	not eval	not evaluated
EAL2	not evaluated	**Informal**	**Informal**	not evaluated
EAL3	not evaluated	**Informal**	**Informal**	not evaluated
EAL4	not evaluated	**Informal**	**Informal**	not evaluated
EAL5	not evaluated	**Semi-Formal**	**Semi-Formal**	**Informal**
EAL6	**Formal**	**Semi-Formal**	**Semi-Formal**	**Informal**
EAL7	**Formal**	**Formal**	**Formal**	**Informal**

Common Criteria Protection Profiles (PPs)

- Controlled Access PP (CAPP)
  - standard OS security, up to EAL3
- Single Level Operating System PP
  - superset of CAPP, up to EAL4+
- Labelled Security PP
  - MAC for COTS OSes
- Multi-Level Operating System PP
  - superset of CAPP, LSPP, up to EAL4+
- Separation Kernel Protection Profile (SKPP)
  - strict partitioning, for EAL6-7

COTS OS Certifications

- EAL3:
  - Mac OS X
- EAL4:
  - 2003: Windows 2000
  - 2005: SuSE Enterprise Linux
  - 2006: Solaris 10 (EAL4+)
    - against CAPP (an EAL3 PP!)
  - 2007: Red Hat Linux (EAL4+)
- EAL6
  - Green Hills INTEGRITY-178B (EAL6+)
    - against SKPP
    - relatively simple hardware platform in TOE

Get regularly hacked!

SKPP on Commodity Hardware

- SKPP:
  - OS provides only separation
- One Box One Wire (OB1) Project
  - Use INTEGRITY-178B to isolate VMs on commodity desktop hardware
  - Leverage existing INTEGRITY certification
    - by “porting” it to commodity platform
  - Conclusion [NSA, March 2010]:
    - SKPP validation for commodity hardware platforms infeasible due to their complexity
    - SKPP has limited relevance for these platforms
  - NSA subsequently dis-endorsed SKPP

Common Criteria Limitations

- Very expensive
  - rule of thumb: EAL6+ costs $1K/LOC
- Too much focus on development process
  - rather than the product that was delivered
- Lower EALs of little practical use for OSes
  - c.f. COTS OS EAL4 certifications
- Commercial Licensed Evaluation Facilities licenses rarely revoked
  - Leads to potential “race to the bottom” (Anderson & Fuloria, 2009)

Formal Verification

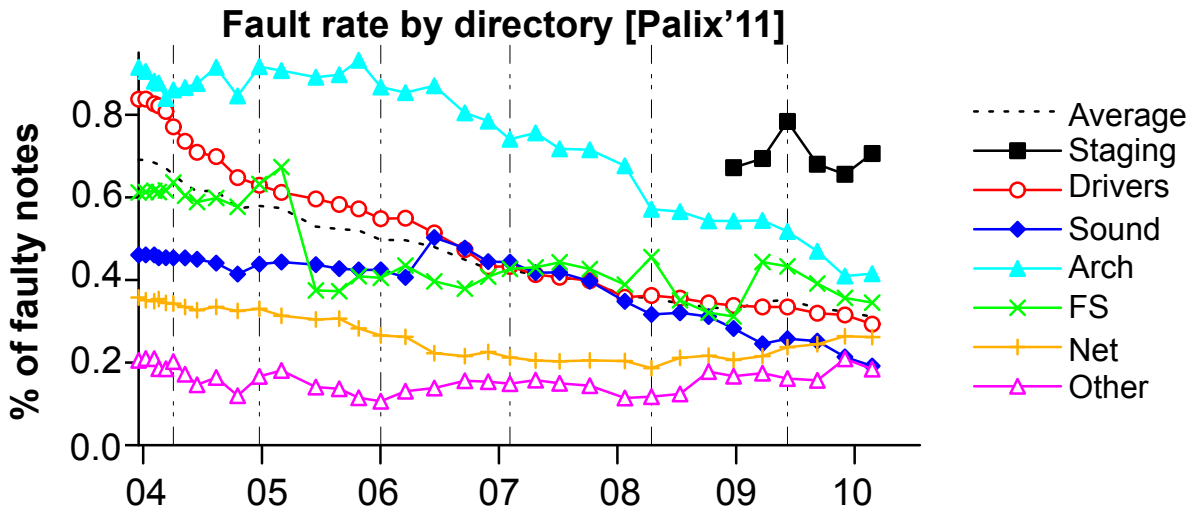
- Based on mathematical model of system
- Two approaches:
  - Automated techniques based on model checking / abstract interpretation
  - Theorem proving (manual or partially automated)

Automatic Analyses

- Algorithms that analyse code to detect certain kinds of defects
  - Usually static analysis
- Cannot generally “prove” code is correct
  - Only certain properties
  - False positives
  - False negatives
- Can be sound: guaranteed to detect all potential bugs of a kind
  - No false negatives
- Relatively cheap, often highly scalable (but then typically not sound)
  - Tradeoff between completeness and cost

Static Analysis and Linux: A Sad Story

- Static analysis of Linux source [Chou & al, 2001]
  - Found high density of bugs, especially in device drivers
- Re-analysis 10 years later [Palix & al, 2011]
  - Density of bugs detectable by static analysis had not dropped a lot!



And the Result?

ars TECHNICA BIZ & IT TECH SCIENCE POLICY CARS GAMING & CU

RISK ASSESSMENT —

Unsafe at any clock speed: Linux kernel security needs a rethink

Ars reports from the Linux Security Summit—a[illegible] find[illegible] wo[illegible] that needs to be done.

J.M. PORUP (UK) - 9/27/2016, 10:57 PM

BlueBorne

The Linux kernel today faces an unprecedented safety crisis. Much like when