

Incentives in Social Decision Schemes with Pairwise Comparison Preferences

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Abstract

Social decision schemes (SDSs) map the preferences of individual voters over multiple alternatives to a probability distribution over the alternatives. In order to study properties such as efficiency, strategyproofness, and participation for SDSs, preferences over alternatives are typically lifted to preferences over lotteries using the notion of stochastic dominance (*SD*). However, requiring strategyproofness or strict participation with respect to this preference extension only leaves room for rather undesirable SDSs such as random dictatorships. Hence, we focus on the natural but little understood pairwise comparison (*PC*) preference extension, which postulates that one lottery is preferred to another if the former is more likely to return a preferred outcome. In particular, we settle three open questions raised by Brandt [2017]: (i) there is no Condorcet-consistent SDS that satisfies *PC*-strategyproofness; (ii) there is no anonymous and neutral SDS that satisfies *PC*-efficiency and *PC*-strategyproofness; and (iii) there is no anonymous and neutral SDS that satisfies *PC*-efficiency and strict *PC*-participation. All three impossibilities require $m \geq 4$ alternatives and turn into possibilities when $m \leq 3$.

1 Introduction

Incentives constitute a central aspect when designing mechanisms for multiple agents: mechanisms should incentivize agents to participate and to act truthfully [Nisan *et al.*, 2007; Shoham and Leyton-Brown, 2009; Brandt *et al.*, 2016]. However, for many applications, guaranteeing these properties—usually called *participation* and *strategyproofness*—is a notoriously difficult task. This is particularly true for collective decision making, which studies the aggregation of preferences of multiple voters into a group decision, because strong impossibility theorems show that these axioms are in variance with other elementary properties [Gibbard, 1973; Satterthwaite, 1975; Moulin, 1988]. For instance, the Gibbard-Satterthwaite theorem shows that every strategyproof voting rule is either dictatorial or imposing, and Moulin’s No-Show paradox demonstrates that all Condorcet-consistent voting rules violate participation. A natural escape route in light of

these negative results is to allow for randomization in the output of the voting rule. Rather than returning a single winner, a *social decision scheme (SDS)* selects a lottery over the alternatives and the winner is eventually drawn at random according to the given probabilities.

In order to study properties such as efficiency, strategyproofness, and participation for SDSs, preferences over alternatives are typically lifted to preferences over lotteries using the notion of *stochastic dominance (SD)*, i.e., one lottery is preferred to another lottery if the expected utility of the former exceeds that of the latter for *every* utility representation consistent with the voter’s preferences over alternatives [Gibbard, 1977; Bogomolnaia and Moulin, 2001; Brandl *et al.*, 2018]. When demanding *SD*-efficiency, *SD*-strategyproofness, and anonymity, the only SDS that does the job is *uniform random dictatorship (RD)*, which selects a voter uniformly at random and then returns his favorite alternative [Gibbard, 1977]. Moreover, *RD* satisfies strict *SD*-participation, which means that voters are *strictly* better off participating (unless their top choice already receives probability 1). Unsurprisingly, *RD* has some severe shortcomings. It is often criticized for its inability to compromise: if all voters agree on a second best alternative but disagree on the best one, the uniform random dictatorship will not choose the common second best option. Furthermore, it cannot be extended to weak preferences without giving up *SD*-strategyproofness or *SD*-efficiency [Brandl *et al.*, 2016b]. On top of these criticisms, the representation of preferences over lotteries via expected utility functions has come under scrutiny [Allais, 1953; Kahneman and Tversky, 1979; Anand, 2009].

As an alternative to traditional expected utility representations, some authors have proposed to postulate that one lottery is preferred to another if the former is more likely to return a preferred outcome [Blyth, 1972; Packard, 1982; Blavatsky, 2006]. The resulting preference extension is known as *pairwise comparison (PC)* and represents a special case of Fishburn’s skew-symmetric bilinear utility functions [Fishburn, 1982]. Brandl *et al.* [2019] have shown that the No-Show paradox can be circumvented using *PC* preferences. Moreover, Brandl and Brandt [2020] proved that *PC* preferences constitute the *only* domain of preferences within a rather broad class of preferences over lotteries that allow for preference aggregation that satisfies independence of irrelevant alternatives and efficiency, thus avoiding Arrow’s impossibility.

In both cases, the resulting SDS is the set of *maximal lotteries* (*ML*), which was proposed by Fishburn [1984b] and has recently attracted significant attention [Brandl *et al.*, 2016a; Brandl *et al.*, 2022; Hoang, 2017].

It is known that *PC*-efficiency is stronger than *SD*-efficiency and violated by *RD*. *PC*-strategyproofness and strict *PC*-participation, on the other hand, are weaker than their *SD* counterparts. *ML* satisfies Condorcet-consistency, *PC*-efficiency, and *PC*-participation [Aziz *et al.*, 2018; Brandl *et al.*, 2019]. Furthermore, it is *PC*-strategyproof in preference profiles that admit a Condorcet winner [Hoang, 2017; Brandl *et al.*, 2022]. These encouraging results lead to the natural question of whether there are attractive SDSs that satisfy *PC*-strategyproofness or strict *PC*-participation. We address this question by proving the following theorems, all of which settle open problems raised by Brandt [2017]:

- There is no Condorcet-consistent SDS that satisfies *PC*-strategyproofness.
- There is no anonymous and neutral SDS that satisfies *PC*-efficiency and *PC*-strategyproofness.
- There is no anonymous and neutral SDS that satisfies *PC*-efficiency and strict *PC*-participation.

All three theorems hold for strict preferences and require $m \geq 4$ alternatives; we show that they turn into possibilities when $m \leq 3$ by constructing two new SDSs. The second theorem strengthens Theorem 5 by Aziz *et al.* [2018], which requires weak preferences. Our theorems demonstrate that efficiency, strategyproofness, and strict participation—which are satisfied by *RD* if we extend preferences using *SD*—are not compatible for *PC* preferences. Hence, there is no equivalent of random dictatorships for *PC* preferences. This also means that—unlike with Arrow’s impossibility and the No-Show paradox—*PC* preferences do not help to circumvent the Gibbard-Satterthwaite theorem. As a consequence, we face a tradeoff between efficiency and incentive-compatibility, which implies that *no* SDS can combine the advantages of *ML* and *RD* when using *PC* to compare lotteries. Hence, among the known SDSs, *RD* is the most attractive one when aiming for incentive-compatibility and *ML* when aiming for efficiency (and other properties such as Condorcet-consistency).

2 The Model

Let $A = \{a_1, \dots, a_m\}$ be a finite set of m alternatives and $\mathbb{N} = \{1, 2, 3, \dots\}$ an infinite set of voters. We denote by $\mathcal{F}(\mathbb{N})$ the set of all finite and non-empty subsets of \mathbb{N} . Intuitively, \mathbb{N} is the set of all potential voters, whereas $N \in \mathcal{F}(\mathbb{N})$ is a concrete electorate. Given an electorate $N \in \mathcal{F}(\mathbb{N})$, every voter $i \in N$ has a *preference relation* \succ_i , which is a complete, transitive, and anti-symmetric binary relation on A . In particular, we do not allow for ties (which only makes our results stronger). We write preference relations as comma-separated lists and denote the set of all preference relations by \mathcal{R} . A *preference profile* R on an electorate $N \in \mathcal{F}(\mathbb{N})$ contains a preference relation \succ_i for every voter $i \in N$, i.e., $R \in \mathcal{R}^N$. When writing preference profiles, we use sets before preference relations to indicate the voters who report the same preference relation. To this end, we define $[j \dots k] = \{i \in N : j \leq i \leq k\}$ and

note that $[j \dots k] = \emptyset$ if $j > k$. For instance, $[1 \dots 3]$: a, b, c means that voters 1, 2, and 3 prefer a to b to c . We omit the brackets for singleton sets. Given a preference profile $R \in \mathcal{R}^N$, the *majority margin* between two alternatives $x, y \in A$ is $g_R(x, y) = |\{i \in N : x \succ_i y\}| - |\{i \in N : y \succ_i x\}|$, i.e., the majority margin indicates how many more voters prefer x to y than *vice versa*. Furthermore, we define $n_R(x)$ as the number of voters who prefer alternative x the most in the profile R . Next, we denote by $R_{-i} = (\succ_1, \dots, \succ_{i-1}, \succ_{i+1}, \dots, \succ_n)$ the profile derived from $R \in \mathcal{R}^N$ by removing voter $i \in N$. Finally, $\mathcal{R}^{\mathcal{F}(\mathbb{N})}$ is the set of all possible preference profiles.

In this paper, we study *social decision schemes* (*SDSs*), which map preference profiles to lotteries over the alternatives. A *lottery* p is a probability distribution over the alternatives, i.e., a function $p : A \rightarrow [0, 1]$ such that $p(x) \geq 0$ for all $x \in A$ and $\sum_{x \in A} p(x) = 1$. The set of all lotteries on A is denoted by $\Delta(A)$. Then, an SDS f formally is a function of type $f : \mathcal{R}^{\mathcal{F}(\mathbb{N})} \rightarrow \Delta(A)$. We define $f(R, x)$ as the probability assigned to x by $f(R)$ and extend this notion to sets $X \subseteq A$ by $f(R, X) = \sum_{x \in X} f(R, x)$.

2.1 Fairness and Decisiveness

Next, we formalize desirable properties of SDSs. Two basic fairness notions are anonymity and neutrality, which require that voters and alternatives are treated equally, respectively. Formally, an SDS f is *anonymous* if $f(\pi(R)) = f(R)$ for all electorates $N \in \mathcal{F}(\mathbb{N})$, preference profiles $R \in \mathcal{R}^N$, and permutations $\pi : N \rightarrow N$. Here, $R' = \pi(R)$ is defined by $\succ'_i = \succ_{\pi(i)}$ for all $i \in N$. Analogously, *neutrality* requires of an SDS f that $f(\pi(R)) = \pi(f(R))$ for all electorates $N \in \mathcal{F}(\mathbb{N})$, preference profiles $R \in \mathcal{R}^N$, and permutations $\pi : A \rightarrow A$. This time, $R' = \pi(R)$ is the profile such that for all $i \in N$ and $x, y \in A$, $\pi(x) \succ'_i \pi(y)$ if and only if $x \succ_i y$. Another fairness condition is *cancellation*, which demands that the outcome does not change if two voters with inverse preferences join the electorate. Hence, an SDS f satisfies *cancellation* if $f(R) = f(R')$ for all preference profiles $R, R' \in \mathcal{R}^{\mathcal{F}(\mathbb{N})}$ such that R' is derived from R by adding two voters with inverse preferences.

A natural desideratum in randomized social choice is *decisiveness*: randomization should be avoided whenever possible. For instance, Condorcet-consistency formalizes this idea. We say an alternative x is a *Condorcet winner* in a profile R if $g_R(x, y) > 0$ for all $y \in A \setminus \{x\}$. Then, *Condorcet-consistency* requires of an SDS f that the Condorcet winner is chosen with probability 1 whenever it exists, i.e., $f(R, x) = 1$ for all preference profiles $R \in \mathcal{R}^{\mathcal{F}(\mathbb{N})}$ with Condorcet winner x . A weaker decisiveness condition is the *absolute winner* property. An *absolute winner* is an alternative x that is top-ranked by more than half of the voters in $R \in \mathcal{R}^N$, i.e., $n_R(x) > \frac{|N|}{2}$. Then, the *absolute winner property* requires that $f(R, x) = 1$ for all profiles $R \in \mathcal{R}^{\mathcal{F}(\mathbb{N})}$ with absolute winner x . Since absolute winners are also Condorcet winners, Condorcet-consistency implies the absolute winner property.

2.2 PC and SD Preferences

We assume that the voters’ preferences over alternatives are lifted to preferences over lotteries via the pairwise comparison

(*PC*) extension (see, e.g., [Aziz *et al.*, 2015; Aziz *et al.*, 2018; Brandt, 2017; Brandl and Brandt, 2020]). According to this notion, a voter prefers lottery p to lottery q if the probability that p returns a better outcome than q is at least as large as the probability that q returns a better outcome than p , i.e.,

$$p \succsim_i^{PC} q \iff \sum_{x,y \in A: x \succ_i y} p(x)q(y) \geq \sum_{x,y \in A: x \succ_i y} q(x)p(y).$$

The relation \succsim_i^{PC} is known to be complete but intransitive. An appealing interpretation of *PC* preferences is *ex ante* regret minimization, i.e., given two lotteries, a voter prefers the one which is less likely to result in *ex post* regret. Despite the simple definition, *PC* preferences are quite difficult to work with and even simple notions such as *PC*-efficiency are little understood [Aziz *et al.*, 2015].

Another well-known way to compare lotteries is *stochastic dominance* (*SD*) (e.g., [Gibbard, 1977; Brandl *et al.*, 2018]):

$$p \succsim_i^{SD} q \iff \forall x \in A: \sum_{y \in A: y \succ_i x} p(y) \geq \sum_{y \in A: y \succ_i x} q(y).$$

It follows from a result by Fishburn [1984a] that $p \succsim_i^{SD} q$ implies $p \succsim_i^{PC} q$ for all preference relations \succ_i and all lotteries p and q (see also [Aziz *et al.*, 2015]). In other words, the *SD* relation is a subrelation of the *PC* relation. For both $\mathcal{X} \in \{PC, SD\}$, we say a voter *strictly* \mathcal{X} -prefers p to q , denoted by $p \succ_i^{\mathcal{X}} q$, if $p \succsim_i^{\mathcal{X}} q$ and not $q \succsim_i^{\mathcal{X}} p$. Note that $p \succ_i^{SD} q$ implies $p \succ_i^{PC} q$.

2.3 Efficiency and Incentives

We can now define efficiency, strategyproofness, and participation. All of these axioms can be defined for both *SD* and *PC*; we thus define the concepts for $\mathcal{X} \in \{PC, SD\}$. First, we discuss efficiency, which requires that no lottery is unanimously preferred to the lottery chosen by the SDS. To formalize this, we say a lottery p \mathcal{X} -dominates another lottery q in a profile $R \in \mathcal{R}^N$ if $p \succsim_i^{\mathcal{X}} q$ for all voters $i \in N$ and $p \succ_{i^*}^{\mathcal{X}} q$ for some voter $i^* \in N$. Conversely, a lottery p is \mathcal{X} -efficient in R if it is not \mathcal{X} -dominated, and an SDS f is \mathcal{X} -efficient if $f(R)$ is \mathcal{X} -efficient for all preference profiles $R \in \mathcal{R}^{\mathcal{F}(N)}$. Both *PC*-efficiency and *SD*-efficiency imply *ex post* efficiency. For introducing this concept, we say an alternative x *Pareto-dominates* another alternative y in a profile $R \in \mathcal{R}^N$ if $x \succ_i y$ for all voters $i \in N$.¹ Then, *ex post efficiency* requires that $f(R, x) = 0$ for all profiles $R \in \mathcal{R}^{\mathcal{F}(N)}$ and alternatives $x \in A$ that are Pareto-dominated in R .

Next, we introduce strategyproofness, which demands that no voter can benefit by lying about his true preferences. Formally, an SDS f is \mathcal{X} -strategyproof if $f(R) \succsim_i^{\mathcal{X}} f(R')$ for all electorates $N \in \mathcal{F}(N)$, voters $i \in N$, and preference profiles $R, R' \in \mathcal{R}^N$ with $R_{-i} = R'_{-i}$.² Conversely, an SDS is \mathcal{X} -manipulable if it is not \mathcal{X} -strategyproof. Since strategyproofness does not require a variable electorate, we usually specify the electorates for which an SDS is strategyproof or manipulable. Similarly to strategyproofness, participation requires that voters should not be able to benefit

¹Recall that ties in \succ_i are not allowed.

²Another version is to require that $f(R') \not\succeq_i^{\mathcal{X}} f(R)$; these two versions coincide for *PC* because the *PC* extension is complete.

by abstaining from the election. Hence, an SDS f satisfies \mathcal{X} -participation if $f(R) \succsim_i^{\mathcal{X}} f(R_{-i})$ for all electorates $N \in \mathcal{F}(N)$, voters $i \in N$, and preference profiles $R \in \mathcal{R}^N$. In this paper, we are interested in *strict* \mathcal{X} -participation introduced by Brandl *et al.* [2015], which demands of an SDS f that, for all $N \in \mathcal{F}(N)$, $i \in N$, and $R \in \mathcal{R}^N$, it holds that $f(R) \succsim_i^{\mathcal{X}} f(R_{-i})$ and, moreover, $f(R) \succ_i^{\mathcal{X}} f(R_{-i})$ if there is a lottery p with $p \succ_i^{\mathcal{X}} f(R_{-i})$. That is, if possible, a voter strictly benefits from voting compared to abstaining.

Since $p \succ_i^{SD} q$ implies $p \succ_i^{PC} q$ and $p \succ_i^{SD} q$ implies $p \succsim_i^{PC} q$, the concepts of *SD*-efficiency, *SD*-strategyproofness, and *SD*-participation are related to the analogous concepts for *PC*: *PC*-efficiency entails *SD*-efficiency, whereas *SD*-strategyproofness and strict *SD*-participation are stronger than the corresponding notions for *PC* [Brandt, 2017]. See Figure 1 for an overview of these axioms.

2.4 Random Dictatorship and Maximal Lotteries

The following two important SDSs help to put our results into perspective: the uniform random dictatorship (*RD*) and maximal lotteries (*ML*). These SDSs are well-known and all subsequent claims are taken from the survey by Brandt [2017]. The *uniform random dictatorship* (*RD*) assigns probabilities proportional to $n_R(x)$, i.e., $RD(R, x) = \frac{n_R(x)}{\sum_{y \in A} n_R(y)}$ for every alternative $x \in A$ and preference profile $R \in \mathcal{R}^{\mathcal{F}(N)}$. *RD* is known to satisfy *SD*-strategyproofness, strict *SD*-participation, and *SD*-efficiency. Even more, when additionally imposing anonymity, it is the only SDS that satisfies these axioms. Since *SD*-strategyproofness and strict *SD*-participation imply the corresponding concepts for *PC*, *RD* satisfies our incentive axioms also if we extend preferences using *PC*. However, *RD* fails *PC*-efficiency and Condorcet-consistency, as can be seen in the following profile.

$$R: \quad 1: a, b, c \quad 2: b, a, c \quad 3: c, a, b$$

For this profile, $RD(R, x) = \frac{1}{3}$ for all $x \in A$, but a is the Condorcet winner and the lottery that puts probability 1 on a *PC*-dominates $RD(R)$.

In order to define *ML*, let $ML(R) = \{p \in \Delta(A) : \sum_{x,y \in A} p(x)q(y)g_R(x,y) \geq 0 \text{ for all } q \in \Delta(A)\}$ be the set of maximal lotteries for profile R . $ML(R)$ is non-empty by the minimax theorem and almost always a singleton. For all our claims about *ML*, it does not matter how ties are broken and any maximal lottery can be returned for a profile that admits multiple maximal lotteries. *ML* satisfies *PC*-efficiency, *PC*-participation, and Condorcet-consistency. However, *ML* fails *PC*-strategyproofness and strict *PC*-participation. The former can be seen by considering the following profiles.

$$\begin{aligned} R: & \quad \{1, 2\}: a, b, c \quad \{3, 4\}: b, c, a \quad 5: c, a, b \\ R': & \quad \{1, 2\}: a, b, c \quad 3: b, c, a \quad \{4, 5\}: c, a, b \end{aligned}$$

The unique maximal lotteries in R and R' , respectively, are p and q with $p(a) = q(c) = \frac{3}{5}$ and $p(b) = p(c) = q(a) = q(b) = \frac{1}{5}$. Since $\succ_i = \succ'_i$ for all $i \in \{1, 2, 3, 5\}$ and $q \succ_4^{PC} p$, voter 4 can *PC*-manipulate by deviating from R to R' . This raises the question of whether there is an SDS that unifies the advantages of *ML* and *RD*. As we show, this is not the case.

3 Results

We are now ready to present our results. The results for *PC*-strategyproofness are given in Section 3.1 while those for strict *PC*-participation are given in Section 3.2. Due to space restrictions, we defer most proofs to the full version of this paper [Brandt *et al.*, 2022] and discuss proof sketches instead.

3.1 *PC*-strategyproofness

In this section, we show that every Condorcet-consistent and every anonymous, neutral, and *PC*-efficient SDS is *PC*-manipulable when there are $m \geq 4$ alternatives. These results show that no SDS simultaneously satisfies *PC*-strategyproofness and the desirable properties of maximal lotteries. Moreover, since *PC*-strategyproofness is weaker than *SD*-strategyproofness, the incompatibility of *PC*-strategyproofness and Condorcet-consistency is a strengthening of the well-known impossibility of Condorcet-consistent and *SD*-strategyproof SDSs. Perhaps more surprising is the impossibility involving *PC*-efficiency: while anonymity, neutrality, *SD*-strategyproofness, and *SD*-efficiency characterize the uniform random dictatorship, the axioms become incompatible when moving from *SD* to *PC*. Since both impossibilities require $m \geq 4$ alternatives, we also show that they turn into possibilities if $m \leq 3$.

We start by discussing the impossibility of Condorcet-consistent and strategyproof SDSs.

Theorem 1. *Every Condorcet-consistent SDS is *PC*-manipulable if $|N| \geq 5$ is odd and $m \geq 4$.*

Proof. Assume for contradiction that there is a Condorcet-consistent and *PC*-strategyproof SDS f for $m \geq 4$ alternatives. Subsequently, we focus on the electorate $N = \{1, \dots, 5\}$ because we can generalize the result to any larger electorate with an odd number of voters by adding pairs of voters with inverse preferences. These voters do not change the Condorcet winner and hence will not affect our analysis.

As the first step, consider the profiles R^1 to R^4 . The * symbol is a placeholder for all missing alternatives.

R^1 :	1: $a, b, d, c, *$	2: $d, b, a, c, *$	3: $a, d, c, b, *$
	4: $*, c, d, b, a$	5: $c, b, a, d, *$	
R^2 :	1: $a, b, d, c, *$	2: $d, b, a, c, *$	3: $a, b, c, d, *$
	4: $*, c, d, b, a$	5: $c, b, a, d, *$	
R^3 :	1: $a, b, d, c, *$	2: $d, b, a, c, *$	3: $a, d, c, b, *$
	4: $*, c, d, a, b$	5: $c, b, a, d, *$	
R^4 :	1: $a, b, d, c, *$	2: $d, b, a, c, *$	3: $d, a, c, b, *$
	4: $*, c, d, b, a$	5: $c, b, a, d, *$	

Note that b is the Condorcet winner in R^2 , a in R^3 , and d in R^4 . Thus, Condorcet-consistency entails that $f(R^2, b) = f(R^3, a) = f(R^4, d) = 1$. In contrast, there is no Condorcet winner in R^1 and we use *PC*-strategyproofness to derive $f(R^1)$. For instance, this axiom postulates that $\sum_{x, y \in A: x \succ_{\frac{2}{3}} y} f(R^2, x) f(R^1, y) \geq \sum_{x, y \in A: x \succ_{\frac{2}{3}} y} f(R^1, x) f(R^2, y)$ as voter 3 can *PC*-manipulate by deviating from R^2 to R^1 otherwise. By

substituting $f(R^2, b) = 1$ and $f(R^2, x) = 0$ for $x \in A \setminus \{b\}$, we thus derive that

$$f(R^1, a) \leq f(R^1, A \setminus \{a, b\}). \quad (1)$$

Analogously, *PC*-strategyproofness between R^1 and R^3 and between R^1 and R^4 entails the following inequalities because voter 4 needs to *PC*-prefer $f(R^3)$ to $f(R^1)$ and voter 3 needs to *PC*-prefer $f(R^1)$ to $f(R^4)$.

$$f(R^1, A \setminus \{a, b\}) \leq f(R^1, b) \quad (2)$$

$$f(R^1, A \setminus \{a, d\}) \leq f(R^1, a) \quad (3)$$

Chaining the inequalities together, we get $f(R^1, A \setminus \{a, d\}) \leq f(R^1, a) \leq f(R^1, A \setminus \{a, b\}) \leq f(R^1, b)$, so $f(R^1, A \setminus \{a, b, d\}) = 0$. Simplifying (1), (2), and (3) then results in $f(R^1, a) \leq f(R^1, d) \leq f(R^1, b) \leq f(R^1, a)$, so $f(R^1, a) = f(R^1, b) = f(R^1, d) = \frac{1}{3}$.

Next, we analyze the profiles R^5 to R^8 .

R^5 :	1: $a, b, d, c, *$	2: $b, d, a, c, *$	3: $a, d, c, b, *$
	4: $*, c, d, b, a$	5: $c, b, a, d, *$	
R^6 :	1: $a, b, d, c, *$	2: $b, d, a, c, *$	3: $a, d, c, b, *$
	4: $*, c, d, b, a$	5: $b, c, a, d, *$	
R^7 :	1: $a, b, d, c, *$	2: $b, d, a, c, *$	3: $a, d, c, b, *$
	4: $*, c, d, a, b$	5: $c, b, a, d, *$	
R^8 :	1: $a, b, d, c, *$	2: $b, c, d, a, *$	3: $a, d, c, b, *$
	4: $*, c, d, b, a$	5: $c, b, a, d, *$	

Just as for the profiles R^1 to R^4 , there is no Condorcet winner in R^5 , whereas b is the Condorcet winner in R^6 , a in R^7 , and c in R^8 . Consequently, Condorcet-consistency requires that $f(R^6, b) = f(R^7, a) = f(R^8, c) = 1$. Next, we use *PC*-strategyproofness to derive $f(R^5)$. In particular, we infer the following inequalities as voter 5 needs to *PC*-prefer $f(R^5)$ to $f(R^6)$, voter 4 needs to *PC*-prefer $f(R^7)$ to $f(R^5)$, and voter 2 needs to *PC*-prefer $f(R^8)$ to $f(R^5)$.

$$f(R^5, A \setminus \{b, c\}) \leq f(R^5, c) \quad (4)$$

$$f(R^5, A \setminus \{a, b\}) \leq f(R^5, b) \quad (5)$$

$$f(R^5, b) \leq f(R^5, A \setminus \{b, c\}) \quad (6)$$

Analogous computations as for R^1 now show that $f(R^5, a) = f(R^5, b) = f(R^5, c) = \frac{1}{3}$. Finally, note that R^1 and R^5 only differ in the preferences of voter 2. This means that voter 2 can *PC*-manipulate by deviating from R^5 to R^1 since he *PC*-prefers $f(R^1)$ to $f(R^5)$. Hence, f fails *PC*-strategyproofness, which contradicts our assumptions. \square

It is open whether Theorem 1 also holds if $|N|$ is even.

Next, we turn the focus to our second impossibility result: every anonymous and neutral SDS that satisfies *PC*-efficiency is *PC*-manipulable. For proving this theorem, we first show that every SDS that satisfies the absolute winner property and *PC*-efficiency is *PC*-manipulable.

Lemma 1. *Every *PC*-efficient SDS that satisfies the absolute winner property is *PC*-manipulable if $|N| \geq 3$, $|N| \notin \{4, 6\}$, and $m \geq 4$.*

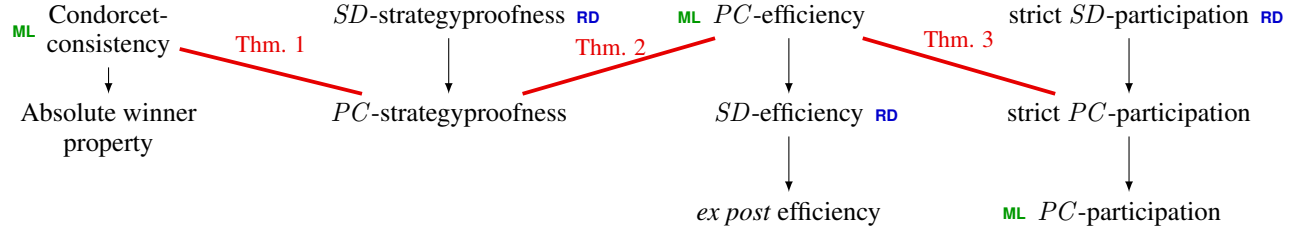


Figure 1: Overview of results. An arrow from an axiom X to another axiom Y indicates that X implies Y . The red lines between axioms represent impossibility theorems. Note that Theorems 2 and 3 additionally require anonymity and neutrality. Axioms labeled with ML are satisfied by maximal lotteries, and axioms labelled with RD are satisfied by the uniform random dictatorship.

Proof sketch. Assume for contradiction that there are an SDS f and an electorate $N \in \mathcal{F}(\mathbb{N})$ that satisfy the requirements of the lemma. In this sketch, we focus on $m = 4$ alternatives as we can extend the construction to more alternatives by adding Pareto-dominated ones. Moreover, we suppose that $n = |N|$ is odd; the argument for even n is similar but more involved. Now, consider the profiles R and R' shown below.

$$\begin{array}{l}
 R: \quad [1 \dots \frac{n-1}{2}]: a, d, b, c \quad \frac{n+1}{2}: b, c, d, a \\
 \quad \quad [\frac{n+3}{2} \dots n]: c, a, d, b \\
 R': \quad [1 \dots \frac{n-1}{2}]: a, d, b, c \quad \frac{n+1}{2}: b, d, c, a \\
 \quad \quad [\frac{n+3}{2} \dots n]: c, a, d, b
 \end{array}$$

The goal is to show that $f(R, a) = f(R, b) = f(R, c) = \frac{1}{3}$ and $f(R', a) = f(R', c) = f(R', d) = \frac{1}{3}$. Then, voter $\frac{n+1}{2}$ can PC -manipulate by deviating from R' to R because he PC -prefers $f(R)$ to $f(R')$. We proceed in three steps to derive the lottery $f(R)$: first, we use the absolute winner property and PC -strategyproofness to show that $f(R, c) > 0$. Next, we infer from PC -efficiency that $f(R, d) = 0$. Finally, we repeatedly apply PC -strategyproofness and the absolute winner property to prove that $f(R, a) \geq f(R, c) \geq f(R, b) \geq f(R, a)$, which implies that $f(R, x) = \frac{1}{3}$ for $x \in \{a, b, c\}$. An analogous argument can be used to derive $f(R')$. \square

Note that Lemma 1 is a rather strong impossibility itself. Next, we use it to prove that every anonymous, neutral, and PC -efficient SDS is PC -manipulable.

Theorem 2. *Every anonymous and neutral SDS that satisfies PC -efficiency is PC -manipulable if $|N| \geq 3$, $|N| \notin \{4, 6\}$, and $m \geq 4$.*

Proof sketch. We prove this theorem by showing that the given axioms imply the absolute winner property; then, Lemma 1 implies the impossibility. Consider an arbitrary SDS f that satisfies all given axioms and an electorate $N \in \mathcal{F}(\mathbb{N})$ with $n = |N| \geq 3$. Moreover, we focus on three alternatives because, as in Lemma 1, we can extend the argument to more alternatives by assuming that these are Pareto-dominated. We proceed with a case distinction with respect to the parity of n and, in this sketch, restrict attention to even n . Consider the following profile R .

$$\begin{array}{l}
 R: \quad 1: a, b, c \quad 2: a, c, b \\
 \quad \quad [3 \dots \frac{n}{2} + 1]: b, a, c \quad [\frac{n}{2} + 2 \dots n]: c, a, b
 \end{array}$$

Anonymity and neutrality require that $f(R, b) = f(R, c)$. Next, PC -efficiency implies that $f(R, b) = f(R, c) = 0$ and

hence $f(R, a) = 1$. Based on this insight, one can prove that a is chosen with probability 1 whenever the voters $i \in [1 \dots \frac{n}{2} + 1]$ report it as their best alternative. Due to anonymity and neutrality, this statement is equivalent to the absolute winner property. \square

Since both Theorems 1 and 2 require $m \geq 4$ alternatives, we can still hope for a possibility if $m \leq 3$. Indeed, for $m = 2$, ML satisfies Condorcet-consistency, PC -efficiency, PC -strategyproofness, anonymity, and neutrality. However, as shown in Section 2.4, ML fails PC -strategyproofness if $m = 3$. Thus, we construct another SDS that satisfies all given axioms. To this end, let $CW(R)$ be the set of Condorcet winners in R , and let $WCW(R) = \{x \in A: g_R(x, y) \geq 0 \text{ for all } y \in A \setminus \{x\}\}$ be the set of weak Condorcet winners if $CW(R) = \emptyset$, and $WCW(R) = \emptyset$ otherwise. Then, define the SDS f^1 as follows.

$$f^1(R) = \begin{cases} [x : 1] & \text{if } CW(R) = \{x\} \\ [x : \frac{1}{2}; y : \frac{1}{2}] & \text{if } WCW(R) = \{x, y\} \\ [x : \frac{3}{5}; y : \frac{1}{5}; z : \frac{1}{5}] & \text{if } WCW(R) = \{x\} \\ [x : \frac{1}{3}; y : \frac{1}{3}; z : \frac{1}{3}] & \text{otherwise} \end{cases}$$

It is easy to see that f^1 is Condorcet-consistent. The next proposition characterizes f^1 as the *only* SDS that satisfies cancellation and the axioms of Theorem 2.

Proposition 1. *For $m = 3$, f^1 is the only SDS that satisfies PC -efficiency, PC -strategyproofness, neutrality, anonymity, and cancellation.*

Proof sketch. The definition of f^1 immediately implies that this SDS satisfies anonymity, neutrality, and cancellation. Moreover, tedious case distinctions establish that f^1 satisfies PC -strategyproofness and PC -efficiency. For the reverse direction, we show first that every SDS that satisfies all given axioms is Condorcet-consistent. Building on this insight, we use the given axioms to infer the outcomes for all profiles. \square

Remark 1. Of all the axioms in this section besides cancellation, ML only fails PC -strategyproofness, dictatorships only fail anonymity and Condorcet-consistency, and the uniform random dictatorship only fails PC -efficiency and Condorcet-consistency. This shows that all axioms of Theorem 2 but neutrality are required for the result, and that both axioms are required for Theorem 1. We conjecture that Theorem 2 holds even without neutrality. Proposition 1 shows that $m \geq 4$ alternatives are required for both Theorems 1 and 2.

3.2 Strict PC -participation

In this section, we investigate strict PC -participation and prove that this axiom is incompatible with PC -efficiency. This result is rather surprising given that multiple SDSs are known to satisfy SD -efficiency and strict SD -participation [Brandl *et al.*, 2015]. Moreover, our impossibility can be seen as a complement to the work of Brandl *et al.* [2019] which shows that ML satisfies both PC -participation and PC -efficiency. In particular, our result demonstrates that maximal lotteries satisfy a maximal degree of participation subject to PC -efficiency. Finally, since Theorem 3 requires $m \geq 4$ alternatives, we construct an SDS that satisfies all our requirements when $m \leq 3$.

We first discuss the impossibility theorem.

Theorem 3. *No neutral and anonymous SDS satisfies both PC -efficiency and strict PC -participation if $m \geq 4$.*

Proof. Assume for contradiction that there is a neutral and anonymous SDS f that satisfies both PC -efficiency and strict PC -participation. In what follows, we focus on the case $m = 4$ because we can generalize our construction to $m > 4$ by adding $m - 4$ alternatives at the bottom of all voters' preference rankings. Since these dummy alternatives are Pareto-dominated by the original alternatives, PC -efficiency requires these alternatives to be assigned probability 0, and thus they will not affect our subsequent analysis.

First, consider the following profile with ten voters.

R^1 :	1: a, b, c, d	2: a, b, d, c	3: a, c, b, d
	4: a, c, d, b	5: a, d, b, c	6: a, d, c, b
	7: b, a, c, d	8: b, a, d, c	
	9: c, a, b, d	10: c, a, d, b	

Observe that d is Pareto-dominated by a , so by PC -efficiency, $f(R^1, d) = 0$. Moreover, since b and c are symmetric in this profile, neutrality and anonymity imply that $f(R^1, b) = f(R^1, c)$. If $f(R^1, b) = f(R^1, c) > 0$, then f is not PC -efficient because all voters weakly prefer the degenerate lottery that puts probability 1 on a , with voters 1–6 strictly preferring this lottery. Hence, $f(R^1, b) = f(R^1, c) = 0$ which means that $f(R^1, a) = 1$.

Next, consider profile R^2 , which is obtained by adding voter 11 with the preference d, a, b, c to R^1 . We infer from strict PC -participation that $f(R^2, d) > f(R^2, b) + f(R^2, c)$.

Finally, consider profile R^3 , which is obtained by adding voter 12 with the preference d, a, c, b to R^2 . Observe that b, c , and d are symmetric in R^3 , so by neutrality and anonymity, $f(R^3, b) = f(R^3, c) = f(R^3, d)$. If $f(R^3, b) = f(R^3, c) = f(R^3, d) > 0$, then f is not PC -efficient because all voters strictly prefer the degenerate lottery that puts probability 1 on a . Hence, $f(R^3, b) = f(R^3, c) = f(R^3, d) = 0$, which means that $f(R^3, a) = 1$. Since $f(R^2, d) > f(R^2, b) + f(R^2, c)$, voter 12 has a disincentive to participate in R^3 , thereby contradicting the strict PC -participation of f . \square

Since Theorem 3 requires $m \geq 4$, a natural question is whether the impossibility also holds for $m \leq 3$. As we demonstrate, the impossibility ceases to hold. If $m = 2$, it is easy to see that the uniform random dictatorship satisfies all axioms of Theorem 3. For $m = 3$, however, the uniform random dictatorship fails PC -efficiency (see Section 2.4). In light of

this, we construct a new SDS that satisfies all axioms used in Theorem 3. To this end, let B denote the set of alternatives that are never bottom-ranked. Then, the SDS f^2 is defined as follows: return the uniform random dictatorship if $|B| \in \{0, 2\}$; otherwise, we delete the alternatives $x \in A \setminus B$ that minimize $n_R(x)$ (if there is a tie, delete both alternatives) and return the outcome of the uniform random dictatorship for the reduced profile. As the following proposition demonstrates, f^2 indeed satisfies all axioms of Theorem 3 if $m = 3$.

Proposition 2. *For $m = 3$, f^2 satisfies anonymity, neutrality, PC -efficiency, and strict PC -participation.*

Proof sketch. The definition of f^2 immediately implies that this SDS is anonymous and neutral. Next, for proving that f^2 is PC -efficient, we consider the case distinction used in the definition of this SDS: if $|B| \in \{0, 2\}$, randomizing over the top-ranked alternatives is PC -efficient. On the other hand, if $|B| = 1$, f^2 is PC -efficient as it ignores one of the top-ranked alternatives. Finally, the strict PC -participation of f^2 follows from a tedious case distinction with respect to B . \square

Remark 2. Each of PC -efficiency and strict PC -participation is by itself compatible with anonymity and neutrality, as witnessed by maximal lotteries and the uniform random dictatorship, respectively. Hence, these two axioms are required for Theorem 3. In contrast, we do not know whether anonymity and neutrality are needed for this result. Proposition 2 shows that $m \geq 4$ is required.

Remark 3. Strict PC -participation is also incompatible with Condorcet-consistency, i.e., a statement analogous to Theorem 1 holds. This follows from the fact that a single voter cannot always change the Condorcet winner by joining the electorate, even if it is his least preferred outcome. In such cases, Condorcet-consistency implies that the outcome does not change while strict PC -participation requires the opposite.

4 Conclusion

We have studied incentive properties of social decision schemes (SDSs) based on the pairwise comparison (PC) lottery extension, and answered open questions raised by Brandt [2017] by proving three strong impossibilities. In particular, we showed that PC -strategyproofness and strict PC -participation are incompatible with PC -efficiency and Condorcet-consistency (see also Figure 1). We highlight three important aspects and consequences of our results. Firstly, when moving from the standard approach of stochastic dominance (SD) to PC , previously compatible axioms become incompatible. Secondly, our results show that—unlike with other classical impossibilities— PC does not help to circumvent the Gibbard-Satterthwaite theorem. Finally, our impossibilities identify a tradeoff between incentive-compatibility and efficiency. In light of this tradeoff, two SDSs seem particularly appealing: the uniform random dictatorship because it satisfies PC -strategyproofness and strict PC -participation, and maximal lotteries because it satisfies PC -strategyproofness in all profiles that admit a Condorcet winner, Condorcet-consistency, PC -efficiency, and PC -participation.

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